File Number 1130-30 Order Number GC26-5929-8





Systems Reference Library

IBM 1130 Subroutine Library

Ninth Edition (May 1974)

This is a reprint of GC26-5929-7 incorporating changes released in the following newsletter GN34-0182.

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This edition applies to version 2, modification 12 of IBM 1130 Disk Monitor System and to all subsequent modifications until otherwise indicated in new editions or Technical Newsletters. Changes are continually made to the specifications herein; before using this publication in connection with the operation of IBM systems, consult the latest SRL Newsletter, Order No. GN20-1130, for the editions that are applicable and current.

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Preface

The publication describes how the programmer can use the IBM 1130 library subroutines to increase the efficiency of his programs and decrease his writing and testing time. The libraries include the following programs:

- Interrupt Level Subroutines.
- Interrupt Service Subroutines.
- RPG Subroutines.
- FORTRAN I/O Subroutines.
- Data Code Conversion Subroutines.
- Arithmetic and Functional Subroutines.
- Selective Dump Subroutines.
- Utility Programs.

The subroutines are available for use with the 1130 Assembler, the 1130 FORTRAN Compiler, and the 1130 RPG Compiler. The Utility Programs are executable under Monitor control.

In Assembler language, the user calls the subroutines via a calling sequence. The appropriate subroutine calls are generated by the FORTRAN Compiler whenever a read, write, arithmetic, or CALL statement is encountered. The RPG Compiler generates the appropriate subroutine linkages. This publication describes each subroutine and the required calling sequence. All subroutines in the 1130 libraries are included in the lists that appear in Appendix A.

It is assumed that the reader is familiar with the methods of data handling and the functions of instructions in the IBM 1130 Computing System. He must also be familiar with the Assembler or Compiler used in conjunction with the subroutines. The following IBM publications provide the prerequisite information:

<u>IBM 1130 Functional Characteristics</u>, Order No. GA26-5881. IBM 1130 Operating System, Order No.

GA26-5717.

IBM 1130 Assembler Language, Order No. GC26-5927.

IBM 1130/1800 Assembler Language, Order No. GC26-3778.

IBM 1130/1800 Basic FORTRAN IV Language, Order No. GC26-3715.

IBM 1130 RPG Specifications, Order No. GC21-5002.

The operating procedures manuals for the programming systems also provide information on subroutine usage. These manuals are:

<u>IBM 1130 Card/Paper Tape Programming</u> <u>System Operator's Guide</u>, Order No. GC26-3629.

IBM 1130 Disk Monitor System, Version 2, Programmer's and Operator's Guide, Order No. GC26-3717.

MACHINE CONFIGURATION

The use of the library subroutines requires the following minimum machine configuration:

IBM 1131 Central Processing Unit with 4096 words of core storage.

IBM 1442 Card Read Punch, or IBM 1134 Paper Tape Reader with IBM 1055 Paper Tape Punch.

<u>Note</u>: RPG, available only with the DM2 system, requires 8192 words of core storage.

In addition, the following input/output units and features can be controlled by the input/output subroutines.

Console Printer/Keyboard Single Disk Storage 1132 Printer 1627 Plotter 1403 Printer (DM2 only) 2310 Disk Storage (DM2 only) 2311 Disk Storage (DM2 card system only) 2501 Card Reader (DM2 only) 1231 Optical Mark Page Reader (DM2 only) Synchronous Communications Adapter (DM2 only) Plotter subroutines are described in IBM

<u>1130/1800 Plotter Subroutines</u>, Order No. GC26-3755.

SCA subroutines are described in <u>IBM</u> <u>1130 Synchronous Communications Adapter</u> <u>Subroutines</u>, Order No. GC26-3706.

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It is often necessary to repeat a group, or block, of instructions many times during the execution of a program (examples include conversion of decimal values to equivalent binary values, computation of square roots, and the reading of data from a card reader). It is not necessary to write the instructions each time a function is required. Instead, the block of instructions is written once, and the main program transfers to that block each time it is required. Such a block of instructions is called a subroutine. Subroutines normally perform such basic functions that they can assist in the solution of many different kinds of problems.

When a main program uses a subroutine several times, which is the common situation, the block of instructions constituting the subroutine need appear only once. Control is transferred from a main program to the subroutine by a set of instructions known as a calling sequence, or basic linkage. A calling sequence transfers control to a subroutine and, through parameters, gives the subroutine any control information required.

The parameters of a calling sequence vary with the type of subroutine called. An input/output subroutine requires several parameters to identify an input/output device, storage area, amount of data to be transferred, etc., whereas an arithmetic/functional subroutine usually requires one parameter representing an argument. Each calling sequence used with subroutines in the 1130 system consists of a CALL or LIBF statement (whichever is required to call the specific subroutine), followed by the DC statements that make up the parameter list. The calling sequences for the various subroutines in the libraries are presented later in the manual. Each subroutine is self-contained, so that only those subroutines required by the current job are in core storage during execution.

In addition to the subroutines described in this publication, subroutines are available for use with the Disk Monitor System, Version 2 that are not provided in the system library for that version. These subroutines are contained in the following separately available programs:

- <u>Graphic Subroutine Package</u>, which enables the FORTRAN IV or Assembler language programmer to display images in the form of lines, points, and characters on the screen of a 2250 Model 4 Display Unit attached to the 1130 system. The program also provides for communication between the 2250 operator and the user's program. It is described in the publication <u>IBM</u> <u>1130/2250 Graphic Subroutine Package</u> for Basic FORTRAN IV, GC27-6934.
- Data Transmission Subroutines, which enable the FORTRAN IV or Assembler language programmer to transmit data between a program being processed by the Disk Monitor System Version 2 and a program being processed by a remote System/360 Operating System. These subroutines permit an 1130 program to use the high-speed computational capability and large storage capacity of the IBM System/360 Operating Communication between the two System. systems is accomplished in binary synchronous mode via telecommunication lines. The data transmission subroutines are described in the publication IBM System/360 Operating System and 1130 Disk Monitor System: System/360-1130 Data Transmission for FORTRAN, GC27-6937.
- <u>Satellite Graphic Job Processor</u>, which enables the user at a 2250 Model 4 Display Unit attached to the 1130 to easily start the processing of related programs in a remote System/360 Operating System. This allows the 2250 user to access the high-speed computational capability and large storage capacity of the IBM System/360 Operating System. Use of the Satellite Graphic Job Processor requires the data transmission subroutines discussed in the preceding paragraph. The Satellite Graphic Job Processor is described in the publication IBM System/360 Operating System and 1130 Disk Monitor System: User's Guide for Job Control from an IBM 2250 Display Unit Attached to an IBM 1130 System, GC27-6938.

Interrupt Service Subroutines

The interrupt service subroutines (ISSs) transfer data from and to the various input/output devices attached to the computer. These subroutines handle all the details peculiar to each device, including the usually complex interrupt functions, and can control many input/output devices at the same time by overlapping their operations.

ISS Characteristics

To fully understand subsequent descriptions of each ISS, the user should be familiar with the following characteristics, which are common to all ISSs:

- Methods of data transfer.
- Interrupt processing.
- ILS (interrupt level subroutine) operation.
- ISS (interrupt service subroutine) operation.
- General error-handling procedures.
- Basic calling sequence.

METHODS OF DATA TRANSFER

IBM 1130 I/O devices and their related subroutines can be differentiated according to their methods of transmitting and/or receiving data.

Direct Program Control

Serial I/O devices operate via direct program control, which requires a programmed I/O operation for each word or character transferred. A character interrupt occurs whenever a character I/O operation is completed. Direct program control of data transfer is used for the following system I/O devices: 1442 Card Read Punch, 1442 Card Punch, 1134 Paper Tape Reader and 1055 Paper Tape Punch, Console Printer, Keyboard, 1132 Printer, and 1627 Plotter.

Data Channel

Other system I/O devices operate via a data channel, which requires an I/O operation only to initiate data transfer. These devices are provided with control information, word counts, and data from the

user's I/O area. Once initiated, data transfer proceeds concurrently with program execution. An operation-complete interrupt signals the end of an I/O operation when all data has been transferred. All disk drives, the 1403 Printer, and the 2501 Card Reader operate via a data channel.

INTERRUPT PROCESSING

Interrupt processing is divided into two parts, level processing and device processing. The flow of logic in response to an interrupt is: user program interrupted, level processing begun, device processing begun and completed, level processing completed, and user program continued.

Level Processing

Level processing consists of selecting the correct device processing subroutine, performing certain housekeeping functions, and clearing the level by a BOSC instruction when interrupt processing is complete.

Level processing is done by the ILSs (interrupt level subroutines). Entered by interrupts, ILSs give temporary control to device processing subroutines (ISSs) and eventually return control to the user program. The interrupt entrance address is stored during the loading of a core load or program, in the appropriate interrupt branch address; location 8 for interrupt level zero (ILS00), location 9 for interrupt level one (ILS01),..., location 12 (/000C) for interrupt level four (ILSO4). The device processing entrance address is computed during the loading of a core load from identifying information that is a part of the ILS.

In the card/paper tape system, the device processing entrance address is stored during the loading of a program from identifying information stored in the ILS, in the compressed ISS header card, and in the loader interrupt Transfer Vector.

Device Processing

Device processing consists of operating an I/O device, processing the interrupts, and clearing the device by an XIO (sense DSW) instruction when interrupt processing is complete.

Device processing is done by the ISSs (interrupt service subroutines). The ISSs can be entered by a calling instruction (LIBF or CALL), which either requests certain initialization to be done or requests an I/O device operation. They can also be entered by an ILS as part of the interrupt processing. The calling entry point is specified in the ISS statement. The interrupt entry points are set up in the ISS and identified in the ILS. They are entered indirectly through a branch address table.

ILS OPERATION

The ISS/ILS package services all input/output interrupts.

Description

There is one ILS for each interrupt level used. Each subroutine determines which device on its level caused a particular interrupt; preserves the contents of the Accumulator, the Accumulator Extension, Index Register 1 (XR1), and the Carry and Overflow indicators; and transmits identifying information to the ISS. Disk Monitor ILSs also save Index Register 2 (XR2). The special ILSX subroutines in DM2 save and restore Index Register 3.

Interrupt service subroutines are loaded first so that the loader loads only the ILSs that are required. For example, if a main program does not call the 1132 Printer subroutine, the subroutine for interrupt level 1 (ILS01) need not be loaded because no interrupts will occur on that level. An ILS cannot be called; it is included in a core load or program only if requested by an ISS. If you use the 1130 Card/Paper Tape system, see "ISS-Define Interrupt Service Entry Point" in IBM 1130 Assembler Language. If you use the 1130 Disk Monitor, Version 2, system (DM2), see "Define Interrupt Service Subroutine Entry Point" in IBM 1130/1800 Assembler Language.

When the 1LSs are loaded, the core addresses assigned to them are incorporated into the appropriate locations in the Interrupt Transfer Vector (decimal words 8-13). Interrupts occurring during execution of a user program cause a Branch Indirect, via the interrupt branch address, to the correct ILS.

Recurrent Subroutine Entries

Recurrent entries to a subroutine can result from interrupts. For example, during execution of the Console Printer subroutine, a disk interrupt can start execution of a subroutine to handle the condition that caused the disk interrupt. If this handling includes calling the Console Printer subroutine, certain information is destroyed, the most important of which is the return address of the program that originally called the Console Printer.

To prevent the loss of data resulting from such a recurrent entry, the <u>user</u> must provide the programming required to save the return address and any other data needed to continue an interrupted subroutine after an interrupt has been serviced.

Note: All ISSs were written with the assumption that all LIBFs to them would be executed on the mainline level (i.e., not while on the interrupt level). There are <u>no</u> provisions in any ISSs to handle recurrent entries. See Appendix G for information on user-written re-enterable code.

ISS OPERATION

This section briefly describes the operation of the ISSs. This description, along with some basic flowcharts, should make it easier for the reader to understand the descriptions of individual subroutines presented later.

The disk subroutines are included here as ISSs even though in the Disk Monitor System they are not truly ISSs. They do however, have most of the characteristics of an ISS.

LSS Subdivision

Each ISS is divided into a call portion and an interrupt response portion. The call portion is entered when a user's calling sequence is executed; the interrupt response portion is entered as a result of an I/O interrupt.

Call Processing

The Interrupt Service Subroutines -- with the exception of those used by FORTRAN -save and restore the contents of the Accumulator, Index Registers 1 and 2, and the Carry and Overflow Indicators. However, the contents of the Accumulator will be destroyed if a preoperative error is detected. The call portion, illustrated in Figure 1, has four basic functions:

- Determines if any previous operations on the specified device are still in progress.
- 2. Checks the calling sequence for legality.
- 3. Saves the calling sequence.

4. Initiates the requested I/O operation.

The flow diagram (Figure 1) is not exact for any one ISS. It is only a general picture of the internal operation of the call portion of an ISS.

<u>Tetermine Status of Previous Operation</u>. This function can be performed by using a programmed subroutine-busy indicator to determine if a previous operation is complete. The CARD1 subroutine is a good example. When an operation is started on the 1442, a subsequent LIBF CARD1 for the 1442 is not honored until the subroutine-busy indicator is turned off. A call to any other ISS subroutine, such as TYPE0, is not affected by the fact that the CARD1 subroutine is busy.

Each ISS, except PAPTN and DISKN, can use one programmed subroutine-busy indicator to determine if a previous operation is complete. The PAPTN subroutine uses two busy indicators, one for the paper tape reader and one for the If an operation is started on the punch. reader, a subsequent LIBF PAPTN for the reader is not honored until the Reader Busy indicator is turned off. However, a LIBF PAPTN for the paper tape punch is treated in the same manner as a call to any other ISS and is not affected by the fact that the paper tape reader is busy. The subroutine DISKN uses five busy indicators, one for each disk drive. (Each disk drive corresponds to a certain bit in \$DBSY.) This provides the possibility to operate all of the disk drives simultaneously.

<u>Check Legality of Calling Sequence</u>. Calling sequences are checked for such items as illegal function character, illegal device identification code, zero or negative word count, etc.

<u>Save Calling Sequence</u>. The call portion saves, within itself, all of the calling sequence information needed to perform an I/O operation. The user can modify a calling sequence, even though an I/O operation is not yet complete.

<u>Note</u>: The I/O data area should be left intact during an operation because the ISS is continually accessing and modifying that area. <u>Initiate I/O Operation</u>. The call portion only initiates an I/O operation. Subsequent character interrupts cr operation complete interrupts are handled by the interrupt response routine.

Interrupt Response Processing

The I/O interrupt response portion of an ISS is illustrated in Figure 2.

Operation. An I/O interrupt causes a user program to exit to an interrupt level subroutine, which in turn exits to the I/O interrupt response portion of an ISS. The interrupt response portion checks for errors, does any necessary data manipulation, initiates character operations, and initiates retry operations in case of errors. It then returns control to the interrupt level subroutine, which returns control to the user.

<u>Character Interrupts</u>. These interrupts occur for devices under direct program control whenever data can be read or written, e.g., a card column punched or a paper tape character read.

Operation Complete Interrupts. These interrupts occur in disk and card operations when a specified block of data has been read or written, e.g., a disk record read.

Error Detection and Recovery Procedures. These procedures are an important part of an ISS. However, little can be done about reinitiating an operation until a character interrupt or operation complete interrupt occurs. Therefore, error indicators are not examined until one of these interrupts occurs.

<u>Recoverable Device</u>. This is an I/O device that can be easily repositioned by a subroutine or by an operator and an I/O operation reinitiated. If a device is not recoverable, or if an error cannot be corrected after a specified number of retries, the user is informed of the error condition. If a device is recoverable, the user may request, via his error subroutine, that the operation be reinitiated.





Figure 2. Interrupt Response Portion of an ISS

• Figure 1. Call Portion of an ISS

GENERAL ERROR-HANDLING PROCEDURES

Each ISS has its own error detecting portion, which determines the type of error and chooses an error procedure. (In this context, the term error includes such conditions as last card, channel 9, channel 12, etc.) Errors fall into one of two categories: those that are detected before an I/O operation is initiated, and those that are detected after an I/O operation has been initiated. Appendix B contains a list of the errors detected by the ISSs; Appendix C contains descriptions of the actions taken by each ISS after the return from user-written error subroutines.

Preoperative Error Detection

Before an ISS initiates an I/O operation, it checks the device status and the legality of calling parameters. If a device is not ready or a parameter is in error, the Interrupt Service Subroutine will wait at \$PRET+1 displaying an error indicator that defines the error (see Appendix E). This error indicator consists of four hexadecimal digits that are defined below.

\$PRET is entered via a Branch and Store Instruction Counter (BSI) instruction in the following subroutines: DISKZ, DISK1, DISKN, OMPR1, PLOTX, and the ISSs used by FORTRAN. All other ISSs store the address of the LIBF statement in \$PRET and then branch to \$PRET+1 to wait and display the error; i.e., when PROGRAM START is pressed, the call to the subroutine is retried.

Digit 1 identifies the ISS subroutine called:

C/PT System	DM2 System
1-CARD0 or CARD1	1- CARDO, CARD1, or CARDZ; PNCH0, PNCH1, or PNCHZ
2-TYPE0 or WRTY0	2- TYPEO or TYPEZ, WRTYO, or WRTYZ
3-PAPE1 or PAPTN	3- PAPT1, PAPTN, PAPTX or PAPTZ
	4- READO, READ1, or READZ
5-DISK0, DISK1, or DISKN	5- DISKZ, DISK1, or DISKN
6- PRNT 1	6- PRNT1 or PRNTZ

7-PLOT1

8-SCAT1, SCAT2 or SCAT3

9- PRNT3 or PRNZ

A- OMPR1

Digits 2 and 3 are reserved.

- Digit 4 identifies the error:
 - 0- device not ready.
 - 1- illegal LIBF parameter or illegal specification in the I/O area.

There is a WAIT instruction in core location \$PRET+1 and a branch instruction (BSC I \$PRET) in the next location. Therefore, the LIBF may be executed again (after the error condition has been corrected) by pressing PROGRAM START on the console. The user can, if he chooses, replace these two instructions with an exit to his own error subroutine.

Postoperative Error Detection

After an I/O operation has been started, certain conditions may be detected about which the user should be informed. The conditions might be card jams for which manual intervention is needed before the operation can continue; read checks that have not been corrected after a specified number of retries; or indications of equipment readiness, such as last card or channel 12 indicators. All these conditions are detected by the interrupt response portion (see "ISS Operation").

No Error Parameter. If no error parameter is included in the calling sequence that initiated the I/O operation and a postoperative error condition is detected, the card/paper tape system subroutine initiates a Wait procedure (programmed loop), which continues until the operator corrects the detected condition.

The DM2 system does not use a programmed loop, but rather branches to a postoperative error trap that is similar to the preoperative error trap. Each interrupt level (1-4) has its own postoperative error trap with accompanying WAIT address.

Level	1	-	\$PST1	(0081)
Level	2	÷	\$PST2	(0085)
Level	3	-	\$PST3	(0089)
Level	4	-	\$PST4	(008D)

Processing resumes -- at the address immediately following -- after the operator corrects the detected condition and presses PROGRAM START.

Error Parameter Included. If an error parameter is included in the calling sequence, a Branch and Store Instruction Counter (BSI) instruction to the user's error subroutine specified in the calling sequence is executed. Identifying information is placed in the Accumulator and Extension (see Appendix B). When the user's error subroutine returns control to the ISS using the return link (see "Basic ISS Calling Sequence"), the subroutine examines the Accumulator. If the user has cleared the Accumulator before returning to the subroutine, he is requesting that the error condition be ignored and the operation terminated. If the user has not cleared the Accumulator, he is requesting that the operation be restarted, in which case the subroutine reinitiates the operation before returning to the user's main program.

<u>Implications of the User's Error</u> <u>Subroutine</u>. It is important to note that a user's error subroutine (entered via the LIBF error parameter address) is executed as part of the interrupt processing. The interrupt level is still on, preventing recognition of other interrupts of the same or lower priority. This has the following implications:

- Return must be made to the ISS subroutine via the return link (set up by the BSI instruction executed by the ISS subroutine). Otherwise, normal processing cannot be continued because the ISS must return to the ILS to restore the contents of the Accumulator and Extension, Status Indicators, and Index Registers.
- 2. Return must be made with a BSC instruction, not a BOSC instruction. Otherwise, the interrupt level is turned off, setting up the possibility that another interrupt could occur on the same level, thus destroying the return address to the user from the ILS.
- 3. A LIBF or CALL to another subroutine from the user's error subroutine can cause a recurrent-entry problem. If that subroutine is already in use when the interrupt occurs, the user's new LIBF or CALL destroys the original return address and disrupts operation of the called subroutine.
- 4. A LIBF or CALL to another ISS can cause an endless loop if the new I/O device operates on the same or lower

priority interrupt level than the device that caused the error.

Note: A call to WRTYO to type an error message can be made only if the user does not wait for the completion of typing or for operator intervention before returning to the ISS. A test loop on level 4 (typewriter) or a WAIT loop will both block the clearing of the level that caused the interrupt to the user's error subroutine.

5. The user should have a separate error subroutine for each device to prevent errors on several devices (on different levels) from causing recurrent-entry problems in the _ser's error subroutine.

<u>Note</u>: The error codes in the Accumulator may not distinguish between ISSs, as the preoperative error codes do.

Since the ILS saves Index Register 1 as part of its interrupt processing, the user's error subroutine can also use this index register without saving and restoring it. However, the user cannot depend on the contents of Index Register 1 unless he initializes it as part of his error subroutine. The DM2 ILSs also save Index Register 2. The special ILSX subroutines in the DM2 save and restore Index Register 3.

<u>Programming Techniques - User's Error</u> <u>Subroutine Exits</u>. Some programming techniques that can be used in conjunction with the ISS error exit are as follows:

1. To try the operation again:

Label	Opera	ation	F	т	Operands & Remarks
<u>n 8</u>	.·	ر	\mathbf{v}_{i}	11	5. 40 45 50
US ER.	D.C.				0
	B.S.	C,	τ		USER .

2. To terminate the operation:

Label	Operation	FT			Operands & Rema	rks	
·		9 0	n 4	*		·	
U.S.E.R.	e.c.		6				
	S.RA.		16	. , T.O.	CLEAR T	HELACOUN	ULATOR .
Level 1	B.S.C.	I	USER				
	1		Luna				<u></u>

3. To indicate that a condition ("last card" or "channel 9") was detected and that the normal program flow should be altered:

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Label	Operation	1	1	11	Operands & Remarks
· · · ·	11 V	.,	: 11		a <u>a a a</u> <u>a</u> <u>a</u>
	L.D.	-		Ц	I.N.D.I.C.
لمتصب	B.S.C.	1		Ц	NEW, Z ALLTER PROGRAM FLOOM
	LIBE		1		CIARD 1
l	D.C.				1.1.0.0.0.0
	lo.c				LINPUT READ ONE CARD
1	O.C.		1		U.S.E.R
	·		Т	Π	•••••••••••••••••••••••••••••••••••••••
	•		Т	Π	•••••••••••••••••••••••••••••••••••••••
			Т	Π	•••••••••••••••••••••••••••••••••••••••
I.S.F.R.	D.C.			П	
	B.S.C.	,	-	П	U.S.E.B 7.
	L.O.			П	DØØØ1.
	S.T. Q		1		IND.T.C.
EX.I.T.	8.S.C.	,	·		U.S.F.B.
	••••	T	1		•
LNDIC	D.C.		П		<u> </u>
					•••••••••••••••••••••••••••••••••••••••
V.E.W.	S.R.A.		11		1.6
	5.7.0		П	Π	
		+	Н	H	**************************************
		1	H		••••••••••••••••••••••••••••••••••••••
		-	+ +		The hard and the second of the

BASIC ISS CALLING SEQUENCE

Each ISS described in this manual is entered via a calling sequence. These calling sequences follow a basic pattern. In order not to burden the reader with redundant descriptions, this section presents the basic calling sequences and describes those parameters that are common to most of the subroutines.

Basic Calling Sequence

LIBF	Name
DC	Control parameter
DC	I/O area
DC	Error subroutine

The above calling sequence, with the parameters shown, is basic to most of the ISSs. Detailed descriptions of the above four parameters are omitted when the subroutines are described later in the manual. Unless otherwise specified, the subroutine returns control to the instruction immediately following the last parameter.

Name Parameter

Each subroutine has a symbolic name that must be written in the LIBF statement exactly as listed in Figures 3 and 4.

Device	Subroutine
1442 Card Read Punch	CARDO, CARD1, or CARDZ
Disk	DISKO, DISK1 DISKN
1132 Printer	PRNT 1 OF PRNTZ
Keyboard/Console Printer	TYPEO or TYPEZ
Console Printer	WRTYO or WRTYZ
1134/1055 Paper Tape	PAPT1, PAPTN or PAPTZ
1627 Plotter	PLOT 1 or PLOTX
Synchr. Comm. Adapter	SCAT1, SCAT2, or SCAT3

Figure 3. C/PT System ISS Names

For some devices more than one subroutine is available, although only one can be selected for use in any one program (including called subroutines).

NAMEO. The NAMEO subroutine is the shortest and least complicated. The NAMEO version is the standard subroutine for the 1442, 2501, and Console Printer/Keyboard. The NAMEO version of the Disk routine (DISKO) can be used if transfer of data is 320 words or less (C/PT system only).

Device	Subroutine
1442 Card Read Punch	CARDZ, CARDO, or CARD1
2501 Card Reader	READZ, READO, or READ1
1442 Card Punch	PNCHZ, PNCHO or PNCH1
Disk	DISKZ, DISK1, or DISKN
1132 Printer	PRNTZ, PRNT1, or PRNT2
1403 Printer	PRNZ, Or PRNT3
Keyboard/Console Printer	TYPEZ, or TYPE0
Console Printer	WRTYZ, or WRTYO
1134/1055 Paper Tape Reader Punch	PAPTZ, PAPT1, PAPTN, or PAPTX
1627 Plotter	PLOT1, or PLOTX
1231 Optical Mark Page Reader	OMPR 1
Synchr. Comm. Adapter	SCAT1, SCAT2, or SCAT3
2250 Display Unit, Model 4	DSP YN

Figure 4. DM2 System ISS Names

NAME1. The NAME1 version is the standard subroutine for the disk, 1132, 1403, 2501, 1134/1055, 1231, and 1627. It may be used if a user error exit is needed rather than the internal looping and retries by the NAME0 subroutine.

NAMEN. The NAMEN version is available to operate the 1134/1055 Paper Tape Reader and Punch simultaneously and to minimize extra disk revolutions when transferring more than 320 words to or from the disk. The NAMEN subroutine offers more options than the NAME1 subroutine. In DM2, it also operates as many as 5 disks simultaneously.

<u>NAMEZ</u>. The NAMEZ version is designed for use in an error-free environment. It provides no preoperative parameter checking. The FORTRAN formatting subroutines use these ISSs but they do <u>not</u> use the calling sequence listed below (see "Subroutines Used by FORTRAN").

<u>PRNT2</u>. The PRNT2 version is used when the 1132 is used with the SCA.

PRNT3. The PRNT3 version is used with the 1403.

Control Parameter

The control parameter, in the form of four hexadecimal digits, conveys necessary control data to the ISSs by specifying the desired function (read, write, etc.), the device identification, and similar control information. Most subroutines do not use all four digits.

A typical control parameter is illustrated below.

Hexadecimal Digits

		1st	2nd	3rd	4th
		4	ŧ	ŧ	+
1/0	Function		-		1
Not	Used				

Device Identification ____

Since the I/O function and device identification digits are used in most subroutines, a description of the purpose of each is given here.

I/O Function

The function digit in the calling sequence specifies which I/O operation the user is requesting. Three of these functions-read, write, and test-- are used in most subroutines.

<u>Read</u>. The read function causes a specified amount of data to be read from an input device and placed in a specified input area. Depending upon the device, an interrupt signals the subroutine either when the next character is ready or when all requested data has been read. When the specified number of characters has been read, the subroutine becomes available for another call to that device.

<u>Write</u>. The write function causes a specified amount of data from the user's output area to be written, i.e., printed or punched, by an output device. As with the read function, an interrupt signals the subroutine when the device can accept another character, or when all characters have been written. When the specified number of characters has been written, the subroutine becomes available for another call to that device.

<u>Test</u>. The test function causes a check to be made as to the status of a previous operation initiated on an I/O device. If

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the previous operation has been completed, the subroutine branches to the LIBF+3 core location; if the previous operation has not been completed, the subroutine branches to the LIBF+2 core location. The test function is illustrated below:

LIBF Name

LIBF+1 DC Control Parameter (specifying Test function)

LIBF+2 OP Code xxxx....

LIBF+3 OP Code xxxx....

<u>Note</u>: Specifying the test function requires two statements (one LIBF and one DC), except in disk subroutines, where three statements are required.

The test function is useful in situations in which input data has been requested, but no processing can be done until the data is available.

Device Identification

This digit should be zero except for the Test function with the PAPTN (paper tape) subroutine.

<u>Note</u>: For all disk subroutines, this digit appears in the I/O area rather than in the control parameter.

I/O Area Parameter

The I/O area for a particular operation consists of one table of control information and data. This table is composed of a data area preceded by a control word (two control words for disk operations) that specifies how much data is to be transferred. The area parameter in the calling sequence is the address (symbolic or actual) of the first control word that precedes the data area.

The control word contains a word count referring to the number of data words in the table. It is important to remember that the number of words in the table is not always the number of characters to be read or written, because some codes pack two characters per word. The disk subroutines require a second control word, which is described along with those subroutines.

Error Parameter

TTDD

BUR ME

The error parameter is the means by which an ISS can give temporary control to the user in the event of conditions such as error, last card, etc. This parameter is not required for the NAMEO subroutines for the 2501, 1442, Console Printer, or Keyboard. The instruction sequence for setting up the error subroutine is shown below.

•	•	•
1	DC	ERROR (error parameter)
ERROR I	DC	0 (return link)
•	RSC T	. (error routine) EREOR (branch to return link)

The return link is the address in the related ISS to which control must be returned upon completion of the error subroutine. The link is inserted in location ERROR by a BSI from the ISS when the subroutine branches to the error subroutine.

The types of errors that cause a branch to the error address are listed in Appendix B.

<u>Note</u>: The user's error subroutine is executed as part of the interrupt response handling. The interrupt level is still on and remains on until control is returned to the ISS (see "General Error-Handling Procedures").

Assignment of Core Storage Locations (C/PT System)

The portion of core storage used by the ISS and ILS subroutines is defined below. Care should be used in altering any of these locations (see Figure 5).

The areas illustrated in Figure 5 are described below.

Interrupt Branch Addresses

<u>ILS Subroutines</u>. When required, the address of ILS00 is always stored in location 8, ILS01 in location 9,..., ILS05 in location 13 (/000D).

<u>Interrupt Trap</u>. The address of the interrupt trap is stored in any location for which no ILS is loaded.

1132 Printer

This area is used by 1132 Printer.

Preoperative Error Trap

This exit is used whenever a preoperative error (illegal LIBF or device not ready) is detected by an ISS. To retry the call, press START.

ISS Exit

The ISS exit results from pressing the Keyboard Interrupt Request key. The TYPEO, and WRTYO subroutines execute a BSI 1 /002C whenever a keyboard operator request is detected. Note that interrupt level 4 is still on.

The user-written subroutine must return to the TYPEO or WRTYO subroutine in order to allow interrupts of equal or lower priority to occur. Also a call executed to any subroutine might cause a recurrententry problem unless the user can guarantee that the subroutine was not in use when the keyboard interrupt occurred.

Location /002C is initialized with the address of the interrupt trap in case the user fails to store an address in the interrupt trap to process Keyboard operator requests.

Interrupt Trap

This trap is entered when an interrupt occurs for which there is no ILS and/or no ISS assigned to the pertinent bit in the Interrupt Level Status Word (ILSW).

Interrupts of higher priority will be processed before the system finally halts with the IAR displaying /002F.

ISS Counter

The ISS counter is incremented by +1 every time an ISS initiates an interrupt-causing I/O operation and is decremented by +1 when the operation is complete. A positive value in this location indicates the number of interrupt (s) pending. This counter should never be negative.



Figure 5. ISS and ILS Core Locations for the C/PT System

Assignment of Core Storage Locations (DM2 System)

The portion of core storage used by the ISS and ILS subroutines is defined below. Care should be used in altering any of these locations (see Figure 6).

The areas illustrated in Figure 6 are described below.

Interrupt Branch Addresses

<u>ILS Subroutines</u>. The address of ILS00 is always stored in location 8, ILS01 in location 9,..., ILS05 in location decimal 13.

Interrupt Trap. The address of the Program Stop Key trap (\$STOP-location /0091) is stored in any location for which no ILS is loaded.

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<u>Reserved Areas</u>

These locations are reserved for the DM2 system.

1132 Printer

This area is used by 1132 Printer.

Preoperative Error Trap

This exit is used whenever a preoperative error (illegal LIBF or device not ready) is detected by an ISS. To retry the call, press START.

1SS Counter

The ISS counter is incremented by +1 every time an ISS initiates an interrupt-causing I/O operation and is decremented by +1 when the operation is complete.

A positive value in this location indicates the number of interrupts (s) pending. This counter should never be negative.

Interrupt Request Branch Address

The subroutine ILS04 or ILSX4 executes a BSI I \$IREQ whenever a Keyboard operator request is detected.

\$IREQ (location /002C) is initialized with the address \$I420 in Resident Monitor. This allows the user to terminate the job by pressing the Interrupt Request key (INT REQ).

Note the following when writing an interrupt request subroutine:

- Interrupt level 4 is still on.
- An XIO instruction sensing Keyboard with reset must be performed.
- Return to ILSO4 or ILSX4 to exit address +6.
- ILSO4 or ILSX4 will turn off the interrupt.
- Subroutines that are in use when the interrupt occurs may not be called.

For examples of INT REQ see <u>IBM 1130</u> <u>Disk Monitor System</u>, Version 2, <u>Programmer's</u> and Operator's Guide.

Hex	Decimo	
8 9 8 0 0	8 9 10 11 12 13	(ILS00) (ILS01) (ILS02) (ILS03) (ILS03) (ILS04) (ILS05)
E	14	Reserved for Monitor System
1F	31	
20	32	Reserved for 1132 Printer
27	39	
28 29 2A	40 41 42	SPRET DC *-* WAI1 BSC I SPRET
2Ċ	44	\$IREQ DC + - + } Interrupt Request Branch Address
2D	45	T T Reserved for Monitor System
31	49	
32	50	SIOCT DC * - * }ISS Counter
33	51	Reserved for Monitor System
3E	62	
40	64	Reserved for Monitor System
80	128	
82 83	130 131	WAIT BSC 1 \$PST1 Postoperative Error Trap for Level 1
85 86 87	133 134 135	\$PST2 DC * - * WAIT BSC 1 \$PST2 }
89 8A 8B	137 138 139	SPST3 DC * - * WAIT BSC I SPST3 Postoperative Error Trap for Level 3
8D 8E 8F	141 142 143	\$P\$T4 DC * - * WAIT BSC I BSC I \$P\$T4
91 92 93	145 146 147	SSTOP DC * - * WAIT BOSC I SSTOP

Figure 6. ISS and ILS Core Locations for the DM2 System

Postoperative Error Traps

These traps are entered when a devicenot-ready condition is detected prior to the initiation of an I/O operation in the interrupt response portion of an ISS subroutine. Each interrupt level (1-4) has its own postoperative error trap. The system will WAIT with the IAR displaying the address of \$PST1+2, \$PST2+2, \$PST3+2, or \$PST4+2, depending on the interrupt level of the device.

Description of Interrupt Service Subroutines

Note that the subroutine READ0, READ1, PNCH0, PNCH1, PRNT3, and OMPR1 are available only with the DM2 system.

1442 CARD READ PUNCH SUBROUTINES (CARDO AND CARD 1)

The card subroutines perform all I/O functions relative to the IBM 1442 Card Read Punch: read, punch, feed, and stacker select.

<u>CARDO</u> <u>Subroutine</u>. The CARDO subroutine is shorter and less complicated than CARD1 and is the standard subroutine for the 1442.

CARDO can be used if the error parameter is not needed. When an error occurs, the subroutine loops (DM1 and C/PT system) or will WAIT at \$PST4+1 (DM2 system) until the operator takes corrective action. Last card conditions cause preoperative not-ready exits.

<u>CARD1 Subroutine</u>. The CARD1 subroutine can be used for the Card Read Punch if a user error exit is neeled, rather than the error procedures of the CARD0 subroutine.

Calling Sequence

Label	Operation	F	T				Operands &	Remarks	
.) 15	27 30	11	- 13		35 40	45	50	55	0a
لىبىت	LIBF				CARD AL	L. C.A.L	LICIAR	0.110	<u></u>
	0.C				/1b.Ø.Ø.Ø.	C.O.N	TROLL	P.A.R.A.N	ETER
h	0.0				I QAR	1.1.1.0	AREA	PARA	METER
	D.C.			L	EIRIRIOR .	. ERR	QIR PAN	RAMET	ER
	<u></u>								
	<u></u>								
ERRIGIR	D _i C _i			L	<u>*</u>	BIEIT	URMAN	DIDIRIES	יייי צי
	· · · · ·	-	\downarrow	L					
	· · · · ·	_	╞				· · · · · · · ·		
	<u></u>		+-+	┝			. <u></u> .		<u> </u>
	BISIC	I	4		ERRIDIR	RET	URNIT	QC.A.L	1.L.E.R
	•••••		┢	-				_ ب ا	
	•••••		┢	⊢	++				
لىبيە	••••		+	┢					
I.O.A.R.	C	+	┢	⊢	firme	WOR	יימי יכיסימי	NITLL	بتعتب
	B.S.S.		+	ŀ.	hurre	I /.Q	AREA		

where

a is 0 or 1,

b is the I/O function digit,

f is the number of columns to be read from or punched into the card,

h is the length of the I/O area. h must be equal to or greater than f.

The calling sequence parameters are described in the following paragraphs.

Control Parameter

This parameter consists of four hexadecimal digits as shown below:

	· · · · · · · · · · · · · · · · · · ·	1 3	2	3	4
		•		• .	ŧ
1/0	Function				
			١		
Not	Used				

I/C Function

The I/O function digit specifies the particular operation to be performed on the 1442 Card Read Funch. The functions, associated digital values, and required parameters are listed and described below.

Function	Digita <u>Value</u>	l Required <u>Parameters</u> ¹	
Test	0	Control	
Read	1	Control, I/O Area, Error ²	
Punch	2	Control, I/O Area, Error ²	
Feed	3	Control, Error ²	
Stacker Select	4	Control	
¹ Any parameter not required for a particular function must be omitted. ² Error parameter not required for CAREO.			

<u>Test</u>. Branches to LIBF+2 if the previous operation has not been completed, to LIBF+3 if the previous operation has been completed.

<u>Read</u>. Reads one card and transfers a specified number of columns of data to the user's input area. The number of columns read (1-80) is specified by the user in the first location of the I/O area. The

subroutine clears the remainder of the I/O area and stores a 1 in bit position 15 of each word, initiates the card operation, and returns control to the user's program. When each column is ready to be read, a column interrupt occurs. This permits the card subroutine to read the data from that column into the user's input area (clearing bit 15), after which the user's program is again resumed. This sequence of events is repeated until the requested number of columns has been read, after which the remaining column interrupts are cleared (no data read).

When an operation complete interrupt occurs, the card subroutine checks for errors, informs the user if an error occurred (CARD1 only), and sets up to terminate (CARD1 only) or retry the operation.

The data in the user's input area is in a code identical to IBM Card Code format; that is, each 12-bit column image is left-justified in one 16-bit word.

<u>Punch</u>. Punches into card the number of columns of data specified by the word count found at the beginning of the user's output area. The punch operation is similar to the read operation. As each column comes under the punch dies, a column interrupt occurs; the card subroutine transfers a word from the user's output area to the punch and then returns control to the user's program.

This sequence is repeated until the requested number of columns has been punched, after which an Operation Complete interrupt occurs. At this time the card subroutine checks for errors, informs the user if an error occurred (CARD1 only), and sets up to terminate (CARD1 only) or retry the operation. The character punched is the image of the leftmost 12 bits in the word.

<u>Feed</u>. Initiates a card feed cycle. This advances all cards in the machine to the next station, i.e., a card at the punch station advances to the stacker, a card at the read station advances to the punch station, and a card in the hopper advances to the read station. No data is read or punched as a result of a feed operation and no column interrupts occur. This effectively skips a card when used in conjuction with a Read or Punch function.

When the card advance is complete, an Operation Complete interrupt occurs. At this time the card subroutine checks for errors, informs the user if an error occurred (CARD1 only), and sets up to terminate (CARD1 only) or retry the operation. <u>Stacker Select</u>. Selects stacker 2 for the card currently at the punch station. After the card passes the punch staticn, it is directed to stacker 2.

I/O Area Parameter

The 1/0 area parameter is the label of the control word that precedes the user's 1/0 area. The control word consists of a word count that specifies the number of columns of data to be read or punched, always starting the count at column 1. The word count must be in the range of 1-80.

Error Parameter

<u>CARDO</u>. CARDO has no error parameter. If an error is detected while an operation complete interrupt is being processed, the subroutine loops (C/PT system) or will WAIT at \$PST4+1 (DM2) with interrupt level 4 on, waiting for operator intervention. When the condition has been corrected, the 1442 made ready, and PROGRAM START pressed, the subroutine retries the operation.

<u>CARD1</u>. CARD1 has an error parameter. If an error is detected, the user can request the subroutine to terminate (clear subroutine-busy indicator and the interrupt level) or to loop (C/PT system) or WAIT at \$PST4+1 (DM2) for operator intervention (interrupt level 4 on). See "Basic Calling Sequence".)

Protection of Input Data

Since the CARD subroutines read data directly into the user's I/O area, the user can manipulate the data before the entire card has been processed. This procedure is inherently dangerous because, if an error occurs, the data may be in error and error-recovery procedures will cause the operation to be tried again. The exit via the error parameter is the only method of informing the user that an error has occurred. Therefore, do not manipulate data before the entire card has been processed when using CARDO.

When using CARD1, the following precautions should be taken:

 Do not store converted data back into the read-in area.

- Do not take any irretrievable action based on the data until the card has been read correctly; i.e., be prepared to convert the data or perform the calculations a second time.
- when data manipulation is complete, check the user-assigned error indicator that is set when a branch to the user-written error subroutine occurs. The data conversion or calculations can then be reinitiated, if necessary.

Last Card

When the last card has been detected, a branch to the user error routine with /0000 in the Accumulator will occur. An operation requested after the last card has been fed from the hopper causes an exit to \$PRET. When the 1442 is made ready and the PROGRAM START key is pressed, the last card will be processed.

2501 CARD READER SUBROUTINES (READ 0 AND READ1)

These card subroutines, available only with the DM2 system, perform read and test functions relative to the IBM 2501 card reader.

<u>READO Subroutine</u>. READO is shorter than READ1, provides no error parameter, and is the standard subroutine for operation of the 2501 card reader. On an error, READO branches to \$PSTY, then a WAIT for operation intervention will occur. The last card condition causes a branch to \$PRET.

<u>READ1 Subroutine</u>. READ1 is used for operation of the 2501 card reader if a user error exit is required.

Calling Sequence

Lavet	Operation	F	T	Operands & Remarks
	L.I.B.F	T		R.E.A.D.A. , C.A.L.L. C.A.R.D. I.N.P.U.T.
	D.C.	_		1 16. 0. 0. 0 C.O.N.T.R.O.L P.A.R.A.ME.T.E.R.
	P.C.			I.O.A.R
	0.0			ERROR ERBOR PARAMETER
	- ! · · · ·	-		, , , , , , , , , , , , , , , , , , , , ,
	┶╍╼┟		-]
		+-		
ERROR	0.6	4-1		#
	·		-	_
		_	-	· · · · · · · · · · · · · · · · · · ·
	1	4		<u> </u>
<u> </u>	B.S.C.	I	-	EIRIRIOR RIEITURNI TO CALLER.
	1			
	·			
		+	+	
I.O.A.R.	D.C.	-		Fr
<u> </u>	BSS.	+	4	be
	1			

where

a is 0 or 1,

b is the I/O function digit,

f is the number of columns to be read from the card,

h is the length of the I/O area. h must be equal to or greater than f.

The calling sequence parameters are described in the following paragraphs.

Control Parameter

This parameter consists of four hexadecimal digits as shown below:

	1		3 4	•
	+	1	} 4	
I/0	Function			
			1	
Not	Used			

I/O Function

The I/O function digit specifies the particular operation to be performed on the 2501 Card Reader. The functions, associated digital values, and required parameters are listed and described below.

Function	Digital <u>Value</u>	Required Parameters ¹		
Test Read	0 0	Control Control, 1/0 Area, Error ²		
¹ Any parameter not required for a particular function must be omitted. ² The error parameter is not required for READO.				

<u>Test</u>. Branches to LIBF+2 if the previous operation has not been completed, to LIBF+3 if the previous operation has been completed.

<u>Read</u>. Reads one card and transfers a specified number of columns of data to the user's input area. The number of columns read (1-80) is specified by the user in the first location of the input area. The subroutine initiates the read function and returns control to the user's program.

When an Operation Complete interrupt occurs, the card subroutine checks for errors. If an error occurred, READO exits to \$PST4; READ1 informs the user of the error and sets up to terminate or retry the operation.

The data in the user's input area is in IBM Card Code format; that is, each 12-bit column image is left-justified in one 16-bit word.

There is no separate feed function. However, a feed can be obtained by a read function with a word count of zero.

I/O Area Parameter

The I/O area parameter is the label on the control word that precedes the user's input area. The control word consists of a word count that specifies the number of columns of data to be read, always starting with column 1. The word count must be in the range of 0-80.

Error Parameter

<u>READO</u>. READO has no error parameter. If an error is detected while an Operation Complete interrupt is being processed, the subroutine branches to \$PST4, with interrupt level 4 on, waiting for operator intervention. When the condition has been corrected, the 2501 made ready, and PROGRAM START pressed, the subroutine attempts the operation again.

<u>READ1</u>. READ1 has an error parameter. If an error is detected, the user can request the subroutine to terminate (that is, to clear the subroutine's busy indicator and turn off the interrupt level) or retry. Prior to a retry, the subroutine checks to see if the unit is ready. If the unit is not ready, the subroutine branches to \$PST4 with interrupt level 4 on, waiting for operator intervention.

Last Card

A read function requested after the last card has been fed from the hopper causes an exit to \$PRET. When the reader is made ready and the PROGRAM START key pressed, the last card is read and fed into the stacker.

1442 CARD PUNCH SUBROUTINES (PNCHO AND PNCH1)

These card subroutines, available only with the DM2 system, perform all I/O functions relative to the IBM 1442-5 Card Punch, that is, punch and feed. These subroutines may also be used with the 1442-6 or 1442-7 Card Read Punch for punch and feed functions.

<u>PNCHO</u>. The PNCHO subroutine is shorter than PNCH1, provides no error parameter, and is the standard subroutine for operation of the 1442 card punch. On an error, PNCHO branches to \$PST4, then a WAIT for operator intervention will occur. The last card condition causes a branch to \$PRET.

<u>PNCH1</u>. PNCH1 can be used for operation of the 1442 card punch if a user error exit is desired.

Calling Sequence

Lotel	Operation	FT	Operands & Remarks
		1.1.1	· · · · · · · · · · · · · · · · · · ·
	LIBE	\downarrow	PINCHE , CALL CARD OUTPUT.
	D.C.		VID. O.O.O.N.T.R.O.L. PARAMETER.
	D.C.		I.O.A.R
	p.c.	┥┥╷	E.R.R.O.RER.R.O.R. P.A.R.A.M.E.T.E.R.
	<u></u>	<u>_</u>	<u></u>
		\downarrow \downarrow \downarrow	<u></u>
	_ <u>-</u>		
ERROR	D.C	\rightarrow	#1-# , , , , , , , , , , , RIEITIURINI ADDRIEISISI , , , , ,
	1		<u></u>
· ·	_ <u>_</u> !	+++	
	· · · · · · · · · · · · · · · · · · ·	+++	<u></u>
<u> </u>	B.S.C.	I	ERROR RETURN TO CALLER .
	··· · · · · · · · · · · · · · · · · ·	+++	
			<u></u>
		+	
I.O.A.R.	p.c	+++	FL WIDIR D. COULINT
	8.5.5	<u></u>	he contracted Julia, A.R.E.A
			<u></u>

where

a is 0 or 1,

b is the I/O function digit,

f is the number of columns to be punched into the card,

h is the length of the I/O area. h must be equal to or greater than f.

The calling sequence parameters are described in the following paragraphs.

Control Parameter

This parameter consists of four hexadecimal digits as shown below:

	1	2	2 3	34	
	•	4	t i	t t	
1/0	Function				
Not	used				

I/O Function

The I/O function digit specifies the particular operation to be performed on the 1442 Card Punch. The functions, associated digital values, and required parameters are listed and described below.

Function	Digita <u>Value</u>	l Required <u>Parameters</u> ⁴		
Test Punch Feed	0 2 3	Control Control I/O Area, Error ² Control, Error ²		
⁴ Any parameter not required for a particular function must be omitted. ² The error parameter is not required for PNCH0.				

<u>Test</u>. Branches to LIBF+2 if the previous operation has not been completed, to LIBF+3 if the previous operation has been completed.

<u>Punch</u>. Punches into one card the number of columns of data specified by the word count found at the beginning of the user's output area. As each column comes under the punch dies, a column interrupt occurs, the subroutine transfers a word from the user's output area to the punch, and then returns control to the user's program. The character punched is the image of the leftmost 12 bits in the word.

This sequence is repeated until the requested number of columns has been punched, after which an Operation Complete interrupt occurs. At this time the card subroutine checks for errors. If an error occurred, PNCH0 exits to \$PST4; PNCH1 informs the user of the error and sets up to terminate or retry the operation.

Feed. Initiates a card feed cycle. This function advances all cards in the machine to the next station; that is, a card at the punch station advances to the stacker, a card at the read station advances to the punch station, and a card in the hopper advances to the read station. No data is punched as a result of a feed function and no column interrupts occur.

When the card advance is complete, an Operation Complete interrupt occurs. At this time the card subroutine checks for errors. If an error occurred, PNCH0 exits to \$PST4; PNCH1 informs the user of the error and sets up to terminate or retry the operation.

I/O Area Parameter

The I/O area parameter is the label of the control word that precedes the user's output area. The control word consists of a word count that specifies the number of columns of data to be punched, always starting with column 1. The word count must be in the range of 1-80.

Error Parameter

<u>PNCHO</u>. PNCHO has no error parameter. If an error is detected while an Operation Complete interrupt is being processed, the subroutine branches to \$PST4 with interrupt level 4 on, waiting for operator intervention. When the condition has been corrected, the 1442 made ready, and PROGRAM START pressed, the subroutine retries the operation.

<u>PNCH1</u>. PNCH1 has an error parameter. If an error is detected, the user can request the subroutine to terminate (that is, to clear the subroutine-busy indicator and turn off the interrupt level) or retry. Prior to a retry, the subroutine checks to see if the unit is ready. If the unit is not ready, the subroutine branches to \$PST4, with interrupt level 4 on, waiting for operator intervention.

DISK SUBROUTINES (C/PT SYSTEM)

The disk subroutines perform all reading and writing of data relative to disk storage. This includes the major functions: seek, réad, and write, in conjunction with readback check, file protection, and defective cylinder handling.

<u>DISKO</u>. The DISKO subroutine is the shortest version of the disk subroutine and can be used if not more than 320 words are to be read or written at one time.

DISK1. The DISK1 version is the standard subroutine for the disk and allows more than 320 words to be read or written; however, a full disk revolution might occur between sectors. DISK1 requires more core storage than DISK0.

<u>DISKN</u>. The DISKN subroutine minimizes extra disk revolutions in transferring more than 320 words. The DISKN subroutine requires more core storage than DISK1.

The major difference between DISK1 and DISKN is the ability of DISKN to read or write consecutive sectors on the disk without taking an extra revolution. If a full sector is written, the time in which the I/O command must be given varies. DISKN is programmed so that the extra revolution will not occur the majority of the time; DISK1 approximately 50 percent of the time.

All three disk subroutines have the same error-handling procedures.

Sector Numbering and File Protection

In the interest of providing disk features permitting versatile and orderly control of disk operations, programming conventions have been adopted concerning sector numbering, file protection, and defective cylinder handling. Successful use of the disk subroutines can be expected only if user programs are built within the framework of these conventions.

The primary concern behind these conventions is the safety of data recorded on the disk. To this end, the file-protection scheme plays a major role, but does so in a manner that is dependent upon the sector-numbering technique. The latter contributes to data safety by allowing the disk subroutine to verify the correct positioning of the access arm before it actually performs a write operation. This verification requires that sector identification be prerecorded on each sector and that subsequent writing to the disk be done in a manner that preserves the existing identification. The disk subroutines have been organized to comply with these requirements.

Sector Numbering. The details of the numbering scheme are as follows: each disk sector is assigned an address from the sequence 0,1,...,1623, corresponding to the sector position in the ascending sequence of cylinder and sector numbers from cylinder 0 sector 0 (outermost), through cylinder 202 sector 7 (innermost). The user can address cylinders 0 through 199. The remaining three cylinders are reserved for defective cylinder handling.

Each cylinder contains eight sectors and each sector contains 321 words. The sector address is recorded in the first word of each sector and occupies the rightmost eleven bit positions. Of these eleven positions, the three low-order positions identify the sector (0-7) within the cylinder. Utilization of this first word for identification purposes reduces the per sector availability of data words to 320; therefore, transmission of full sectors of data is performed in increments of 320 words. The sector addresses must be initially recorded on the disk by the user and are thereafter rewritten by the disk subroutines as each sector is written (see "Disk Initialization" in this section).

<u>File Protection</u>. File protection is provided to guard against the inadvertent destruction of previously recorded data. By having the write functions (except write immediate) uniformly test for the file-protect status of sectors that they are about to write, this control can be achieved.

This convention is implemented by assigning a file-protected area to each disk. The address of the first unprotected sector (0000-1623) on each disk is stored within the disk subroutine. Every sector below this one is file-protected, i.e., no writing is permitted below this address.

Defective Cylinder Handling

A defective sector is one in which, after ten retries, a successful writing operation cannot be completed. A cylinder having one or more defective sectors is defined as a defective cylinder. The disk subroutines can operate when as many as three cylinders are defective.

Since there are 203 cylinders on each disk, the subroutine can "overflow" the normally used 200 cylinders when defective cylinders are encountered (see "Effective Address Calculation" in this section).

The address of each defective cylinder is stored within the disk subroutines by the user (see "Disk Initialization" in this section).

If a cylinder becomes defective during an operation, the user can move the data in that cylinder and each higher-addressed cylinder into the next higher-addressed cylinder. Then the address of the new defective cylinder can be stored in DISKx +16, +17, or +18 and normal operation continued. Thus the user should not store the new defective cylinder address in DISKx and then continue normally because the effective sector address computation then yields a sector address eight higher than is desired (see "Effective Address Calculation" in this section).

If there are no defective cylinders, all three words in the defective cylinder table contain /0658. If, for example, only sector 0009 is defective, the table would contain /0008 (cylinder 1), /0658, and /0658.

Calling Sequence

Label	Operation	F	т		Operands & Remarks
21 r	:	17	33		35 40 45 50 55 y
	L.I.B.F				DISKA CALL DISK ILO
	O.C	_			/ 16. C. d. C. O.N.T.R.O.L. P.A.R.A.METTER
	p.c.				I.O.A.R., I.I.I.I.O. A.R.E.A. P.A.R.A.M.E.T.E.R.
	0.0	1	Ц		ERROR ERROR PARAMETER
	<u> </u>	+			
	1	+			
ERROR	O.C.			_	#1-# ADDRESS
	1	+			1. 1. <u>1. 1. 1. 1. 1. 1. 1. 1. 1. 1. 1. 1. 1. 1</u>
	- <u> </u>	+			
	1	+			
	BISIC	I	Ц		EIRIROR
	····	+-	Ц	_	<u></u>
	<u></u>	+-	Ц		· · · · · · · · · · · · · · · · · · ·
4 ب ب ب	- <u>*</u>	-	\square	_	
I.O.A.R.	<i>D.C.</i>		\square		FL
<u> </u>	P.C.		\square	_	SI ALL ALL SIECITIOR ADDRESS
	B.S.S.	-			here is the state of the state
	1		$\lfloor \rfloor$		

where

a is 0, 1, or N.

b is the I/O function digit,

c is in DISKN test function, the logical drive number. Otherwise c is 0.

d is the Seek option digit,

e is the Displacement option digit,

f is the number of words to be transferred to or from the disk,

g is the sector address at which the transfer is to begin,

h is the length of the I/O area. h must be equal to or greater than f.

The calling sequence parameters are described in the following paragraphs.

Control Parameter

This parameter consists of four hexadecimal digits as shown below:

	1	2	3	4
I/O Function		1	t	t
Not Used				
Seek Option				
Displacement Option.				

1/0 Function

The 1/O function digit specifies the operation to be performed on disk storage. The functions, their associated digital value, and the required parameters are listed and described below.

Function	Digita <u>Value</u>	l Required <u>Parameters</u> '		
Test	0	Control, I/O Area		
Read	1	Control, I/O Area, Error		
Write wit Out RBC	th- 2	Control, I/O Area, Error		
Write with RBC	3	Control, I/O Area, Error		
Write Immediate	e 4	Control, I/O Area		
Seek	5	Control, I/O Area, Error		
Any parameter not required for a particular function must be omitted.				

Test. Branches to LIBF+3 if the previous operation has not been completed, to LIBF+4 if the previous operation has been completed.

Note: This function requires the I/O area parameter even though it is not used.

<u>Read</u>. Positions the access arm and reads data into the user's I/O area until the specified number of words has been transmitted. Although sector-identification words are read and checked for agreement with expected values, they are neither transmitted to the I/O data area nor counted in the number of words transferred.

If, during the reading of a sector, a read check occurs, up to ten retries are attempted. If the error persists, the function is temporarily discontinued, an error code is placed in the Accumulator, the address of the faulty sector is placed in the Extension, and an exit is made to the error subroutine specified by the error parameter.

Upon return from the error subroutine, that sector operation is reinitiated or the function is terminated, depending on whether the Accumulator is nonzero or zero.

Write With Readback Check. This function first checks whether or not the specified sector address is in a file-protected area. If it is, the subroutine places the appropriate error code in the Accumulator and exits to location /0028.

If the specified sector address is not in a file-protected area, the subroutine positions the access arm and writes the contents of the indicated I/O data area into consecutive disk sectors. Writing begins at the designated sector and continues until the specified number of words has been transmitted. A readback check is performed on the data written.

If any errors are detected, the operation is retried up to ten times. If the function still cannot be accomplished, an appropriate error code is placed in the Accumulator, the address of the faulty sector is placed in the Extension, and an exit is made to the error subroutine designated in the error parameter.

Upon return from this error subroutine, the same sector operation is reinitiated or the function is terminated depending upon whether the contents of the Accumulator is nonzero or zero.

As each sector is written, the subroutine supplies the sectoridentification word. The identification word for the first sector is obtained from the I/O area, although it and subsequently generated identification words are not included in the word count. Writing less than 320 words on any sector sets the remaining words in that sector to zero.

<u>Write Without Readback Check</u>. This function is the same as the function described above except that no readback check is performed.

<u>Write Immediate</u>. Writes data with <u>no</u> attempt to position the access arm, check for file-protect status, or check for errors. Writing begins at the sector number specified by the rightmost three bits of the sector address. This function is provided to fulfill the need for more rapid writing to the disk than is provided in the previously described write functions. Primary application will be found in the "streaming" of data to the disk for temporary bulk storage.

As each sector is written, the subroutine supplies the sectoridentification word. The identification word for the first sector is obtained from the I/O area, although it and subsequently generated identification words are not included in the word count. Writing less than 320 words sets the remainder of the sector to zero. <u>Seek</u>. Initiates a seek as specified by the seek option digit. If any errors are detected, the operation is retried up to ten times.

Seek Option

If zero, a seek is executed to the cylinder whose sector address is in the disk I/O area control word; if nonzero, a seek is executed to the next cylinder toward the center of the disk, regardless of the sector address in the disk I/O area control word. This option is valid only when the seek function is specified.

The seek function requires that the user set up the normal I/O area parameter (see "I/O Area Parameter" in this section) even though only the sector address in the I/O area is used. The I/O area control (first) word is ignored.

Displacement Option

If zero, the sector address word contains the absolute sector identification; if nonzero, the file-protect address for the specified disk is added to bits 4-15 of the sector address word to generate the effective sector identification. The file-protect address is the sector identification of the first unprotected sector.

I/O Area Parameter

The I/O area parameter is the label of the first of two control worls which precede the user's I/O area.

The first word contains a count of the number of data words that are to be transmitted during the disk operation. If the DISK1 or DISKN subroutine is used, this count need not be limited by sector or cylinder size, since these subroutines cross sector and cylinder boundaries, if necessary, in order to process the specified number of words. However, if the DISKO subroutine is used, the count is limited to 320.

The second word contains the sector address where reading or writing is to begin. Bits 0-3 are used for device identification and must be zero. Bits 4-15 specify the sector address. Following the two control words is the user's data area. Note: The I/O area parameters are not available to the user until the requested operation is completed. The word count and sector addresses may be altered during a requested disk operation but are restored at the completion of the operation.

Error Parameter

Refer to the section "Basic ISS Calling Sequence".

Important Locations

The relative locations within the DISKO, DISK1, and DISKN subroutines are defined as follows:

DISKx	+0 -	entry point from calling
		transfer vector when LIBF
		DISKx is executed.
	+2 -	loader stores address of
	-	first location (in the
		calling transfer vector)
		accigned to DICKy
	×8 _	assigned to Diskx.
	T4 -	entry point from 115
		nandling Disk Storage
	_	interrupts.
	+7 -	area code for Disk Storage.
	+8 -	zero.
	+9 -	zero.
	+10 -	cylinder identification
		(bits 4-12) of the cylinder
		currently under the disk
		read/write heads (loaded as
		+202).
	+11 -	unused.
	+12 -	reserved.
	+13 -	sector address (bits 4-15)
		of the first
		non-file-protected sector
		for disk storage (loaded as
	+14 -	reserved.
	+15 -	reserved
	+16 -	sector address of the first
		defective culinder for dick
		storage deaded as +1620)
	.17	scorage (roaded as +1024).
	+17 -	sector address of the second
		derective cylinder for disk
	.10	storage (Loaded as +1624).
	+18 -	sector address or the third
		detective cylinder for disk
		storage (Loaded as +1624).

Effective Address Calculation

An effective disk address is calculated as follows:

- Start with the user-requested sector address (found in the sector address word of the I/O area).
- If the displacement option (found in the control parameter) is nonzero, add the sector address of the first non-file-protected sector (found in DISKx +13).

Note: This starting address will cause a preoperative error exit to location /0029 if over 1599.

- If the resulting address is equal to or greater than the sector address of the first defective cylinder (found in DISKx +16), add 8.
- 4. If the resulting address is equal to or greater than that of the second defective cylinder (found in DISKx +17), add 8 more.
- 5. If the resulting address is equal to or greater than that of the third Note: In no ca defective cylinder (found in DISKx +18), the DM2 System. add 8 more.

The address obtained from steps 1-5 is the effective sector address.

Disk Initialization

It is the user's responsibility to correctly load DISKx +13, +16, +17, and +18 at execution time and whenever a new disk is initialized. The following programs can be used to perform these functions.

Disk Pack Initialization Routine (DPIR). The functions of this program are to write sector addresses on a disk, to detect any defective cylinders, and to store defective cylinder information, file-protect addresses, and a disk label in sector 0 of the disk. The operating procedures for DPIR are located in the publication IBM 1130 Card/Paper Tape Programming System Operator's Guide.

Set Pack Initialization Routine (SPIR0, SPIR1, and SPIRN). The function of these subroutines is to store defective cylinder information and the file-protect address from sector 0 of the disk into the appropriate DISKx subroutine. If the above subroutines are not used, the starting address of the DISKx routine can be loaded into an index register for easy use in reaching the specified locations:

Lateri	Operation	11	Operandy & Remarks
	L.D.	11	4.I.B.F.
	S.L.A.	-+-+-+	BL
	S.R.T.	-+ +-+	BI
	S.T.X		LOAD to be a second a second descent of the second se
	Harry 1		LOAD+1 ADD IN. TV. ADDRESS.
	4	-+ + +	D. Q. Q. Q. Z. I ADD. CONSITANTS ITO REACH
	SILO		LOAD+1
LOAD		12	$\mathbf{\mathcal{A}}_{1} = \mathbf{\mathcal{A}}_{1} = \mathbf{\mathcal{A}}_{1} + \mathbf{\mathcal{A}}_{2} = \mathbf{\mathcal{D}}_{1} + \mathbf{\mathcal{A}}_{2} + \mathbf{\mathcal{A}}_{2} = \mathbf{\mathcal{D}}_{1} + \mathbf{\mathcal{A}}_{2} + $
	┝┤╩╾╼╶╺╌┥	+++	and ^a nd an
حد م د د .	┟┽╩┷╍╍╸┟	╶┿╍┊╞	and <mark>Baranan and a brain and a second s</mark>
<u></u>	- •··· + -		an <mark>te de la companya de la</mark>
D.ØØ.Ø2	p.c	4+.+	a the second seco
LIBE.	ASI.	- 3	$n_{1} \dots n_{1} \dots n_{1} \dots n_{1} \dots n_{1} \dots n_{1} \dots \dots n_{1} \dots \dots$
	╞┽╩╍╸┕╡	-+++	المراجع والمراجع والمراجع والمراجع فالمراجع والمراجع والمراجع والمراجع والمراجع والمراجع والمراجع والمراجع
	- · · · · · · · · · · · · · · · · · · ·	-+ +-+	an <mark>t ^aharan dar kanda dar dar dar dar dar dar der kander der kander der der der der der einen der der der der der</mark>
<u> </u>		+++	الفريد فارد فارد بالمالة الاستخلاط المتعاط فالمتعاط فالمتعاط فالمتعاط فالمتعاد والمالية المالية الم
Citille	D.C	- ┽╌┽╺╂	$[g_{1,1,1}, \dots, f_{1,n-1}, \dots, f_{n}, g_{n}, g_{n},$
- A & L -	┢╍╋╍┿╾╇┈╇╌╇	┽┼┽	<mark></mark>

The SPIR is a special-purpose utility subroutine. It is not called by LIBF as are the other disk subroutines described in this section. SPIR0 must be used if DISK0 is called, SPIR1 if DISK1 is called, or SPIRN if DISKN is called.

Note: In no case should SPIR be used with the DM2 System.

The SPIR reads sector 0000 from the disk and stores the first four words into the disk ISS that is in core. Therefore, the SPIR subroutine should be called before any calls are made to the disk ISS.

The calling sequence for SPIR is as follows:

CALL SPIRx DC /0000

The four words read from sector 0000 are described under "Disk Pack Initialization Routine" in the publication <u>IBM 1130</u> <u>Card/Paper Tape Programming System</u> <u>Operator's Guide</u>.

DISK SUBROUTINES (DM2 SYSTEM)

All disk subroutines used by the DM2 system (including DISKZ) reside in the IBM System area on the monitor disk. The disk subroutines are stored in a special core image format in this area rather than in the System Library, since the DM2 system always requires a disk I/O subroutine. The required version is fetched by the Core Image Loader just prior to execution. The disk subroutines used with the Monitor system are DISKZ, DISK1, and DISKN.

DISKZ. DISKZ is intended for use in a FORTRAN environment in which FORTRAN I/O is used. DISKZ makes no preoperative parameter checks and offers no file protection. It is the shortest of the three disk I/O subroutines and requires a special calling sequence (see "DISKZ-Disk I/O Subroutine"). This calling sequence can also be used with DISK1 and DISKN. DISKZ is also used by the RPG disk subroutines.

<u>DISK1</u>. DISK1 is intended for use by Assembler language programs in which the core storage requirement is of more importance than the execution time. DISK1 is longer than DISKZ but is the shorter of the two subroutines intended for use in Assembler language programs (DISK1 and DISKN). However, DISK1 does not minimize extra disk revolutions when transferring more than 320 words.

DISKN. DISKN minimizes extra disk revolutions in transferring more than 320 words. It provides all the functions DISK1 does and also operates as many as 5 drives simultaneously.

Two versions of DISKN are distributed with the Disk Monitor System. Both versions are called by the same calling sequence. The difference between them is the way they control disk drives. One version of DISKN, shown in the next drawing, can control as many as 5 singledisk drives simultaneously. This version of DISKN is for systems having only 2315 Disk Cartridges (mounted in 2310 Disk Storage Drives and/or the 1131 CPU).



The other version of DISKN, shown in the next drawing, can simultaneously control a single-disk drive in the 1131 CPU and two 2311 Disk Storage Drives (only one of the disks in each pack). This version of DISKN is for systems having 1316 Disk Storage Packs (mounted in 2311 Disk Storage Drives), and--optionally--a 2315 Disk Cartridge mounted in the 1131 CPU.



During loading of the Disk Monitor System, the 2310 version of DISKN is automatically placed into the IBM System Area on disk. If your system contains 2311s, however, you must replace this version with the 2311 version of DISKN before you load the Disk Monitor System card deck. (See "Monitor System Initial Load and System Reload" in IBM 1130 Disk Monitor System, Version 2, Programmer's and Operator's Guide.)

<u>Note</u>: Both DISK1 and DISKN can be specified on the Monitor XEQ record for use with FORTRAN programs. However, they offer no real advantage over DISKZ if they are called by the disk FORTRAN I/O subroutine.





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One of the major differences among the disk subroutines is the ability to read or write consecutive sectors on the disk without taking extra revolutions. If full sectors are written, the time in which the I/O command must be given varies. DISKN is programmed so that transfers of more than 320 words are made with a minimum number of extra revolutions occuring between sectors.

DISK1 and DISKN have the same errorhandling procedures.

Note: In the DM2 system, the disk I/O subroutines are not stored in the System Library; consequently they do not have LET entries.

Sector Numbering and File Protection

In the interest of providing disk features permitting versatile and orderly control of disk operations, programming conventions have been adopted concerning sector numbering, file protection, and defective sector handling. Successful use of disk I/O subroutines can be expected only if user programs are built within the framework of these conventions. The primary concern behind the conventions is the safety of data recorded on the disk. To this end, the file-protection scheme plays a major role, but does so in a manner that is dependent upon the sector-numbering technique. The latter contributes to data safety by allowing the disk I/O subroutine to verify the correct positioning of the access arm before it actually performs a write operation. This verification requires that sector identification be prerecorded on each sector and that subsequent writing on the disk be done in a manner that preserves the existing identification. The disk I/O subroutines support these requirements.

Sector Numbering. Each disk sector is assigned an address from the sequence 0, 1, ...,1623, corresponding to the sector position in the ascending sequence of cylinder and sector numbers from cylinder 0, sector 0 (outermost), through cylinder 202, sector 7 (innermost). The user can address cylinders 0 through 199. The remaining three cylinders are reserved for defective cylinder handling.

Each cylinder contains eight sectors and each sector contains 321 words, counting the sector address. The sector address is recorded in the first word of each sector and occupies the rightmost eleven bit positions. Of these eleven positions, the three low-order positions identify the sector (0-7) within the cylinder. Utilization of this first word for identification purposes reduces the per sector availability of data words to 320; therefore, transmission of full sectors of data is performed in increments of 320 words.

Sector addresses must be initially recorded on the disk by the user (via DISC or DCIP: see IBM 1130 Disk Monitor System, Version 2, Programmer's and Operator's <u>Guide</u>), and are thereafter rewritten by the disk I/O subroutines as each sector is written.

<u>Note</u>: Although not actually written on the disk, the logical drive code must be part of the sector address parameter (bits 1-3) which is stored in the second word of the I/O area. Bit 0 must always be zero.

File Protection. File protection is provided to prohibit the inadvertent destruction of previously recorded data. This control is achieved by having all write functions (except write immediate) test for the file-protection status of sectors they are about to write. Each cartridge has a file-protect address in COMMA. This address is the address of the first unprotected sector, i.e., the address of the beginning of Working Storage. Every sector, from sector 0 up to the sector address maintained in COMMA, is file-protected. The initial assignment of the file-protect address is performed by the disk initialization program DCIP or DISC: see IBM 1130 Disk Monitor System, Version 2, Programmer's and Operator's Guide. Subsequent updating of the file-protect address is performed by the Monitor programs.

Defective Sector Handling. A defective sector is a sector on which a read or write function cannot be successfully completed during initialization of the cartridge. A cylinder having one or more defective sectors is defined as a defective cylinder. The disk I/O subroutines can accommodate as many as three defective cylinders per cartridge. Since there are 203 cylinders on each disk, the disk I/O subroutines can "overflow" the 200 cylinders normally used when defective cylinders are encountered (see "Effective Address Calculation" in this section). g is the sector address, including the logical drive code, at which the transfer is to begin,

h is the length of the I/O area. h must be equal to or greater than f.

Control Parameter

This parameter consists of four hexadecimal digits, shown below:

	1	2	3	4		
I/O Function		Å	1	1		
Logical Drive Code (DISKN Test function	Only	n				
Seek Option						
Displacement Option_						

I/O Function

The I/O function digit specifies the operation to be performed on disk storage. The functions, their associated digital value, and the required parameters are listed and described below.

Function	Digita <u>Value</u>	l Required <u>Parameters</u> ⁴		
Test	0	Control, I/O Area		
Read	1	Control, I/O Area, Error		
Write wit out RBC	:h- 2	Control, I/O Area, Error		
Write with RBC	3	Control, I/O Area, Error		
Write Immediate	e 4	Control, I/O Area		
Seek	5	Control, I/O Area, Error		
[†] Any parameter not required for a particular function must be omitted.				

<u>Test</u>. Branches to LIBF+3 if the previous operation on the drive has not been completed, to LIBF+4 if the previous operation has been completed.

Note: This function requires the I/O area parameter even though it is not used.

Calling Sequence

Label	Operation	F	F	Γ	Operands & Remarks
2 25	2 30	32	n	1	25 40 45 50 55 60
	LIBF			[DITSKALL CALL DISK J/O
	D.C.				/1b.Ode. (CONTROL PARAMETER .
	D.C.		L		I.O.A.R. , , , , I.I.O. A.R.E.A. P.A.R.A.METLER
	D.C.		L	L	ERROR ERRIGER PARAMETER
	·	_		L	<u></u>
L				L	, , , , , , , , , , , , , , , , , , , ,
Lul	·		L	L	<u> </u>
ERROR	D.C.			L	KITK RETURN ADDRESS
[••••			L	<u></u>
	•		Ŀ	L	<u></u>
	. .			L	
	BSC	I	Ĺ	L	ERRIDIR REFUERING TO CALLER
	. <u>.</u>		L	L	
	<u> </u>		1	L	<u></u>
	<u></u>			L	<u>, , , , , , , , , , , , , , , , , , , </u>
I.O.A.R.	D.C.	4		L	FILLER WORD COUNTLELE
	D.C.			L	9 SIEICITOIR ADDRIESS
Lund	855	Ш	L	L	here is a start of the ABIE A start of the start
			L	Ļ	<u></u>

where

a is 1 or N. Note that LIBF DISKO is equivalent to LIBF DISK1.

b is the I/O function digit,

d is the Seek option digit,

e is the Displacement option digit,

f is the number of words to be transferred to or from the disk,

<u>Read</u>. Positions the access arm and reads data into the user's I/O area until the specified number of words has been transmitted. Although sector identification words are read and checked for agreement with expected values, they are neither transmitted to the I/O area nor counted in the number of words transferred.

If, during the reading of a sector, a read check occurs, up to 16 retries are attempted. If the error persists, the function is temporarily discontinued, an error code is placed in the Accumulator, the address of the faulty sector is placed in the Extension, and an exit is made to the error subroutine specified by the error parameter.

Upon return from the error subroutine, the operation is either reinitiated or terminated, depending on whether the Accumulator is nonzero or zero, respectively.

<u>Write With Readback Check</u>. Checks whether or not the specified sector address is in a file-protected area. If it is, the subroutine places the appropriate error code in the Accumulator and exits to \$PRET.

If the specified sector address is not in a file-protected area, the subroutine positions the access arm and writes the contents of the indicated I/O area onto the disk. Writing begins at the designated sector and continues until the specified number of words have been transmitted. A readback check is performed on the data written.

Writing less than 320 words on any sector sets the remaining words in that sector to zero.

If any errors are detected, the operation is retried up to 16 times. If the function cannot be accomplished, an appropriate error code is placed in the Accumulator, the address of the faulty sector is placed in the Extension, and an exit is made to the error subroutine designated by the error parameter.

Upon return from this error subroutine, the operation is either reinitiated or terminated, depending upon whether the Accumulator is nonzero or zero, respectively.

As each sector is written, the subroutine supplies the sector-identification word. The identification word for the first sector is obtained from the I/O area, although it and subsequently generated identification words are not included in the word count. Write Without Readback Check. Functions the same as Write With Readback Check except that no readback check is performed.

Write Immediate. Writes data with no attempt to position the access arm, check for file-protect status, or check for errors. Writing begins at the sector number specified in the user's I/O area. This function provides more rapid writing to the disk than is provided in the previously described Write functions; it provides, for example, the ability to "stream" data to the disk for temporary bulk storage or to write addresses in Working Storage (see "System Library Mainline Programs (DM2 System) ADRWS").

Writing less than 320 words on any sector sets the remaining words in that sector to zero.

As each sector is written, the subroutine supplies the sector-identification word. The identification word for the first sector is obtained from the I/O area, although it and subsequently generated identification words are not included in the word count.

Seek. Initiates a seek as specified by the seek option digit. If any errors are detected, the operation is retried up to 16 times.

The seek function requires that the user set up the normal I/O area parameters (see "I/O Area Parameter" in this section) even though only the sector address in the I/O area is used.

Logical Drive Code. Digit 2 defines the logical drive code (0, 1, 2, 3, or 4). This digit is used only with the DISKN test function.

Seek Option. If digit 3 of the control parameter is zero, a seek is executed to the cylinder whose sector address is in the I/O area; if nonzero, a seek is executed to the next nondefective cylinder toward the center, regardless of the sector address in the I/O area. This seek to the next nondefective cylinder must be taken into consideration when planning for the "streaming" of data. This option is valid only when the seek function is specified.

Displacement Option. If digit 4 of the control parameter is zero, the sector address word contains the absolute sector identification; if nonzero, the file-protect address for the specified cartridge is added to bits 4-15 of the sector address word to generate the effective sector identification. The file-protect address is the sector identification of the first unprotected sector, i.e., the address of the first sector of Working Storage.

1/0 Area Parameter

The I/O area parameter is the label of the first of two control words which precede the user's I/O area. The first word contains the number of data words that are to be transferred during the disk operation. This number need not be limited by sector or cylinder size, since the subroutines cross sector and cylinder boundaries, if necessary, in order to transmit the specified number of words.

The second word contains the sector address at which reading or writing is to begin. Bit 0 must be zero. Bits 1-3 are the device identification (logical drive code) and must be 0, 1, 2, 3, or 4. Bits 4-15 specify the sector address. The user's I/O area follows the two control words.

Note: The I/O area parameters are not available to the user until the requested operation is completed. The word count and sector addresses may be altered during a requested disk operation but are restored at the completion of the operation.

Error Parameter

If an error is detected, the user can request the subroutine to terminate (that is, to clear the subroutine's busy indicator and turn off interrupt level 2) or to branch to \$PST2, with interrupt level 2 on, waiting for operator intervention.

Effective Address Calculation

An effective disk address is calculated as follows:

- Obtain the sector address found in the sector address word of the I/O area.
- 2. If the displacement option digit in the control parameter is nonzero, add the sector address of the first sector that is not file-protected.

Note: This address causes an exit to \$PRET if it exceeds 1599.

- 3. If the resultant address is equal to or greater than the sector address of the first defective cylinder, add 8.
- 4. If the resultant address is equal to or greater than that of the second defective cylinder, add 8 more.

If the resultant address is equal to or greater than that of the third defective cylinder, add 8 more.

The address obtained from <u>steps 1-5</u> is the effective sector address. Defective cylinders are handled in this manner for all operations, including seek and write immediate.

Monitor Entry Point

Both DISK1 and DISKN can be entered by a BSI L /00F2, the monitor entry point (see calling sequence of DISK2). This entry point is used by the system programs and by FORTRAN programs when DISK1 or DISKN is specified in the XEQ record.

Reading begins at the designated sector where the access arm reads data into the user's I/O area until the specified number of words has been transmitted.

Writing begins at the designated sector and continues until the specified number of words have been transmitted. A readback check is performed on the data written on the disk. When DISK1 and DISKN are entered via /00F2, however, there is no check for writing in the file-protect area.

There is no possibility of performing a seek operation when using the monitor entry point. A word count of zero will result in a preoperative error wait. All postoperative errors will cause a branch to \$PST2 (see Appendix B).

<u>Disk Initialization</u>

Before the DM2 system is stored on a cartridge, the Disk Cartridge Initialization Program (DCIP) must be executed. This program writes sector addresses on the disk cartridge, detects any defective cylinders, stores defective cylinder information and a cartridge ID in sector 0 of cylinder 0, and initializes DCOM. The operating procedure for DCIP is listed in the publication IBM <u>1130</u> Disk Monitor System, Version 2, Programmer's and Operator's Guide.

DISKZ - DISK I/O SUBROUTINE

The DISKZ subroutine offers no file protection, no preoperative parameter checks, no write immediate function, and no
write without readback check function. It is intended for use by the DM2 programs, RPG programs, and FORTRAN programs in which disk FORTRAN I/O is used. Although DISKZ has many of the characteristics of an ISS, it is assembled as though it were a mainline and is stored in a special Core Image format in the System Device Subroutine area.

Calling Sequence

Lovel	Operation	F T	Operands & Remorks
		+++	
	L.D.D.	+++	LEUSIL
╺┻╼┻╴╍╼┾╺╄╴	-18-8-1	- ┺┼┼	
	╇╧╍┶┡	+++	┉ <mark>╬╷╫_{╺╋┙╊┙}┺╌╽╴╫╶╎╶╎╴╎╴╢╴╢╺╅┉╅┉╋┉╇╖╋┙╋╺╋┙╋╺╋┥╋╺╋┥╋╸╋╺╋┥╋╺╋┥╋╺╋╸╋╸╋╸╋</mark>
<u> </u>	∔⊷⊷∔	-4-4-4	
	alian ali	-+	<u>. </u>
	8.5.5	E	M
LUS.L.	D.C.		1,000
	הם ה	L	LOAR I I I A AREA PARAMETER
	•		
	1	111	
	466	6	a
7000	100		
LUNCO.D.	40.00	++	
- A	D.C.	+++	
┹┹┻┻┣	BS.S.		<u> du</u>
	1	++-	
	· · · · ·		<u></u>
ليبيب	1		
DZØØØ	E.Q.U.	11	/.ø.ø.F. 2
- 1-11		++1	
	╺╋╼┺╼╄╌┵┊	+++	
		+++	<mark>┍╴╪╶╉┉╪╶┿╍╪╌┿┉╋┉╪╴┶╍┽╴┶╼┫┉╋╍┶╶┽┄┿╼╣╍╅┍╅┊╿╶┟╍╢┉╅╸┿╺╢┉╢┉╅╴</mark> ┿╍ <mark>╢┉╅╴</mark> ┼┍╹╵┥╴╿╴╹

where

a is the I/O function digit: 0 indicates a read, 1 a write.

b is the number of words to be transferred to or from the disk,

c is the sector address at which the transfer is to begin,

d is the length of the I/O area. d must be equal to or greater than b.

The word count (first word of the buffer) must be nonnegative and must be on an even core boundary. The sector address must be the second word of the buffer. The logical drive code (0, 1, 2, 3, or 4), as defined by the // JOB DM2 control record, is in bits 1-3 of the sector address. Bit zero is always zero.

A word count of zero indicates a seek to the cylinder denoted in the sector address. File protection is not provided. If the access arm is not positioned at the cylinder addressed, DISKZ seeks to that cylinder before performing the requested function. A read follows each seek to verify that the seek was successful. No buffer is required for this read.

Buffer Size. Unlimited

<u>Operation</u>. DISKZ performs read, seek, and write with readback check functions. Each function returns control to the user after it has been initiated. To determine the completion of a disk operation, the user may test \$DBSY (location /00 EE in COMMA) until it is cleared to zero. DISKZ itself tests this word before initiating an operation. Following a write, this subroutine performs a readback check on the data just written. If it detects an error, it reexecutes the write. Similarly, if a sector is not located or if an error is detected during a read, DISKZ repeats the operation. All operations are attempted 16 times before DISKZ indicates an unrecoverable error.

If a partial sector (less than 320 words) is written, the remaining words of the sector are set to zero.

<u>Subroutines Required</u>. No other subroutines are required by DISKZ.

<u>Note</u>: It is important to realize that the DISKZ subroutine is designed to operate in an error-free environment; it is not recommended for general usage. The user should therefore use DISK1 or DISKN whenever possible.

1132 PRINTER SUBROUTINE (PRNT1)

The printer subroutine PRNT1 handles all print and carriage control functions relative to the IBM 1132 Printer (see also "1132 Printer/Synchronous Communications Adapter Subroutine (PRNT2)"). Only one line of data can be printed, or one carriage operation executed, with each call to the printer subroutine. The data in the output area must be in EBCDIC form, packed two characters per computer word. Any code other than those defined for the 1132 will be interpreted by the PRNT1 subroutine as a blank. (See "Appendix D. Character Code Chart".)

Calling Sequence

t abe I	Operation		F	τİ	Operands & Remarks
: r	2 30		32	11	35 40 45 50 35 er
	LIBF	L	4		P.R.N.T.I., CALL PRIMTER OUTPUT
	0.0			1	LIBICID . C.D.N.T.R.O.L. PARAMETER
	0.0			\downarrow	I.O.A.R. I.I.I.I.A. A.R.E.A. PARAMETTER
	0.0			1	ERRORERROR PARAMETER
	·			+	
-+	·	Ц			
- _	·				<u></u>
ERROR	<i>b</i> ,c, ,				RETURN ADDRESS
	••••			\perp	
	•• • •			\perp	
	·				<u></u>
	B.S.C.		I	1	ERROR , RETURN TO CALLER
	·				
	••••				
Liul	1. L.				
I.O.A.R.	D.C.			L	FILL WORD COUNTLE
	B.S.S.			L	h

where

b is the I/O function digit,

c is the "immediate" carriage operation digit,

d is the "after-print" carriage operation digit,

f is number of words to be printed on the 1132 Printer,

h is the length of the I/O area. h must be equal to or greater than f.

The calling sequence parameters are described in the following paragraphs.

Control Parameter

This parameter consists of four hexadecimal digits which are used as shown below.



I/O Function

The I/O function digit specifies the operation to be performed on an 1132 Printer. The functions, their associated digital values, and the required parameters are listed and described below.

Function	Digita <u>Value</u>	l Required <u>Parameters</u> ⁴							
Test	0	Control							
Print	2	Control, I/O Area, Error							
Control Carriage	3	Control							
Print Numeric	4	Control, I/O Area, Error							
¹ Any parameter not required for a particular function must be omitted.									

<u>Test</u>. Branches to LIBF+2 if the previous operation has not been completed, to LIBF+3 if the previous operation has been completed.

<u>Print</u>. Prints characters from the user's I/O area, checking for channel 9 and 12 indications. If either of these conditions is detected, the subroutine branches to the user's error subroutine after the line of data has been printed (see Appendix B for error codes). Upon return from this error subroutine, a skip to channel 1 is initiated or the function is terminated, depending upon whether the Accumulator is nonzero or zero.

<u>Control Carriage</u>. Controls the carriage as specified by the carriage control digits listed in Figure 7.

<u>Print Numeric</u>. Prints only numerals and special characters from the user's I/O area and checks for channel 9 and channel 12 indications. See "Print" above.

<u>Carriage Control</u>. Digits 2 and 3 specify the carriage control functions listed in Figure 7. An "immediate" request is executed before the next print operation; an "after-print" request is executed after the next print operation and replaces the normal space operation.

If the I/O function is print, only digit 3 is examined; if the I/O function is control, and digits 2 and 3 both specify carriage operations, only digit 2 is used.

If channel 9 or channel 12 is encountered during a carriage control function, a branch is made to the user's error subroutine at completion of the next print function.

Note: An after-print request will be lost if it is followed by an immediate request or by a print with spacing suppressed. If a series of after-print requests is given, only the last one will be executed. A skip operation must not be less than four lines.

Digit #2: Immediate Carriage Operations
Print Functions Not Used
Control Function 1 - Immediate Skip To Channel 1 2 - Immediate Skip To Channel 2 3 - Immediate Skip To Channel 3 4 - Immediate Skip To Channel 4 5 - Immediate Skip To Channel 5 6 - Immediate Skip To Channel 6 9 - Immediate Skip To Channel 9 C - Immediate Skip To Channel 12 D - Immediate Space Of 1 E - Immediate Space Of 2 F - Immediate Space Of 3
Digit #3: After-Print Carriage Operations
Print Functions 0 - Space One Line After Printing 1 - Suppress Space After Printing
Control Function 1 - Skip After Print To Channel 1 2 - Skip After Print To Channel 2 3 - Skip After Print To Channel 3 4 - Skip After Print To Channel 4 5 - Skip After Print To Channel 5 6 - Skip After Print To Channel 6 9 - Skip After Print To Channel 9 C - Skip After Print To Channel 12 D - Space 1 After Print E - Space 2 After Print F - Space 3 After Print

Figure 7. Carriage Control Operations for 1132 Printer

I/O Area Parameter

The I/O area parameter is the label of the control word that precedes the user's I/O area. The control word consists of a word count that specifies the number of computer words of data to be printed. The data must be in EBCDIC format, packed two characters per computer word. The word count must be in the range of 1-60. (See "Descriptions of Data Codes".)

Error Parameter

See "Basic ISS Calling Sequence".

1132 PRINTER/SYNCHRONOUS COMMUNICATIONS ADA PTER SUBROUTINE (PRNT2)

The printer subroutine PRNT2 is an additional printer subroutine for the IBM 1132 Printer, specifically provided to permit concurrent operation of the 1132 and the Synchronous Communications Adapter. PRNT2 handles all print and carriage control functions related to the 1132.

Only one line of data can be printed, or one carriage operation executed, with each call to the printer subroutine. The data in the output area must be in EBCDIC form, packed two characters per word. Any code other than those defined for the 1132 will be interpreted by the PRNT2 subroutine as a blank.

Restriction. The PRNT1 and PRNT2 subroutines are mutually exclusive; i.e., both subroutines can not be in core at the same time. Thus, if the Synchronous Communications Adapter is in operation, the PRNT2 subroutine must be used for concurrent operation of the 1132 Printer. If the PRNT2 subroutine is required in a core load for the concurrent operation of the 1132 Printer and the Adapter, all IBMand user-written programs in that core load using the PRNT2 subroutine must be modified to use the PRNT2 subroutine.

Calling Sequence

Label	Operation		1		Operandi & Remarks
n n	<i>v</i> x		1.		35 aC 45 ~ 55 ar aS
	LIBF				P.R.N.T.2. ,
	D.C.			1	1. D.C. d.d. ,
	D.C.		ľ		I.O.BR
	D.C.		T	Γ	E.R.R.O.R E.R.R.O.R. P.D.R.A.ME.T.E.R.
	• • • •	Ц			
	J	Ц	1	1	<u> </u>
	<u></u>	L	1	L	<u> </u>
ERROR	J.C.	Ц			<u>当一,老,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,</u>
	<u></u>	Ш			
<u> </u>	· · · · ·	11	+	1	
	•	Ll	1	1	<u></u>
	8.6.C.		Z	L	EIR.R.O.R. L.
	<u></u>	Ц	╇		<u></u>
	<u></u>	Ц	+	1	<u></u>
	der er	Ц			<u></u>
I.O.AR.	D.C.				full where the second with the second s
	8.5.5.	Ц			h
	<u> </u>	Ш	4	+-	<u></u>
1		11		1	

where

b is the I/O function digit,

c is the "immediate" carriage operation digit,

d is the "after-print" carriage operation digit,

f is the number of words to be printed on the 1132 Printer,

h is the length of the I/O area. h must be equal to or greater than f.

The calling sequence parameters are described in the following paragraphs.

Control Parameter

The control parameter consists of four hexadecimal digits which are used as shown below:



1/0 Function

The I/O function digit specifies the operation to be performed on the 1132 Printer. The functions, their associated digital values, and the required parameters are listed and described below.

Function	Digita Value	l Required <u>Parameters</u> ¹							
Test	0	Control							
Print	2	Control, I/O Area, Error							
Control Carriage	3	Control							
Print Numeric	4	Control, 1/0 Area, Error							
Any parameter not required for a particular function must be omitted.									

<u>Test</u>. Branches to LIBF+2 if the previous operation has not been completed, to LIBF+3 if the previous operation has been completed.

<u>Print</u>. Prints characters from the user's I/O area; checks for channel 9 and 12 indications. If either of these conditions is detected, the subroutine branches to the user's error routine after the line of data has been printed (see Appendix B for error codes). Upon return from this error routine, a skip to channel 1 is initiated or the operation is terminated, depending upon whether the Accumulator is nonzero or zero.

<u>Control Carriage</u>. Controls the carriage as specified by the carriage control digits listed in Figure 7.

<u>Print_Numeric</u>. Prints only numerals and special characters from the user's I/O area and checks for channel 9 and 12 indications. (See "Print" above.)

<u>Carriage Control</u>. Digits 2 and 3 specify the carriage control operations listed in Figure 7. An "immediate" request is executed before the next print operation; an "after-print" request is executed after the next print operation and replaces the normal space operation.

If the I/O function is Print, only digit 3 is examined; if the I/O function is Control Carriage, and digits 2 and 3 both specify carriage operations, only digit 2 is used.

Carriage control functions do not check for channel 9 and channel 12 indications.

1/0 Area Parameter

The I/O area parameter is the label of the control word that precedes the user's I/O area. The control word consists of a word count that specifies the number of words of data to be printed. The data must be in EBCDIC format, packed two characters per word. The word count must be in the range of 1-60.

Error Parameter

See "Basic ISS Calling Sequence".

1403 PRINTER SUBROUTINE (PRNT3)

The printer subroutine PRNT3, available only with the DM2 system, handles all print and carriage control functions relative to the 1403 Printer. Only one line of data can be printed and/or one carriage operation executed with each call to the printer subroutine.

The data in the output area must be in the 1403 character code, as defined in "Descriptions of Data Codes", and packed two characters per word. Each data code consists of seven bits and the total number of bits should always be a valid number. The first bit is the parity bit. If the remaining six bits correspond to a valid 1403 code, that character will be printed. A branch to the user error routine will or will not be made depending upon the validity of the parity bit. The user can specify a retry of the operation, if desired.

Calling Sequence

Lobel	Gueration	1	Operands & Remarks							
ىغ بىغان بەر بەر بەر بىلەردىن	LIBF		P.R.N.T.J. , C.A.L.L. P.R.I.N.T.E.R. OU.T.P.U.T							
			/ b.c.d.O CONTROL. PARAMETER .							
	D.C		I.O.A.R I.I.O. A.R.E.A. P.A.R.AME.T.E.R.							
			ERROR ERROR PARAMETER							
A . A A A	· · · ·		<u></u>							
	11									
	Π•									
ERROR	D.C.	r	*1-1#							
	TT•									
			*							
	1.									
	BSC	T	EPPOR							
	-									
	+++	+++								
·	+++	+++	<u> </u>							
7000		╶╋┊╋								
LUAK.	acc.	┽┽┽								
andradi akashiri	- B.S.S.	+++	the same share a start in the start start is a second start in the same							
	╅╁┶┷╼┪	\rightarrow	<u> </u>							

where

b is the I/O function digit,

c is the "immediate" carriage operation digit,

d is the "after-print" carriage operation digit,

f is the number of words to be printed on the 1403 Printer,

h is the length of the I/O area. h must be equal to or greater than f.

Control Parameter

This parameter consists of four hexadecimal digits which are used as shown below.



I/O Function

The I/O function digit specifies the operation to be performed on the 1403 Printer. The functions, their associated

digital values, and the required parameters are listed and described below.

Function	Digita <u>Value</u>	1 Required <u>Parameters</u> ¹							
Test	0	Control							
Print	2	Control, I/O Area, Error							
Control Carriage	3	Control							
¹ Any parameter not required for a particular function must be omitted.									

<u>Test</u>. Branches to LIBF+2 if the previous operation has not been completed, to LIBF+3 if the previous operation has been completed.

<u>Print</u>. Prints characters from the user's I/O area, checking for channel 9 and 12 and error indications. If any of these conditions are detected, the subroutine branches to the user's error subroutine after the line of data has been printed with an error code in the Accumulator (see Appendix B). Upon return from this error subroutine, a skip to channel 1 is initiated and the function is reinitiated or terminated, depending upon the error code and whether the Accumulator is nonzero or zero.

<u>Control Carriage</u>. Controls the carriage as specified by the carriage control listed in Figure 8.

<u>Carriage Control</u>. Digits 2 and 3 specify the carriage control functions listed in Figure 8. An "immediate" request is executed before the next print operation; an "after-print" request is executed after the next print operation and replaces the normal space operation.

If the function is print, only digit 3 is examined; if the function is control, and digits 2 and 3 both specify carriage operations, only digit 2 is used.

Carriage control functions do not check for channel 9 or channel 12 indications.

Note: An "after-print" request is lost if it is followed by an "immediate" request. If a series of "after-print" requests is given, only the last one is executed.

Digit #2: Immediate Carriage Operations
Print Functions Not Used
Control Function 1 - Immediate Skip To Channel 1 2 - Immediate Skip To Channel 2 3 - Immediate Skip To Channel 3 4 - Immediate Skip To Channel 4 5 - Immediate Skip To Channel 5 6 - Immediate Skip To Channel 6 7 - Immediate Skip To Channel 7 8 - Immediate Skip To Channel 8 9 - Immediate Skip To Channel 9 A - Immediate Skip To Channel 10 B - Immediate Skip To Channel 11 C - Immediate Skip To Channel 12 D - Immediate Space Of 1 E - Immediate Space Of 2 F - Immediate Space Of 3
Digit #3: After-Print Carriage Operations
Print Functions 0 - Space One Line After Printing 1 - Suppress Spaces After Printing
Control Function 1 - Skip After Print To Channel 1 2 - Skip After Print To Channel 2 3 - Skip After Print To Channel 3 4 - Skip After Print To Channel 4 5 - Skip After Print To Channel 5 6 - Skip After Print To Channel 6 7 - Skip After Print To Channel 7 8 - Skip After Print To Channel 8 9 - Skip After Print To Channel 9 A - Skip After Print To Channel 10 B - Skip After Print To Channel 11 C - Skip After Print To Channel 12 D - Space 1 After Print F - Space 3 After Print
Figure 8. Carriage Control Operations for

<u>Note</u>: A skip operation must not be less than two lines.

1/O Area Parameter

The I/O area parameter is the label of the control word that precedes the user's I/O area. The control word consists of a word count that specifies the number of words of data to be printed. The data must be in 1403 Printer code, packed two characters per word. The word count must be in the range of 1-60.

Error Parameter

See "Basic ISS Calling Sequence".

KEYBOARD/CONSOLE PRINTER

There are two ISSs for the transfer of data to and from the Console Printer and the Keyboard.

TYPE0

The TYPEO Subroutine handles input and output.

WRTY0

The WRTYO Subroutine handles output only. If a program does not require keyboard input, it is advantageous to use the WRTYO subroutine because it occupies less core storage than the TYPEO subroutine.

Only the TYPEO subroutine is described below; the WRTYO subroutine is identical, except that it does not allow the read-print function.

Calling Sequence

Label	Operation	F	T							_		_ 0	perandi	1 6. Rem	arks			
e	2 1	п	,,,		в		44			45					**		•	<u>،</u>
	LIBF		·		T'IY.F	P.E.	ó	_		C.		Lu	2.R.1	(ALT	ER	J	1.0	2
	0.0				1.6.0	1.0	6.			،Cı	O.N.	TRO	<u></u>	PA	R.A	ME	ιT.E	R
	D.C.				LOA	R				I.	10	نيم.	BE/	l, P	AIR	AA	AE.I	ER
								<u> </u>					1		_ د _			ب ب ا
	· · · · ·						لسب	L.,				. .			<u>.</u>	<u> </u>		
												. ب						
I.O.A.R.	DO.C.			4	fi i		i.		<u> </u>	Ma	O.R.	0.0	.	JAAT	<u>.</u>			
	BSS.				h		. د			I.	10	Al	RE	9	. 1			
											11		1.1				.	i

where

b is the I/O function digit,

f is the number of characters to be printed on the console printer for read-print operations and is 1/2 the number of characters to be printed on a print operation.

h is the length of the I/O area. h must be equal to or greater than f.

Control Parameter

This parameter consists of four hexadecimal digits, as shown below:



I/O Function

The I/O function digit specifies the operation to be performed on the Keyboard and/or Console Printer. The function, their associated digital values, and the required parameters are listed and described below.

Function	Digital <u>Value</u>	Required Parameters								
Test Read-Prin Print	0 1 1 2	Control Control, I/O Area Control, I/O Area								
Any parameter not required for a particular function must be omitted.										

<u>Test.</u> Branches to LIBF+2 if the previous operation has not been completed, to LIBF+3 if the previous operation has been completed.

<u>Read-Print</u>. Reads from the Keyboard and prints the requested number of characters on the Console Printer. The operation sequence is as follows:

- 1. The calling sequence is analyzed by the call portion of the subroutine, which then unlocks the keyboard.
- 2. When a key is pressed, a character interrupt signals the interrupt response portion that a character is ready to be read into core storage.
- 3. The interrupt response portion converts the keyboard data to Console Printer Code (see "Descriptions of Data Codes"). Each character is printed as it is read; the Keyboard is then unlocked for entry of the next character.

4. Printer interrupts occur whenever the Console Printer has completed a print operation. When the interrupt is received, the subroutine checks to determine if the final character has been read and printed. If so, the operation is considered complete. In the C/PT system, if the Console Printer becomes not-ready during printing, the subroutines loop, waiting for the Console Printer to become ready. In the DM2 system they trap to \$PRET or \$PST4 (see "Descriptions of Data Codes").

5. Steps 2 through 4 are repeated until the specified number of characters have been read and printed. The characters read into the I/O area are identical to IBM Card Code; that is, each 12-bit image is left-justified in one 16-bit word.

<u>Print</u>. Prints the specified number of characters on the Console Printer. A printer interrupt occurs when the Console Printer has completed a print operation. When an interrupt is received, the character count is checked. If the specified number of characters has not been written, printing is initiated for the next character. This sequence continues until the specified number of characters has been printed. Data to be printed must be in Console Printer code (see "Descriptions of Data Codes"), packed two characters can be embedded in the message where desired.

In read-print and print operations, printing begins where the printing element is positioned; that is, carrier return to a new line is not automatic when the subroutine is called.

Keyboard Functions

Keyboard functions provide for control by the TYPE0 subroutine and by the operator.

TYPE0 Subroutine Control

Three keyboard functions are recognized by the TYPEO subroutine.

Backspace. The operator presses the backspace key whenever the previous character is in error. The interrupt response portion senses the control character, backspaces the Console Printer, and prints a slash (/) through the character in error. In addition, the subroutine prepares to replace the incorrect character in the I/O area with the next character.

If the backspace key is pressed twice, the character address is decremented by +2, but only the last graphic character is slashed. For example, if ABCDE was entered and the backspace key pressed three times, the next graphic character to be entered replaces the C but only the E is slashed. If XYZ is the new entry, the printout shows ABCDEXYZ, but the buffer contains ABXYZ.

<u>Erase Field</u>. When the interrupt response portion recognizes the erase field control character, it assumes that the entire message is in error and is to be entered again. The subroutine prints two slashes on the Console Printer, restores the carrier to a new line, and prepares to replace the old message in the I/O area with a new message.

The old message in the I/O area is not cleared. Instead, the new message overlays the old, character by character. If the old message is longer than the new, the remainder of the old message follows the NL (new-line) character terminating the new message.

End of Message. When the interruptresponse portion recognizes the end-of-message (EOF) control character, it assumes the message has been completed, stores an NL character in the I/O area, and terminates the operation.

Operator Request Function (C/PT System)

By pressing the interrupt request key (INT REQ) on the Keyboard, the operator can inform the program that he wishes to enter data from the Keyboard or the Console Entry switches. The interrupt that results causes the TYPEO or WRTYO subroutine to execute an indirect BSI instruction to core location /002C, where the user must have previously stored the address of an interrupt request subroutine. Bit 1 of the Accumulator contains the Keyboard/Console Printer identification bit, that is, the device status word, shifted left two bits.

The user's interrupt request subroutine must return to the ISS subroutine via the return link. The user's subroutine is executed as a part of the interrupt handling. The interrupt level remains on until control is returned to the ISS subroutine (see "General Error-Handling Procedures, Postoperative Checks").

Operator Request Function (DM 2)

By pressing the Interrupt Request key (INT REQ) on the keyboard, the operator can inform the program that he wishes to enter data from the keyboard or the Console Entry switches. The interrupt that results causes the ILSO4 or ILSX4 subroutine to execute a BSI I \$IREQ instruction. \$IREQ is initialized with the address \$I420 in Resident Monitor. This allows the operator to terminate the job by pressing INT REQ key. If the user wants control, \$IREQ must be set to the user Interrupt Service subroutine. This subroutine can set indicators or read the Console Entry switches. If keyboard input/output is desired, only one call to ISS can be made. The user-written subroutine must return to exit address plus one, in ILSO4 or ILSX4. This is to turn off the interrupt and return to the program that was interrupted. In no case should the user perform an XIO sense Keyboard/Console with reset.

I/O Area Parameter

The I/O area parameter is the label of the control word that precedes the user's I/O area. The control word consists of a word count that specifies the number of words to be read or printed. This word count is equal to the number of characters if the read-print function is requested and is equal to 1/2 the number of characters if the print function is requested.

PAPER TAPE SUBROUTINES (C/PT SYSTEM)

The paper tape subroutines, PAPT1 and PAPTN, handle the transfer of data from the IBM 1134 Paper Tape Reader to core storage and from core storage to the IBM 1055 Paper Tape Punch. Any even number of characters can be transferred via one calling sequence.

The PAPTN subroutine must be used if simultaneous reading and punching are desired.

The PAPT1 subroutine can operate both devices, but only one at a time.

When called, the paper tape subroutine starts the reader or punch and then, as interrupts occur, transfers data to or from the user's I/O area. Input data is packed two characters per computer word by the subroutine; output data must already be in the packed format when the subroutine is called for a punch function.

Calling Sequence

	· · · ·	-	_		· ·					
Labet	Operation	- Þ	1				Op	erands & Rema	orks	
21 N	17 18		7 13	35	40	45	50		55	**
	L.I.B.F			P.A.	P.T.a.	L. CA	LL P	APER	TAP	E_I/0
	D.C.			1.0	c.Q.c.		NIT RO	L PA	RAME	TER
	DC .			LIO	A.R.	.. /	Q. AR	EA.P	A.R.A.M	ET.E.R.
	0.0			ER	ROR	ER.	ROR	PARA	METE	Runn
	·	4				المسعم		للمسالم	لللمناسب	1111
	<u>.</u>	1	1	ЦĿ				سينت	للتناب	
	·			ļļ.		í		بالمسالية	سيبل	الم الم الم الم
ERROR	P.C.	Ц.	.l	* 1-	8		TURN	_,A,D,D	RESS	
	<u></u>	1				حليتيت	يتناب	ىبىت	بنتنا	
· · · · · .	احد من ا	-		┝╎╍		للمستعب	ىلىئىت	بنبت	سينا	
	<u> • </u>	4	+	\square					لللله	
	8.S.C.	1	<u>د</u>	ER	RIQIR	RE	TURN	т.o	CALL	ER
	<u>to a d</u>			╞╞╼╺					سينات	
	· · · · ·	4	+	ļļ				للمستحيا	ليستعد السا	
	•••••	Ц	+			<u> </u>				. .
LOAR	p.c. ,	Ц	+	<u> </u> €			R.D. C	OUNT	i t i a a	
	BISIS	Ц	+-	h		 /	O_AR	EA		
				<u> </u>			<u> </u>		تتناب	المراجع المراجع

where

- a is a 1 or N,
- b is the I/O function digit,
- c is a check digit,
- e is a device-identification digit,

f is the number of words to be read from or punched into paper tape,

h is the length of the I/O area. h must be equal to or greater than f.

The parameters used in the above calling sequence are described in the following paragraphs.

Control Parameter

This parameter consists of four hexadecimal digits, as shown below:



I/O Function

The I/O function digit specifies the operation to be performed on a paper tape attachment. The functions, their associated digital value, and the required parameters are listed and described below.

Function	Digital <u>Value</u>	Required <u>Parameters</u>
Test Read Punch	0 1 2	Control Control, I/O Area, Error Control, I/O Area, Error
<pre>Any para particul</pre>	ameter n Lar fund	not required for a ction must be omitted.

<u>Test</u>. Branches to LIBF+2 if the previous operation has not been completed, to LIBF+3 if the previous operation has been completed.

<u>Read</u>. Reads paper tape characters into the specified number of words in the I/O area. Initiating reader motion causes an interrupt to occur when a character can be read into core. If the specified number of words has not been read, or the stop character has not been read (see "Check" in this section), reader motion is again initiated.

<u>Punch</u>. Punches paper tape characters into the tape from the words in the I/O area. Each character punched causes an interrupt which indicates that the next character can be accepted. The operation is terminated by transferring either a stop character or the specified number of words.

Check Digit

The check digit specifies whether or not word count checking is desired while completing a read or punch operation as shown below:

- 0 Check
- 1 No Check

<u>Check</u>. This function should be used with the Perforated Tape and Transmission Code (PTTC/8) only (see "Descriptions of Data Codes"). The PTTC/8 code for DEL is used as the delete character when reading. The delete character is <u>not</u> placed in the I/O area and therefore does not enter into the count of the total number of words to be read.

The PTTC/8 code for NL is used as the stop character when doing a read or punch. On a read operation, the NL character is transferred into the I/O area. On a punch operation, the NL character is punched into the paper tape.

When the NL character is encountered before the specified number of words has been read or punched, the operation is terminated. When the specified number of words has been read or punched, the operation is terminated, even though a NL character has not been encountered.

<u>No Check</u>. The read or punch function is terminated when the specified number of words has been read or punched. No checking is done for a delete or stop character.

Device Identification

When the test function is specified, the PAPTN subroutine must be told which device (reader or punch) is to be tested for an Operation Complete indication. (Remember that both the reader and the punch can operate simultaneously.) Therefore, the device identification is used only for the test function in the PAPTN subroutine. If the device-identification digit is a 0, the subroutine tests for a Reader Complete indication; if it is a 1, the subroutine tests for a Punch Complete indication.

1/0 Area Parameter

The I/O area parameter is the label of the control word that precedes the user's I/O area. It consists of a word count that specifies the number of words to be read into or punched from core. Since characters are packed two per word in the I/O area, this count is one-half the maximum number of characters transferred. Because an entire eight-bit channel image is transferred by the subroutine, any combination of channel punches is acceptable. The data can be a binary value or a character code. The code most often used is the PTTC/8 code. (See "Descriptions of Data Codes".)

Error Parameter

See "Basic ISS Calling Sequence".

PAPER TAPE SUBROUTINES (DM2 SYSTEM)

The paper tape subroutines, PAPT1, PAPTN, and PAPTX, handle the transfer of data from the IBM 1134 Paper Tape Reader to core storage and from core storage to the IBM 1055 Paper Tape Punch. Any even number of characters may be transferred via one calling sequence (PAPTX also allows an odd character count).

The PAPTN or PAPTX subroutine must be used if simultaneous reading and punching are desired. The PAPT1 subroutine will operate both devices but only one at a time. The PAPT1 and PAPTN subroutines use only a word count, reading and punching an even number of characters; PAPTX can use a word count or character count, permitting an odd number of characters to be read or punched. PAPTX allows the user to start punching from or reading into the left or right half of a word. One-frame records can be written on tape.

When called, the paper tape subroutine starts the reader or punch and then, as interrupts occur, transfers data to or from the user's I/O area. The data is packed two characters per computer word by the subroutine when reading, and must be in that form when the subroutine is called for a punch function.

Calling Sequence

Label	Operation		FT	Γ	Operands & Remarks
21 25	21. P	-	9 J	4	1 45 43 - · · ·
	L,I.B.F		+		PIA.P.T. Q. I. I. CIA.L.L. PIA.P.E.R. TAPE I. I. O.
	D.C.				/ BED de CONTROL PARAMETER
	D.C.				I.O.A.R I.I.O. A.R.E.A. P.A.R.A.MET.E.R.
	D.C.	Ц	_		ERRORERROR _PARAMETER
L	••••	Ц	1.	1	╏ ┫╴┫╴┢╼╈╴╡╌╄╕┫╸┫╼╃╴┙╶╃╶╇┍╈╴┥╴┥╴┩╴┽╴┩╶┥╺┥┑┫╕╣╕╶┥╺┥ ┑╋╌┷╸┢╷┷
			1		
	• • • •				
ERROR	D.C.				K- K
ليبينا	•••••				<u> </u>
	••••		Τ		
	•				
	8.5.C.		r	Γ	ERROR , RETURN TO CALLER
	I •				
	•				<u> </u>
	•				
I.O.A.R.	D.C.		Т		FILLER WORD COUNTILLE
	RS.S.		Τ		hun III O AREA IN INTER
	T		Т	Т	· · · · · · · · · · · · · · · · · · ·

where

- a is 1, N, or X,
- b is the I/O function digit,
- c is a check digit,
- d is the character mode digit,
- e is a device identification digit,

f is the number of words to be read from or punched into paper tape,

h is the length of the I/O area. h must be equal to or greater than f.

The parameters used in the above calling sequence are described in the following paragraphs.

Control Parameter

This parameter consists of four hexadecimal digits which are used as shown below:



I/O Function

The I/O function digit specifies a particular operation performed on the 1134/1055 Paper Tape attachment. The functions, associated digital values and required parameters are listed and described below.

Function	Digital <u>Value</u>	Required Parameters
Test	0	Control
Read	1	Control, 1/0 Area, Error
Puncn	2	Control, 1/0 Area, Error
Any para particul	ameter n lar func	ot required for a tion must be omitted.

<u>Test</u>. Branches to LIBF+2 if the previous operation has not been completed, to LIBF+3 if the previous operation has been completed.

<u>Read</u>. Reads paper tape characters into the specified number of words in the I/O area. Initiating reader motion causes an interrupt to occur when a character can be read into core. If the specified number of words has not been read or the stop character has not been read (see "Check" in this section), reader motion is again initiated.

<u>Punch</u>. Punches paper tape characters into the tape from the words in the I/O area. Each character punched causes an interrupt which indicates that the next character can be accepted. The operation is terminated either by encountering a stop character (see "Check" in this section) or by transferring the requested number of words.

Check Digit

The check digit specifies whether or not checking is desired while doing a read or punch operation. 0 Check

1 No Check

<u>No Check</u>. The read or punch function is terminated when the specified number of words or characters has been read or punched. No check is made for a delete or stop character.

<u>Check</u>. This function should be used with the Perforated Tape and Transmission (PTTC/8 Code only (see "Descriptions of Data Codes"). The PTTC/8 code for DEL will be used as the delete character when doing a read. The delete character is not placed in the I/O area and therefore is not included in the word or character count.

The PTTC/8 code for NL will be used as the stop character when doing a read or punch. On a read operation, the NL character is transferred into the I/O area and causes the operation to be terminated. On a punch operation, the NL character is punched in the paper tape and causes the operation to be terminated.

When the NL character is encountered before the specified number of words has been read or punched, the operation is terminated. When the specified number of words has been read or punched, the operation is terminated even though an NL character has not been encountered.

Character Mode

This digit is examined by the PAPTX subroutine

- If it is zero, the first word of this I/O area is interpreted as a word count.
- If it is nonzero, the first word of the I/O area is interpreted as a character count:

If the character mode digit is nonzero and even, the first character will be read into or punched from bits 0-7 of the first data word. Bits 8-15 of the last data word will not be altered if the <u>character count</u> is odd.

If the character mode digit is nonzero and odd, the first character will be read into or punched from bits 8-15 of the first data word. Bits 0-7 of the first data word will not be altered. If the <u>character count</u> is even, bits 8-15 of the last data word will not be altered.

Device Identification

When the test function is specified, the PAPTN and PAPTX subroutines must be told

which device (reader or punch) is to be tested for an Operation Complete indication. (Remember that both the reader and the punch can operate simultaneously.) Therefore, the device-identification digit is used for the test function in the PAPTN and PAPTX subroutines only; if it is a 0, the subroutine tests for a Reader Complete indication; if it is a 1, the subroutine tests for a Punch Complete indication.

1/0 Area Parameter

The I/O area parameter is the label of the control word that precedes the user's I/O area. The word count specifies the number of words to be read into or punched from the user's I/O area. Since characters are packed two per word in the I/O area, this count is 1/2 the maximum number of characters transferred. The character count, used only by the PAPTX subroutine if the character mode is nonzero is the maximum number of characters to be read or punched.

Because an entire 8-bit channel image is transferred by the subroutine, any combination of channel punches is acceptable. The data may be a binary value or a character code. The code most often used is the PTTC/8 code (see "Descriptions of Data Codes").

Error Parameter

See "Basic ISS Calling Sequence".

PLOTTER SUBROUTINE (PLOT1)

The Plotter subroutine converts hexadecimal digits in the user's output area into actuating signals that control the movement of the plotter recording pen. Each hexadecimal digit in the output area is translated into a plotter operation that draws a line segment or raises or lowers the recording pen. The amount of data that can be recorded with one calling sequence is limited only by the size of the corresponding output area.

Calling Sequence

Labe?	Operation	F	T	1			Operands	& Remarks	
<u>v</u>	F	37	33	35	*		59	35	6
	LIBF		\square	PLLOT.	1		L PILO	TTER	OUTPUT
	D.C.			1.5.0.0	<u>í</u> de	GOA	TROL	PARAS	AF.T.F.P
	D.C.			LOAR		1.10	AREA	PAR	A.M.F.T.F.P.
	0,0			ERRO	R	ERA	COR PA	RAME	T.E.P.
	- in the second			1					
	- to and			1					
	- <u></u>			4					
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	8.5.5			the com		, 11.C	AREA	<u>.</u> .	
التصحي	1		Т	[

where

b is the I/O function digit,

f is the number of words of plotter data,

h is the length of the I/O area. h must be equal to or greater than f.

The calling sequence parameters are described in the following paragraphs.

Control Parameter

This parameter consists of four hexadecimal digits, as shown below:

			1	2	3	4
			+	4	+	
1/0	Functi	i on				
Not	Used_					

I/O Function

The I/O function digit specifies the operation to be performed on the 1627 Plotter. The functions, their associated digital value, and the required parameters are listed and described below.

Function	Digital <u>Value</u>	Required Parameters
Test	0	Control
Write	1	Control, I/O Area, Error
¹ Any para	ameter n	ot required for a
particu	lar func	tion must be omitted.

<u>Test.</u> Branches to LIBF+2 if the previous operation has not been completed, to LIBF+3 if the previous operation has been completed. <u>Write</u>. Changes hexadecimal digits in the output area into signals that actuate the plotter. Figure 9 lists the hexadecimal digits and the plotter actions they represent. Figure 10 shows the binary and hexadecimal configurations for drawing the letter E.

1/O Area Parameter

The I/O area parameter is the label of the control word that precedes the user's I/O area.

The control word consists of a word count that specifies the number of computer words of data to be used.

Error Parameter

This parameter is not used but must be included because the subroutine will return to LIBF+4. (See "Basic ISS Calling Sequence".)

PLOTTER SUBROUTINE (PLOTX)

The PLOTX subroutine converts the hexadecimal digit in the parameter into a control word. The control word is stored in a buffer inside the PLOTX subroutine. One digit is transferred with each calling sequence. When the plotter is ready to accept control, the movement of the plotter recording pen is controlled by the words in the PLOTX buffer.

Hexadecimal	Plotter Action
Digit	(See Diagram Below)
0	Pen Down
1	Line Segment= +Y
2	Line Segment= +X, +Y
	Line Segment= +X
4	Line Segment= +X, -Y
	Line Segment= -X -Y
7	Line Segment -X
8	Line Segment= $-X$, $+Y$
9	Pen Up
A	Repeat the previous pen
	motion the number of times
1	specified by the next digit
	(Maximum-15 times)
в	Repeat the previous pen
	motion the number of times
· ·	digits (Maximum-255 times)
с	Repeat the previous pen
-	motion the number of times
Ì	specified by the next three
1	digits (Maximum-4095 times)
D	Not Used
E	Not Used
[F	Not Used
ļ	
	+ X I
+X,+	+X,-Y
1	
1	
1	$1 \times 1/4$
+ Y	<u> </u>
i	
İ	୬ እ
l	
- X, +	Y - X,-Y
1	- X
, L	
Figure 9. PI	OTL Control Digits





Calling Sequence

Label	Γ	Operation	Т	F	T	F			Operands &	Remarks	
n 8		77 X		12	30		25 40	45	50	55	60
		LIBF	-				PLOTX	L.C.A.L	LUPLO	TITIER.	OUTPUT
		D.C.	ł				100 de 111		TROL	PARA	

where e is the plotter control digit.

Control Parameter

This parameter consists of four hexadecimal digits:



Plotter Control

The plotter control digit specifies the recording pen action to be taken. This digit is expressed in hexadecimal.

Hexadecimal Digit	Plotter Action
0	Pen down
1	Line segment = +Y
2	Line segment = +X,+Y
3	Line segment = +X
4	Line segment = +X,-Y
5	Line segment = -Y
6	Line segment = -X,-Y
	Line segment = -X
8	Line segment = -X,+1
9 A.F	Not used
+X,+Y +Y	+X +X,-Y -Y -X

Figure 10.1 PLOTX Control Digits .

If there is no room in the buffer for the control digit, the subroutine will loop until there is room.

If the plotter is in a not-ready, not-busy condition, the subroutine exits to \$PRET where the program goes into a wait condition until operator intervention. If the plotter becomes not ready while executing the PLOTX subroutine commands, PLOTX exits to \$PST3 where the program goes into a wait condition until the operator intervenes.

The PLOTX subroutine has no errorhandling capabilities.

1231 OPTICAL MARK PAGE READER SUBROUTINE (OMPR1)

The Optical Mark Page Reader subroutine OMPR1 handles the reading of paper documents eight and one-half inches wide by eleven inches long by the 1231 Optical Mark Page Reader. A maximum of 100 words from one page can be read with one call to the subroutine.

When called to perform a read function, OMPR1 performs a feed function and reads a page into core storage according to the Master Control Sheet (see the publication <u>IBM 1231, 1232 Optical Mark Page Readers</u>, GA21-9012), and the setting of the switches on the reader. Other functions performed by OMPR1 are feed, stacker select, and disconnect.

Calling Sequence

Label		Operation		F	ĩ		Operands & Remarks
n r		v 30		32	33		31 44 45 50 33 ad as
		LIBF	_	_	L		O.M.P.R.1
Luul		D.C				L	VID. CO.B CONTRIOIL PARAMETER.
		0.0			L	1	I.O.A.R. I.I.I.I.O. A.R.E.A. P.A.R.A.M.E.T.E.R.
		ם.כ	_		L	1	ERROR ERROR PARAMETER
	_	· · · ·			-	1	<u></u>
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ERROR		D.C.,		_	L		* - *
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	_	B.S.C.		I			ERROR RETURNI TO CALLERING
		•		L			 - - - - - - - - - - - - -
		••••	_	L	1		
		•					<u></u>
I.O.A.R.		855		L	L		h

where

b is the I/O function digit,

c is the stacker select digit,

e is the timing-mark-check-test digit,

h is the length of the I/O area. h must be equal to or greater than the number of words designated to be read on the Master Control Sheet.

Control Parameter

This parameter consists of four hexadecimal digits:



I/O Function

The I/O function digit specifies the operation to be performed on the 1231 reader. The functions, their associated digital values, and the required parameters are:

Function)igita <u>/alue</u>	l Required <u>Parameters</u> ¶
Test	0	Control
Read	1	Control, I/O Area, Error
Feed	3	Control
Disconnec	t 4	Control
Stacker Select	5	Control
Any paran particula	meter ar fun	not required for a ction must be omitted.

<u>Test.</u> Branches to LIBF+2 if the previous operation has not been completed, to LIBF+3 if the previous operation has been completed.

The operation to be tested is specified by the fourth digit of the control parameter. A zero value in digit 4 specifies a normal device-busy test; that is, a test to determine if there is an operation in progress for which no operation complete interrupt has occurred. The subroutine is "not busy" once the Operation Complete interrupt takes place. A value of one for digit 4 specifies a Timing-Mark-Check-Busy test. This test indicates a "busy" condition as long as the Test-Timing-Mark-Check indicator in the Device Status Word is on. If the user wishes to run with the Timing Mark Switch set on, it is recommended that digit 4 be set to one when performing a test function.

A test function must not directly follow a feed function.

Read. Reads words or segments (response positions 1-5 or 6-10 of any word) from a document page into core storage starting at the I/O area address. The first call to OMPR1 in a program must be a read function. The read feeds the document before reading. When a read function follows a feed, the read begins with the document started by the feed. The number of bits per word read and the number of words per document read depends upon the way in which the Master Control Sheet is programmed (see the publication IBM 1231 Optical Mark Page Readers, GA21-9012). OMPR1 reads a maximum of 100 words. Any word not programmed to be read (mark positions 8 or 18 not penciled on the Master Control Sheet) is skipped. Digit 2 of the control parameter

specifies whether or not the document being read is to be stacker-selected. If digit 2 is set to one, the document is stacker-selected; if digit 2 is set to zero, it is not.

Note: On a feed, or feed as the result of a read, the document is fed from the hopper, the selected data is read into a delay line (and read out on a read), and the document continues through the machine to the stacker.

<u>Feed</u>. Initiates a feed cycle. This function advances a document from the hopper through the read station and into the stacker. Selected information from the document is stored in a delay line. A read function following a feed causes this data to be read. If a feed function is followed by another feed function without an intervening read function, the data read from the document corresponding to the first feed is overlaid in the delay line by the data read from the second document. The first call to OMPR1 in a program must not be a feed function.

A feed function must not be followed directly by a test function.

Disconnect. Terminates the read function on the data currently being read from the delay line. The subroutine-busy indicator is cleared.

Note: If the last document in the hopper is disconnected the hopper empty condition will not be detected.

Stacker Select. Performs a stacker select on the sheet currently being read (and fed), providing the stacker select function is requested while the "OK to select" bit (bit 5) is on in the Device Status Word (DSW). This bit remains on until 50 milliseconds after the read operation is completed. If the request to select arrives too late, the sheet falls in the normal stacker.

1/0 Area Parameter

The I/O area parameter is the label of the user's I/O area.

Error Parameter

There is an error parameter for the read function only. Exits are made to the user's error subroutine when the following conditions are detected:

Master Control Sheet Error

Timing Mark Error

Read Error

Hopper Empty

Document Selected

See "Basic ISS Calling Sequence" and Appendices B and C.

Feed Check

If a feed check is detected during a read or feed operation, exit is made to PST4with an error code of /A002 or /A003. After making device ready and depressing start key, OMPR1 will reinitiate the operation if error code was /A003. No stacker select will be performed on a reinitiated operation. If error code was /A002, the last document has already been processed and the operation is not reinitiated.

2250 DISPLAY UNIT MODEL 4 1/0 SUBROUTINE (DSPYN)

The 2250 I/O subroutine, DSPYN, contains the 2250 Interrupt Service Subroutine. DSPYN controls the interrupt-handling services for the 2250 Display Unit, Model The 2250 ISS contains facilities for handling attentions (graphic interrupts) from four sources; the alphameric keyboard, the programmed function keyboard, the light pen, and the graphic program itself. The DSPYN subroutine is part of the 1130/2250 Graphic Subroutine Package. A complete description of the DSPYN I/O functions can be found in <u>IBM 1130/2250 Graphic</u> Subroutine Package for Basic FORTRAN IV, GC27-6934.

RPG Subroutines (DM2 System)

The DM2 System Library contains a group of subroutines that perform functions required by the RPG Compiler and application programs. These subroutines are divided into two groups, Disk I/O and RPG Object Time Subroutines. The Disk I/O subroutines are available to Assembler language programmers. The other RPG subroutines are for system use only. All RPG subroutines are listed in Figure 24, Appendix A.

Disk File Management Subroutines (DM2 System)

Supplied with 1130 RPG is a group of disk I/O subroutines that will handle all disk file functions. These subroutines can be used by Assembler language programmers directly and are wholly independent of RPG. The subroutines provided are Direct Access, Sequential Access, and Indexed Sequential Access Method (ISAM). The subroutines are stored in the System Library.

Disk I/O Subroutines

The key to the use of the Disk I/O subroutines is an understanding of the basic principles of disk file organization and disk file processing.

FILE ORGANIZATION

File organization is the method of <u>arranging</u> data records on a direct access storage device, i.e., building the file.

The two types of file organization available with DM2 are sequential and indexed sequential (ISAM).

Sequential File Organization

A sequentially organized file is one in which records are placed on the disk in the same order they are read in, one after another. Card files are always organized this way. That is, record six cannot be written until record five is written, record five until record four, etc. Sequential files on disk may be processed sequentially or randomly.

Indexed Sequential (ISAM) File Organization

An indexed sequential file is one in which records are placed on the disk in ascending collating sequence by record key. This key may be a part number, man-number or any other identifying information that is present in the records on the file. In addition, the indexed sequential file uses an index to locate desired records. Each index entry contains a cylinder address and the highest record key on that cylinder. All index entries are formed into an index table. For cylinders that have overflowed, the index entry also contains the overflow sector address and key of the first sector overflowed from that cylinder.

Index tables are analogous to the index card file in a library. If you know the name of a book (record key), you can look in the card file (index table) until you find the card (entry) for that book. On the card you will find a number (cylinder address) where the book (record) is located. You go to the shelf (seek) and find the number (cylinder address) you are looking for. Now you can search for the particular book (record) by title (record key).

Record on an indexed sequential organized file may be processed sequentially or randomly.

FILE PROCESSING

File processing is the method of <u>retrieving</u> data records from the file, i.e., using the file. Four methods of file processing are available with DM2 RPG:

- 1. Sequential processing of sequentially organized files.
- Random processing of sequentially organized files.
- 3. Sequential processing of indexed sequential organized (ISAM) files.
- Random processing of indexed sequential organized (ISAM) files.

Sequential Processing (Sequential Files)

All records in the file are processed in order starting with the first physical record in the file.

Random Processing (Sequential Files)

In random processing the records in a file can be processed in any order. To find a record in a sequentially organized file, the record number must be supplied to the program. The record number indicates the relative position (sequential location) of the record in the file. The disk I/O routine calculates the sector address from the record number and reads the proper record.

Sequential Processing (Indexed Sequential Files)

All records in an ISAM file are available in a sequence determined by record key. Processing may start at the beginning of the file or at any point within the file.

Random Processing (Indexed Sequential Files)

In random processing the records in a file can be processed in any order. To find a random record in an ISAM file, the file index is searched using the record's key. The matching entry in the index points to the cylinder containing the record. That cylinder is then searched for the desired record. The match is again by record key. This kind of processing may be called processing in a random sequence with record keys.

SEQUENTIALLY ORGANIZED DISK ROUTINES

The sequential disk I/O subroutines provided with RPG are sequential access and direct access. A sequentially organized file is built using the sequential access routine or the direct access routine. It may be processed by either the sequential access routine or the direct access routine.

Space for the file is initially established on the disk by using a DUP STOREDATA function. STOREDATA sets aside a specified number of sectors for the file and enters the file name in LET or FLET. This file name must be used in all future references to this file.

Calculating File Size

The number of sectors needed for a file depends on record size and number of records. The records are fixed length and can be defined as any size between 320 words (640 characters) and 1 word (2 characters). Note that records cannot extend across sector boundaries. Thus a 320-word record (one sector) and a 161-word record would each require one sector of disk space. Careful planning is required in calculating optimum record size for your file. When calculating file size, always add one record for the end of file record.

To change record sizes or add records to a sequential file the file must be rebuilt. If the revised file requires additional sectors it must be redefined (*DELETE and *STOREDATA), and rebuilt.

Sequential Files	
Ranges of Record Lengths (in characters)	Records per Sector
1-2	320
3-4	160
5-6	106
7-8	80
9-10	64
11-12	53
13-14	45
15-16	40
17-18	35
19-20	32
21-22	29
21-22	26
25-24	24
23-20	22
27-20	21
27-30	20
31-34 32-34	19
33-34 35 34	10
33 -39 27.40	16
37-40	10
41-42	13
43-44	12
40-50	10
52 50	11
50.44	10
45.70	
71.00	7
/1-00	7
01-70	
91-100 107-199	5
10/-128	5
129-100	4
101-212	
213-320	
321-640	

Figure 10.2 Space Utilization for Various Size Records for Sequential Files.

Sequential Access Routine

This routine allows the programmer to store, retrieve and/or update records on a sequentially organized disk file. The space for the file must have been previously defined by the DUP function STOREDATA.

The sequence of events on a sequential access is open the file, perform the function and close the file. To accomplish these objectives the sequential access routine has three entry points:

SEQOP - open the file SEQIO - read or write a record SEQCL - close the file

The sequential access routine is a part of the System Library. It is called by a LIBF. One parameter must be passed to the routine on each call and that parameter is the address of the Disk File Information (DFI) table. This parameter must immediately follow the LIBF statement.

The coding required to process a data file using the sequential access routine is as follows:



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S.T.A.R.T		TT				
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	L 1.8.F		S.E.Q.O.P	S.E.Q.U.E.N.T.	LALL ELLLE	
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	L, D,		DIF.I.T.B.+19 L.O.A.D.	RETURN I	CIOIDIE	
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	L.D.	Ш	D.F.I.T.B.+,9, L.O.A.D.	RETURN	C.O.D.E.	
	BN.	111	ERR.R.T.N	<u>, E.R.R.O.R.</u>	ROUTLINE	I.F. NEIG
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	LD.	11-1	D.F.I.T.B	RETURN	6,0,D,E,	
	BN.	111	E.R.R.T.N	0, <u>E.R.R.O.R.</u>	R.O.U.T.I.N.E.	I.F. N.E.G
		11	 		يتنتب	
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		┝╂╌┝				
	END		SIT,A,R,T.			

Disk File Information (DFI) Table. A file to be processed by the sequential access routine must be described using a DFI table which is 11 words long. (These words are numbered 0-10.) The DFI table has nine entries, six of which must be filled in by the user. The remaining three entries must be initialized to zero by the user and are filled in by the program during execution.

Figure 11 shows the DFI table for the sequential access routine.

Operation of the Sequential Access Routine. When the routine is entered at the open entry point SEQOP, it checks the validity of the DFI table entries, sets pointers and switches to be used internally by the routine, and sets the return code in the DFI table to the code for file open. For an output function, SEQOP places the address of the record being processed in the DFI table. The routine is then entered at SEQIO to perform the required processing functions.

When the routine is entered at SEQCL, it writes the last sector of data and an endof-file record (for output files) and sets the return code to the code for file closed. The end-of-file record contains a / (slash) and an * (asterisk) in the first word. The remainder of the end-of-file record is set to binary zeros.

The sequential access routine returns to the statement immediately following the parameter that follows the LIBF to the routine for any of the three entry points.

Direct Access Routine

This routine allows the programmer to retrieve and/or update records on a

sequentially organized disk file. The records are accessed by record number relative to the beginning of the file, i.e., the first record in a file is record 1, the second record 2, etc.

The sequence of events on a direct access is open the file, perform the function, and close the file. To accomplish these objectives the direct access routine has three entry points:

DAOPN - open the file DAIO - read or write a record DACLS - close the file

The direct access routine is a part of the System Library. It is called by a LIBF. One parameter must be passed to the routine on each call and that parameter is the address of the Lisk File Information (DFI) table. This parameter must immediately follow the LIPF statement.

The coding required to process a data file using the direct access routine is as follows:

Label	Operation	Π	F	T	Operands & Remarks
0 5	<u> </u>		ы	ונ	31 4. 43 × 33 40 45 70
S.T.A.R.T		П	Τ		
1	1.1.	Π	Т	Т	(U.S.E.R. (C.O.D.E)
	L. I. B.F	П	Т	Т	DAOPN OPEN DIRECT ACCESS FILE.
	D.C	П	1		D.F.I.S D.F.I. ADD.RE.S.S. (REQUIRED)
	L.D.	П	1	1	D.F.I.2.+.9
	BN.	T	T		ERROR
		11	1	1	
	L.1.B.F	П	1	1	DALO
	D.C	П	1	1	D.F. 1.2 D.F.I. ADDRESS. (REQUIRED)
	LD.	П	1		D.F.I.2.+.9
	BN .	П	1	1	ERROR GO TO ERROR ROUTINE IF NEG
		П	7	T	
	1.1.B.F	П	1		DAGLS
	DC .	П	1		D.F.J.L. D.F.I. ADDRESS. (REQUIRED)
	L.D.	Ħ	T		DF.I.2. T. 9
	B.N.			1	ER.R.O.R
[1.1.	П	Т		
*, D.F.I	T,A,B,L	E		F	R. D.I.R.E.C.T. A.C.C.E.S.S. F.I.L.E.
DF.J.S.	D.S.A.	П	Т		F.I.L.E.A
	i.	Π	Т	Т	
E,R,R,O,R	EQU		Т	Τ	5
		П	Τ	Т	
	E.N.D.	IT	1	T	S,T,A,R,T,
	1	T	T	T	
	1	11	T	1	

Disk File Information (CFI) Table. A file to be processed by the direct access routine must be described using a CFI table which is 11 words long. (These words are numbered 0-10.) The CFI table has nine entries, seven of which must be filled in by the user. The remaining two entries must be initialized to zero by the user and are filled in by the program during execution.

Word	Entry	Meaning					
0,1,2	DSA	The first entry in a DFI table is always a DSA statement. The DSA statement allows the programmer to refer symbolically to a disk-stored data file without knowing its actual location. The label is defined as the current value of the Location Assignment Counter when the DSA statement is encountered. The operand is the name of the data file. Further information on DSA may be found in IBM 1130/1800 Assembler Language. Note: The first word of the DSA instruction is used by the sequential access routine as an update-write switch.					
3	DC /0xxx	XXX equals the number of records per sector. This figure is calculated by dividing 320 by the length of a record and ignoring the remainder. The maximum entry is $/0140$ (320 one-word records). The entry in this word must indicate the maximum number of records of X size that will fit on a sector. For example, if the entries in words 3 and 4 of this table indicate 31 ten-word records, a terminal return code of $/8014$, number of records per sector not maximum, will occur during program execution. 32 ten-word records would have to be defined to use all available disk space.					
4	DC /0XXX	XXX equals the length of the record in words. The maximum entry is /0140 (one 320-word record).					
5	DC /000X	Read/Write indicator. For read, set X to zero. For write, set X to one. For an update, set X to zero prior to the read and one prior to the write.					
6	DC LABEL	The address of the data buffer. This address must be on an even word boundary. The length of the data buffer required by the program is calculated by multiplying the number of records per sector (word 3 of this table) by the record length (word 4 of this table) and adding 2. The maximum length of the data buffer is 322 words.					
7	DC .X	Function indicator. X equals I for input, O for output and U for update. The specified character is assembled as right-justified EBCDIC. From the time a file is opened until it is closed this word must not be changed.					
8	DC /0000	Record number. This word must be reserved by the user and is filled in by the subroutine. It will contain the record number of the record being processed.					
9	DC /0000	Return code. This position must be reserved by the user. After each LIBF to any of the three entry points in the sequential access routine it should be checked for the return code. ¹					
10	DC /0000	2 /0000 Record address. This word must be reserved by the user. It will contain the address of the record being processed.					
1 Re	turn codes	s for sequential access are as follows:					
Numbe	er <u>Meaning</u>	Number Meaning					
5555	File is open	8014 Number of records per sector not maximum full 8015 File recersed when not open					
8010	0 Disk file is full 8015 File accessed when not open 1 Write indicator with input file 8016 Buffer not on even-word boundary						
8012	Read indica	ator with output file 8017 Write before read (UPDATE file)					
8013	Record size	exceeds sector size rrrr End of file OFFF File is closed					
All the op	8XXX return cod eration. Proces	les except 8017 are terminal arrors. The file must be reopened to allow program to retry ssing will again start at the first record.					
FFFI the fil	FFFF is a terminal error in the sense that it allows no further processing of the file. It does not, however, prevent the file from being closed in the normal manner.						

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Figure 11. Disk File Information Table for Sequential Access

Figure 12 shows the DFI table for the direct access routine.

Operation of the Direct Access Routine. When the routine is entered at the open entry point DAOPN, it checks the validity of the DFI table entries, sets pointers and switches to be used internally by the routine and sets the return code in the DFI table to the code for file open.

The routine is then entered at DAIO to perform the required processing functions.

When the routine is entered at DACLS, it sets the return code in the DFI table to the code for file closed.

The direct access routine returns to the statement immediately following the parameter that follows the LIBF to the routine for any of the three entry points.

INDEXED SEQUENTIAL ORGANIZED (ISAM) DISK ROUTINES

The indexed sequential disk I/O subroutines provided with RPG are ISAM load, ISAM add, ISAM sequential and random.

Indexed sequential organization gives the programmer a great deal of flexibility in the operations he can perform on a file. He can read or write records whose keys are in ascending collating sequence. He can read and update random records. (This method is not suggested if a large portion of the file is being processed since reading records in this manner is slower than reading according to a collating sequence. The index must be searched for the pointer to each record.) New records can be added to ISAM files. The add routine locates the proper positions for the new record in the file and updates the index accordingly.

ISAM has these advantages:

- It is a file management system specifically designated for direct access storage devices.
- It permits files to be processed in random or sequential order.

- It processes records directly in the I/O area.
- It establishes an index allowing ease of access to any record on the file.
- It uses an efficient chaining method to allow new records to be added to a file.
- It prevents records from being lost if a disk error occurs during an add operation.

ISAM has these restrictions:

- Records must be presorted in ascending collating key sequence before they are loaded on the file.
- Only one I/O area is permitted when a file is loaded or processed.
- All records must contain key areas starting in word one of the record, and all the key areas must be the same length.
- All records on a file must be the same length.
- Only one ISAM function can be performed on an ISAM file in one run. Hence, records cannot be both processed and added in the same run.
- The entire area for an ISAM file must be on one disk.

Contents of an ISAM File

An ISAM file comprises the following: file label, file index, prime data area, overflow area.

The relative position of these components within the ISAM file is as follows:

r 7	1	r				
File	Index	Prime	Data	Area	Overflow	Area
IT abol					over row	
Laber						1
						f

Word	Entry	Meaning		
0,1,2	DSA	The first entry in a DFI table is always a DSA statement. The DSA statement allows the programmer to refer symbolically to a disk-stored data file without knowing its actual location. The label is defined as the current value of the Location Assignment Counter when that DSA statement is encountered. The operand is the name of the data file. For more information on DSA see <u>IBM 1130/1800 Assembler Language</u> .		
3	DC /0XXX	XXX equals the number of records per sector. This entry must be the same as the number of records per sector on the file you are accessing.		
4	DC /0XXX	XXX equals the length of the record in words. This entry must be the same as the length of the records on the file you are accessing.		
5	DC /000X	Read/Write indicator. For read, set X to zero. For an update, set X to zero prior to the read and one prior to the write.		
6	DC LABEL	The address of the data buffer. This address must be on an even word boundary. The length of the data buffer required is calculated by multiplying the number of records per sector (word 3 of this table) by the record length (word 4 of this table) and adding 2. The maximum length of the data buffer is 322 words.		
7 8	DC /0XXX DC /XXXX	Record number. Word 7 and 8 must contain the number of the record on which the operations are to be performed. This number is equivalent to the record's relative location in the file; hence, the 83rd record would be record number 83. The entry is right-justified hexadecimal. Therefore, word 7 will be all zeros for all record numbers less than 65,536. The direct access routine will convert the record number supplied to the actual disk address.		
9	DC /0000	Return code. This word must be reserved by the user. After each LIBF to any of the three entry points in the direct access routine it should be checked for the return code. ¹		
10	DC /0000	Record address. This word must be reserved by the user. It will contain the address of the record being processed.		
Ret	turn codes	for direct access are as follows:		
Hez <u>Nur</u>	Hexadecimal <u>Number Meaning</u>			
5555 File open 8000 Record number not in file 8001 Record size not within limits 8002 Records per sector not maximum 8003 Record number not positive 8004 Write before read 8005 File accessed when not open 8006 Buffer not on even-word boundary 0FFF File closed All 8XXX return codes except 8000, 8003 and 8004 are terminal errors. The file				
must star	be reoper t at the f	ed to allow the program to retry the operations. Processing will again is record.		

Figure 12. Disk File Information Table for Direct Access

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ISAM File Label. The first sector of any ISAM file contains the file label. This label contains information required by the ISAM routines for all future processing of the file. The file label is built by the ISAM load function, updated by ISAM add, and used by ISAM random and sequential. All label operations are performed automatically by the ISAM routines. The user need perform no label operation other than reserving one sector for the label when the file is initially defined.

The format of the ISAM label is shown in Figure 13.

Word Number	Label Entry Description
1 1	Key length
2	Record length
3	Number of index entries per
1	sector
4	Index entry length
5	Number of records per
	sector
6	Record number of last prime
1 1	data record
7	Index entry number of last
	entry in file
8	Sector address of last
	prime data record
9	Sector address of last
	Index entry
1 10	Sector address of next
1 11	Overilow record
	Record number of next
l 	Overriow record

Figure 13. Format of an ISAM Label

<u>ISAM File Index</u>. The ability to read or write records anywhere in a file is provided by the file index. An entry in this index contains a cylinder address and the highest key that is associated with that cylinder. The ISAM routines locate a given record by searching the index for the key and then searching the specified cylinder for the desired record, again searching by key. To increase the efficiency of the ISAM routines, one sector of the index is retained in core storage for each file.

The key may be a part number or an employee name or any other identifying information that is contained in any record of the file. The key entries in the index are the numbers of the highest key on each cylinder in ascending collating sequence. The end-of-record key is the key with the highest possible value, i.e., all bits are ones.

A portion of an index or index table is shown below. Note that each entry contains two sets of the same information. The second set is overlayed to show overflow data when the affected cylinder overflows.

Key 15	First cylinder address	Ke y 15	First cylinder address	Zeros	Кеу 30	Second cylinder address	Key 31	Overflow sector address	Record number
								<u> </u>	

normal entry

	Key 45	Third cylinder address	Key 45	Third cylinder address	Zeros		all 1 bits	nth cylinder address	all 1 bits	nth cylinder address	Zeros
--	-----------	------------------------------	-------------	------------------------------	-------	--	--------------------	----------------------------	----------------------	----------------------------	-------

normal entry

overflow entry

last entry in index

<u>Prime Data Area</u>. This area contains the data records placed in the file by the ISAM load routine. The records must all be the same length (maximum 318 words). ISAM adds a two-word control field to each record. This control field, called the sequence-link control field, is used in the overflow area as a chaining indicator. It is used in the prime data area to indicate whether or not a cylinder has overflowed.

	Data	record	1	word	1	word	
Кеу 75			2	eros	Ze	eros	
					~		7
			S	equence	ce-	link ield.	

Data record on a prime data cylinder.

Data record 1 word 1 word Key Zeros X'FFFF' 520

Sequential-link control field.

Last data record on prime data cylinder that has overflowed.

Overflow Area. When a new record is added to an indexed sequential file, it is placed according to key sequence. If records were to remain in precise physical order, the insertion of each new record would require all records with higher keys to be shifted up. However, because ISAM files have an overflow area, a new record can be entered into its proper position on a cylinder and only cause records with higher keys on that cylinder to be shifted. The record that is forced off the end of the cylinder by the addition of the new record is written in the overflow area.

The index entry of any cylinder that has overflowed points to the overflow sector address and record number of the overflowed record in the overflow area. If two or more records in key order are added, the overflowed records are chained together in the overflow area through the entries in their sequence-link control field. The entry in the first record points to the second, the second to the third, etc. The last overflow record in the chain has a sequence-link control field of all zeros.

The number of cylinders to be allotted to the overflow area must be determined by the programmer when the file is initially defined. Records are placed in the overflow area in the order they have overflowed, not in key sequence.

To illustrate the overflow area, assume that on cylinder six of a defined file the last three entries have keys 150, 152 and 154. Key 154 would identify cylinder six in the index. Now we add a record with key 153, a record on another cylinder and a record with key 151. The overflow area would appear as shown below. Key 152 would identify cylinder six in the index. The overflow entry for cylinder six in the index would point to the overflow area.

Overflow area.

Кеу 154	Zeros	Zeros	Zeros	Zeros	Key 153	Overflow sector address	Rec. 0001
			 		<u> </u>		

First record overflowed. The sequencelink control field is zeros indicating the end of a chain. Record overflowed from another cylinder

Last record overflowed. The sequence-link control field points to the next key in sequence. In this case it is key 154 in the overflow area.-

Creating and Using ISAM Files

An indexed sequential file is built using the ISAM load routine, is expanded using either the ISAM sequential or ISAM random routine.

Space for the file is initially established on the disk by using a DUP STOREDATA function. STOREDATA sets aside a specified number of sectors for that file and enters the file name in LET or FLET. This file name must be used in all future references to this file. Key length is a maximum of 25 words (50 characters). If the length of the key in characters is odd, add one when calculating the number of words, i.e., 49 characters require 25 words.

Overflow sectors = The number of sectors the user wishes to allot to record overflow before the file must be rebuilt. The overflow area is automatically assigned to start at the sector following the last sector of prime data. This assignment is done by the ISAM load (close) routine.

Determining ISAM File Size

The number of sectors required for an ISAM file is computed by the following formula (the remainder in all cases should be disregarded):

Prime data sectors + Index sectors + Overflow sectors + 1 (File label)

where:

Prime data sectors =

Approximate number of records in file + number of records per sector - 1 Number of records per sector

Number of records per sector = 320Record size + 2

The maximum record size is 318 words. Records cannot cross sector boundaries.

Index sectors =

Number of prime data cylinders + number of index entries per sector - 1 Number of index entries per sector

Number of prime data cylinders =

number of prime data sectors + 7

Number of index entries per sector=

320 Index entry size

8

Index entry size = 2 (key length in
words) + 3

When computing file size, always add one sector for the file label.

If desired, an Assembler language program can be used to perform the above calculations. The programmer need only know the index entry size (calculation shown above), the length of a record in words, the approximate number of records in the file and an estimate of the number of sectors of overflow area needed.

Key Length in Characters	Number of Entries on One Sector	Number of file Sectors Accomodated in One Index Sector
1-2	64	512
3-4	45	360
5-6	35	280
7-8	29	232
9-10	24	192
11-12	21	168
13-14	18	144
15-16	16	128
17-18	15	120
19-20	13	104
21-22	12	96
23-24	1 11	88
25-28	10	80
29-30	9	72
31-34	8	64
35-38	7	56
39-44	6	48
45-50	6	48
	1	1

Figure 13.1 ISAM Cylinder Index Chart

Indexed Sequential Files						
Ranges of Record Lengths (in characters)	Records per Sector					
1-2	106					
3-4	80					
5-7	64					
7-8	53					
9-10	45					
11-12	40					
13-14	35					
15-16	32					
17-18	29					
19-20	26					
21-22	24					
23-24	22					
25-26	21					
27-28	20					
29-30	18					
31-32	17					
33-36	16					
37-38	15					
39-40	14					
41-44	13					
45-48	12					
49-54	11					
55-60	10					
61-66	9					
67-76	8					
77-86	7					
87-102	. 6					
103-124	5					
125-156	4					
157-208	3					
209-316	. 2					
317-636	1					
637 -64 0	Invalid					

Figure 13.2 Space Utilization for Various Size Records for Indexed Sequential Files. A program to calculate all values computed above is shown in Appendix J of the <u>IBM 1130 Disk Monitor System</u>, Version 2, Programmer's and Operator's Guide.

ISAM Load Routine

This routine loads presorted records, one after another, into the prime data area of the file. As each prime data cylinder is filled the load routine creates an entry in the file index. After all records are loaded in the prime data area the load routine creates the end of file record and the last index entry. The key for end of file and last index entry are all one bits.

The sequence of events on an ISAM load is open the file, perform the function and close the file. To accomplish these objectives the ISAM load routine has three entry points.

ISLDO - open the file ISLD - write a record ISLDC - close the file

The ISAM load routine is a part of the System Library. It is called by a LIBF. One parameter must be passed to the routine on each call and that parameter is the address of the Disk File Information (DFI) table. This parameter must immediately follow the LIBF statement. The coding required to build a data file using the ISAM load routine is as follows:

Label	Operation	F	T	Γ	Operands & Remarks
n 15	$1c \rightarrow 1$	12	н, 1		ы 60 45 с. – с
START			Γ	L	<u></u>
	. i			Γ	1
	L.1.8.F			L	1, S, L.D. O. ,,, O, P, E, N I, S, A, M. , L.O, A, D. , F. J, L, E
	D. C.		1	L	D.F.I.A.D. , D.F.I. ADD.R.E.S.S. (RE.QUI.R.E.D.)
	L.D.			L	DIF. 1. A.D. +, 9
	BN.			L	ERROR GO. TO. ERROR ROUTINE IF. NEG
	احتضار		1	L	
	L.1.B.F	Ц.	1	L	I.S.L.D. WRITEL RECORD
	D, C,				D.F.I.A.D. D.F.I. ADD.R.E.S.S. (REQUIRED)
	L,D, ,		1	L	D.F.IAD.+.9. LOAD, RETURN CODE
	BN.			L	ERROR
			L		
	L 1 . B.F			L	1.5. L.D.C
	D.C.				D.F.I.A.D. , D.F.I. ADDRESS, (REQUIRED),
	LD,				DF.I.AD. +19 LOAD RETURN CODE
	BN.	T	Т	Γ	ERROR GO. TOLERROR ROUTINE IF. NEG
	1		Г	Γ	
*, D.F.I	T,A,B,L	E	F	a	R. 1.S.A.M. L.O.A.D. R.O.U.T.I.N.E.
D.F.J.A.D	D.S.A.		Т	Γ	FILLEA USER FILLE
			Т	Γ	
E.R.R.O.R	E,QU		Т	Γ	*
	1.5			Γ	
	E,N.D.				S,T,A,R,T,
	1		Т	Γ	
		_	+	•	

<u>Disk File Information (DFI) Table</u>. A file to be loaded by the ISAM load routine must be described using a DFI table which is 21 words long. (These words are numbered 0-20.) The DFI table has nineteen entries, eleven of which must be filled in by the user. The remaining eight entries must be initialized to zero by the user and are filled in by the program during execution.

Figure 14 shows the DFI table for the ISAM load routine.

Operation of the ISAM Load Routine. When the routine is entered at the open entry point ISLDO, it checks the validity of the DFI table entries, sets pointers and switches to be used internally by the routine and sets the return code in the DFI table to the code for file open.

The routine is then entered at ISLD to load a record to the file.

When the routine is entered at ISLEC, it writes the last record in the prime data area, an end-of-file record, the last index entry, and sets the return code to file closed. The end of file record contains all one bits.

The ISAM load routine returns to the statement immediately following the parameter that follows the LIEF to the routine for any of the three entry points.

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	Word	Entry	Meaning						
	0,1,2	DSA	The first entry in a DFI table is always a DSA statement. The DSA statement allows the programmer to refer symbolically to a disk stored data file without knowing its actual location. The label is defined as the current value of the Location Assignment Counter when the DSA statement is encountered. The operand is the name of the data file. For more information on DSA see <u>IBM 1130/1800 Assembler Language</u> .						
	3		XX equals the key length in characters. Maximum is /0032. (50 characters)						
	4	DC /0XXX	XXX equals the length of the record in words. The maximum entry is /0140 (one 320-word record). This includes the two words required for the sequence-link control field.						
	5	DC LABEL 1	The address of the index buffer. This address must be on an even-word boundary. The length of the index buffer is calculated by multiplying the number of index entries per sector by the index entry length and adding 2. The maximum length of this buffer is 322 words.						
	6	DC LABEL 2	The address of the data buffer. This address must be on an even-word boundary. The length of the data buffer is calculated by multiplying the number of records per sector (word 14 in this table) by the record length (word 4 in this table) and adding 2. The maximum length of the data buffer is 322 words.						
	7	DC /XXXX	Routine type code. For ISAM load, XXXX = 1111.						
	8	DC /XXXX	XXXX equals the number of sectors required for the index. See "Determining ISAM File Size" in this section for the methods used to calculate this value.						
	9	DC /0000	Return code. This word must be reserved by the user. After each LIBF to any of the three entry points in the ISAM load routine it should be checked for the return code.						
	10	DC /0000	Address of record being processed. This word must be reserved by the user.						
í	1 Ret	turn codes	for ISAM load are as follows						
į	Hex	adecimal							
	Num	ber <u>Meaning</u>							
i	5555	5 File is ope							
j	8020) Not a load Record siz	1 function						
	8022	2 Key lengt	h greater than maximum						
	8023	3 Index entr	y length not same as length computed from key length						
i	8024	1 Number o	f index entries per sector incorrect						
į	8025	o rrime dat 5 Index area	a area is mil						
ļ	8027	7 File is not	: open						
	8028	3 Index buff	er not on even-word boundary						
	8029	Data buffe	er not on even-word boundary						
İ	8027 0FFI	A Input reco F File is clo	rd out of sequence sed						
ļ	All 8	XXX return code	s except 802A are terminal errors. The file must be reopened to allow the program to retry the						
	operatio	on, Processing w	rill again start at the first record.						
i									

Figure 14. Disk File Information Table for ISAM Load (Part 1 of 2)

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Word	Entry	Meaning
11	DC /0000	Address of the index entry. This word must be reserved by the user.
12	DC /XXXX	XXXX equals the number of index entries per sector. See "Determining ISAM File Size" in this section for the methods used to calculate this value. This value must be the maximum number of index entries that will fit on a sector.
13	DC /XXXX	XXXX equals the index entry length in words. See "Determining ISAM File Size" in this section for the methods used to calculate this value.
14	DC /XXXX	XXXX equals the number of records per sector. See "Determining ISAM File Size" in this section for the methods used to calculate this value. The entry in this word must indicate the maximum number of records that will fit on a sector.
15	DC /0000	Prime data record number. This word must be reserved by the user.
16	DC /0000	Index entry number. This word must be reserved by the user.
17	DC LABEL 3	Address of key hold area. This area is used to hold the key of the previous record so the records can be sequence checked. After the close routine has been executed this word will contain the sector address of the last prime data sector. The key hold area must be as many words long as there are characters in the record key.
18	DC /0000	Sector address of last index sector. This word must be reserved by the user.
19	DC /0000	Sector address of next overflow sector. This word must be reserved by the user.
20	DC /0000	Record number of next overflow record. This word must be reserved by the user.

Figure 14. Disk File Information Table for ISAM Load (Part 2 of 2)

ISAM Add Routine

This routine allows the user to add records to an existing file. The new records are placed in proper order by key sequence in the prime data area. The records forced off the prime data cylinders by the new records are placed in an overflow area. If the record to be added logically falls between the last record presently on the cylinder and the last record originally on the cylinder, it is written directly into the overflow area. If the record being added has a higher key than any record on the file, it is inserted before the end of file record. The add routine will operate most efficiently if the records being added are presorted by key sequence.

It is extremely important that an Add file be <u>closed</u>. This is to insure the file is properly updated for future processing. The add file should be closed before termination of the job as a result of either normal or abnormal EOJ. If the job is abnormally terminated because of a CPU failure or a DASD error (indicated by error code /5004 with DISKZ) when an ADD is being performed, it is possible that a duplicate record may have been generated on the file. If this occurs, the user should check his file and if such a duplicate record exists, it should be deleted.

The sequence of events on an ISAM add is open the file, perform the function and close the file. To accomplish these objectives the ISAM add routine has three entry points:

ISADO - open the file ISAD - write a record ISADC - close the file

The ISAM add routine is a part of the System Library. It is called by a LIBF. One parameter must be passed to the routine on each call and that parameter is the address of the Disk File Information (DFI) table. This parameter must immediately follow the LIBF statement. The coding required to add records to an ISAM file is as follows.

label	Operation	ł	F	I۲		Operande & Remarks
ກ ຮ	<u>a</u> 1		×	33		35 40 43 20 55 60 65 70
S.T.RR.T	1					
1						1, (.U.S.E.R., C.O.D.E.),
	4.1.8F	Γ				I.S.A.DO OPEN ISAM P.D.D. FILE.
	P.C.			Г	Γ	D.F.J.A.D
	2,2,	Г		Г	Г	D.F.I.A.D. 1.9
	8.N.	Г		Γ	T	E.R.S.O.S
		t		ſ	t	
	LZAF		Г		T	T.S.O.D
	DC	t		Γ	t	DFIRD, $DFI (DDDRESS (0FAUTRED))$
	1.2	t	F	F	t	DE TODAS
	8.11	1		F	t	EPPAR
				t	t	
	TAF		H	F	t	TSAJA ALASE TEAM OND BELE
	De	ŀ	-	-	t	$\frac{1}{2} \left[\frac{1}{2} \left$
	1 7	t		\vdash	t	$D = T = D + g \qquad (A = D) = g = T + g = g = g = g = g = g = g = g = g = g$
	10 V			⊢	t	50 2 4 9 11 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
				⊢	t	
	1 2 2 2	t_		-	١.	
AF - 2	1.4.131	E	⊢	۴	10	K ASHAL HOD KOUTANE
NET BO	12128	-	_	⊢	÷-	FIT.LEALT USER FILE
<u></u>	12.1			Ļ.	Ļ.	<u></u>
E.R.R.O.R	E.Q.U.			L	۴.	A LALA LA
					L	
	END	1			L	S.T.A.R.T. L. L. L. L. L. L. L. L. L. L. L. L. L.
ليبين						
		Π			Γ	
		Γ			Γ	
		-	-	-	-	· • • • • • • • • • • • • • • • • • • •

Disk File Information (DFI) Table. The ISAM add routine requires a DFI table describing the file. The DFI table (which is 21 words long, numbered 0-20) has nineteen entries, six of which must be filled in by the user. The remaining thirteen entries must be initialized to zero by the user and are filled in by the program during execution.

Figure 15 shows the DFI table for the ISAM add routine.

Operation of the ISAM Add Routine. When the routine is entered at the open entry point ISADO, it checks the validity of the DFI table entries, sets pointers and switches to be used internally by the routine and sets the return code in the DFI table to the code for file open.

The routine is then entered at ISAD to add a record to the file.

When the routine is entered at ISADC, the label is updated and the return code is set to file closed.

The ISAM add routine returns to the statement immediately following the parameter that follows the LIBF to the routine for any of the three entry points.

ISAM Sequential

The ISAM sequential routine is used to retrieve and update records on an ISAM

file. Processing may start at the first record or at any record within the file.

The programmer can update each record immediately after it is processed by writing it back to the same location from which it was retrieved. This update is accomplished by specifying /0010 in word 7 of the DFI table when the file is opened and modifying word 19 of the DFI to /0001 before issuing the LIBF ISEQ. Word 19 must be restored to /0000 prior to reading the next record. An update is not required if the records are not changed.

The sequence of events for an ISAM sequential operation is open the file, set a low key limit if required, perform the function and close the file. To accomplish these objectives the ISAM sequential routine has four entry points:

ISEQ0 - open the file ISETL - set low key limit (start processing at this record) ISEQ - process a record ISEQC - close the file

The ISAM sequential routine is a part of the System Library. It is called by a LIBF. One parameter must be passed to the routine on each call and that parameter is the address of the Disk File Information (DFI) table. This parameter must immediately follow the LIBF statement.

The coding required to retrieve and update records on an ISAM file starting with the first record is as follows:

Lobel	Operation	FT	Operands & Researks
·	· · · · ·	5 1	30 45 <u>5 35 40 45 5</u>
S.T.A.R.T			
1.4.1.1			(,,,,,,,,,,.
	LIBE		TISEGO OPEN ISAM SEGUENTLAL
	BC .		D.F. LAD DF.I. ADD. R.E.S.S. (REAUIRED)
	1.2.		DEF.I.AD.19
	8.1.		EREDO
	1.TAR	+++	TREA
	DA		DETAIL DET ODDECC (DEAUTARD)
┽┷┹╍┺╋	LAC.		
	12.4	+-+-+	$D_i F_i L_i H_i A_i Y_i , \dots L_i D_i H_i A_i R_i A_i T_i U_i H_i A_i \dots C_i D_i D_i A_i \dots \dots \dots \dots \dots \dots \dots \dots \dots \dots \dots \dots \dots \dots \dots \dots \dots \dots \dots$
	15/1	++-	ERROR I I IGO ITO ERROR ROOM TILAE ILE MEG
+++++++++++++++++++++++++++++++++++++++	++	+ + + + + + + + + + + + + + + + + + +	╺╬╍┺┉┶┉┙┛╸┍╶┙┙┚┚╘╘╘╘╘╘╘╘╘╘╘╘╘╘╘╘╘╘╘╘╘╘╘╘╘╘╘╘╘
<u></u>	4.I.O.F	444	IS.E.G.C
	DC.	4	P.F.J.A.D
<u> </u>	4.2		DF.Z.AD.+19
الغبب	6.N.	11	ERROR GO. TO. ERROR ROWTINE IF. NEG
	1.6.1		
. D.F.I.	TABL	E F	OR. I.S.B.M. SEGUENTIZAL.
P.F.I.A.D	D.S.A.	LT.	F.I.L.E.R
		ΤП	
E.R.R.O.R	EQU		*
	E.U.D.		START 1
	+		<u>·</u> ···································

Word	Entry	Meaning					
),1,2	DSA	The first entry in a DFI table is always a DSA statement. The DSA statement allows the programmer to refer symbolically to a disk stored data file without knowing its actual location. The label is defined as the current value of the Location Assignment Counter when the DSA statement is encountered. The operand is the name of the data file. Further information on DSA may be found in <u>IBM 1130/1800</u> Assembler Language.					
3	DC /00XX	XX equals the key length in characters. Maximum is $/0032$ (50 characters).					
4	DC /0XXX	XXX equals the length of the record in words. The maximum entry is $/0140$ (one 320-word record).					
5	DC LABEL 1	The address of the index buffer. This address must be on an even word boundary. The length of the index buffer is calculated by multiplying the number of index entries per sector by the index entry length and adding 2. The maximum length of this buffer is 322 words.					
6	DC LABEL 2	The address of the record being added to the file.					
7	DC /XXXX	Routine type code. For ISAM add, XXXX = 0000.					
8	DC /0000	Index entry number in process. This word must be reserved by the user.					
9	DC /0000	Return code. This word must be reserved by the user. After each LIBF to any of the three entry points in the ISAM add routine it should be checked for the return code.					
10	DC /0000	Prime data record number in process. This word must be reserved by the user.					
• Ret	turn codes	for ISAM Add are as follows.					
Hexa	adecimal						
Num	ber Meaning						
5555	5 File is ope	en de la companya de la companya de la companya de la companya de la companya de la companya de la companya de					
8030) Not an ad	d function					
8031	Keylengt	: open h in DFI table not same as key length in label					
8032	3 Record le	ngth in DFI table not same as record length in label					
8034	8034 Key is presently on file						
8035	8035 Overflow area is full						
8036	5036 Index builter not on even-word boundary OFFE File is closed						
OFFI	r rite is cit	iseu .					
All 8XXX return codes except 8034 are terminal errors. The file must be reopened to allow the program to retry the operation. Processing will again start at the first record.							

Word Entry Meaning 11 DC /0000 Address of the index entry. This word must be reserved by the user. DC /0000 12 Number of index entries per sector. This word must be reserved by the user. 13 DC /0000 Index entry length in words. This word must be reserved by the user. Number of records per sector. 14 DC /0000 This word must be reserved by the user. 15 DC /0000 Record number of last prime data record processed. This word must be reserved by the user. 16 DC /0000 Number of last index entry for file. This word must be reserved by the user. 17 DC /0000 Sector address of last prime data sector. This word must be reserved by the user. 18 DC /0000 Sector address of last index sector. This word must be reserved by the user. 19 DC /0000 Sector address of next overflow sector. This word must be reserved by the user. Record number of next overflow record. This word must be reserved by 20 DC /0000 the user.

Figure 15. Disk File Information Table for ISAM Add (Part 2 of 2)

The coding required to retrieve and update records on an ISAM file starting at a record other than the first record on the file is as follows. (Note that the record key must be placed in the key hold area in a special format: one character per word, occupying the rightmost eight bits of the word. The coding included below illustrates one way that this can be done. It is required only if the key characters are left-justified.)

Disk File Information (DFI) Table. The ISAM sequential routine requires a DFI table describing the file. The DFI table (which is 21 words long, numbered 0-20) has nineteen entries, eight of which must be filled in by the user. The remaining eleven entries must be initialized to zero by the user and are filled in by the program during execution.

Figure 16 shows the DFI table for the ISAM sequential routine.

Lobel		Operation	Π	F	T	Γ	Operands & Remarks
e 1		2 1		s:	13		ы <u>6 6 т. </u>
S.T.A.B.T							
			Γ			Г	1
		L.I.B.F					ISEQUENTIAL OPEN ISAM SEQUENTIAL
		D.C	Π			Γ	D.F.I.A.D D.F.I. ADDRESS (REQUIRED)
		L.D	Γ			Γ	DIF. LAD. + 9 LOAD, REITURN CODIE,
		B.N.	Γ			Г	ERROR GO. TO ERROR ROUTINE IF. NEG
			Г			F	
4. 9.5.4	5	. 1.N.	Ŀ.	0	w	Г	LIMIT, ADDRESS, THE FOILLOWING EXAMPLE
#. S.H.O	~	SO.N	£		м	ε	THOD. OF CONVERTING THE KEY TO THE
+F.O.R	м	8.5.0		1	8	ε	D. B.Y. J.S.E.T.L.
			Г		r	r	
		L.D.	T		Г		DELADITAL KEY LENGTH IN CHARACTERS
		STO	Г	,	Г		1
		L.D.			\vdash	Г	DELAD+S
		6.T.O			Г	t	2
	1	SL T.		-	h		3.1. ZPRO REGISTERS
1000		1 7	T		2	t	0 JOAD WORD OF KEY
	H	C.R.A		H	F		C. SHIET CHARACTER TO BU WORD
		STO		H	,		A PESTORE WORD OF KEY
		MD.Y.	T		12	T	I
	Η	MD.X	T	F	1		DECREMENT CHARACTER COUNT
		MD.V.	t	1	r		LOOR J.E. NON TERO CONTINUE
	Н		r		F	L	
T. FND		OF. F	1.	6	5	6	LE. THE FOLLOWING CODING IS REQUIRED
).T.8.F	r	ľ	Ĩ.	ľ	T.S.F.T.L. SET LOW KEN LIMIT
		D.C.	t	1	t	t	DELAD DEL ADDRESS (REQUIRED)
	-	1 1 1 1		-	-	<u> </u>	LAT A PUMPLE A TOTAL REPORT OF A PUPPLE AND A PUPPLE A PU

· · · · · ·	1	TT.		·					
Label	Operation	11			Operands &	Remorks			
21 25	2/ 30	14/12	1 ³⁹	40 45	50			62	70
··· ···	L.O. ,		DELAD	<u>419i</u>	0.A.DR.E.	T.U.R.N.	C.O.D.E.	ب بالبين	
	B.N.		ERROR	G u	O. TO E	RIRIQIR	ROVITI	N.E. 11.F.	. N.E.G
			11						
	L.1.8.F		1.5.6.0.		E.A.D. 0.8.		E. L.F.C.	3.8.D.	
	00	П	DELAD	0	E.I. A.D.D.	9.F.S.S.	(.8.6.0.11	18.6.0.)	
		++	0.5	LO		T A .1	CADE		
Freed		++-	DIF.I.M.D	<u> </u>	UAD INC.	LUKIN	COULE		
·····	1041-1	++	FIR'R'D'B	فكالأستين فيشتك	3, 10, 1, 10	R.K.O.R.	KONITI	عليل بسيعه الا	N.E.6
	4 6 1 1	++-		بلبتبيات				بي السب	للسبب
	-Huli B.F	++-	LISEGE	سکر در در دار ا	LOSEI	<u>5 A.N5</u>	E.Q.U.E.N.	TILAL	
	DC.		D.F.I.A.D	باهتينيات	F.I. A.D.O.	RESIS	(R.E.G.U.	(Gr3.8.1	
	1.0.		DIF.I.A.D	+19	O.A.D. R.E.	TNRIN	CODE .		
	6.N.		ERROR	.	0. TO	R.R.O.R.	RA.U.T.L	N.E. ILE	NE
		ТΓ							
F DEI	TABL		00. 1.54	N. TEAUE		<u> </u>			
		• r						_	
7 12-112-0	10.3.12	╉╋┉	PULLEA.	يلاب حجيجات	بابط بهرهرج				السعيب
	4-4-4-4	++-		بالتشييات	بالمستعدية				ـــــ
E.R.R.O.R	16.0.0			L L L L L L	R.R.O.RR.	ONTIA		A 4.4 A	
ليتنبط	المشال								
	END		START						

r			T							
Wor	d	Enti	су	Meaning						
0,1	,2	DSA		The first entry in a DFI table is always a DSA statement. The DSA statement allows the programmer to refer symbolically to a disk stored data file without knowing its actual location. The label is defined as the current value of the Location Assignment Counter when the DSA statement is encountered. The operand is the name of the data file. Further information on DSA may be found in <u>IBM 1130/1800</u> <u>Assembler Language</u> . Note that the first word of the DSA instruction is loaded with the last prime data sector address when the file is openel.						
3		DC /	/00XX	XX equals the key length in characters. Maximum is /0032. (50 characters). The key length must be the same as the key length in the file label.						
4 4 		DC /	⁄0XXX	XXX equals the length of the record in words. The maximum entry is //0140 (one 320-word record). The record length must be the same as the record length in the file label.						
5		DC I	LABEL 1	The address of the index buffer. This address must be on an even-word boundary. The length of the index buffer is calculated by multiplying the number of index entries per sector by the index entry length and adding 2. The maximum length of this buffer is 322 words.						
6		DC I	LABEL 2	The address of the data buffer. This address must be on an even-word boundary. The length of the data buffer is calculated by multiplying the number of records per sector by the record length and adding two. The maximum length of the data buffer is 322 words.						
7	ļ	DC /	/XXXX	Routine type code. For ISAM sequential retrieve, XXXX = 0001. For ISAM sequential update, XXXX = 0010.						
8		DC /	/XXXX	Address of key hold area if processing starts at a point other than the first record in the file. If the entire file is being processed, this word must be /0000.						
9		DC /	/0000	Return code. This word must be reserved by the user. After each LIBF to any of the four entry points in the ISAM sequential routine it should be checked for the return code. ⁴						
10		DC /	/0000	Address of record in process. This word must be reserved by the user.						
1 1		urn	codes	for ISAM sequential are as follows:						
i :	Hexa	decim	al							
ļ .	Num	ber 1	Meaning	İ						
	3355 8040	1	nie is open Not a semie	ntial retrieve or update function						
i	8041	1	ndex buffer	not on even-word boundary						
!	8042	I	Data buffer	not on even-word boundary						
	8043	ł	Key length in DFI table not same as key length in label							
1	8044 8045	I	Record length in Dri table not same as record length in label							
i	8046	1	Write before read on update							
1	FFFF	I	End of file							
	OFFF	I	File is close	d						
Į.	A11 8	XXX re	eturn codes	except 8046 are terminal errors. The file must be reopened to allow the program						
to i	retry	the op	eration. Pr	ocessing will again start at the first record.						
I I I I I I I I I I I I I I I I I I I	FFFF :	is a te i from L	rminal error	in the sense that it allows no further processing of the file. It does not, however, prevent						
L										

Figure 16. Disk File Information Table for ISAM Sequential (Part 1 of 2)
Word	Entry	Meaning
11	DC /0000	Address of the index entry used to locate the record. This word must be reserved by the user.
12	DC /0000	Number of index entries per sector. This word must be reserved by the user.
13	DC /0000	Index entry length in words. This word must be reserved by the user.
14	DC /0000	Number of records per sector. This word must be reserved by the user.
15	DC /0000	Update-write indicator. This word must be reserved by the user.
16	DC /0000	Number of index entry in process. This word must be reserved by the user.
17	DC /0000	ISETL switch to indicate low-limit record found. This word must be reserved by the user.
18	DC /0000	Internal switch used to indicate that last record in overflow area has been found. This word must be reserved by the user.
19	DC /XXXX	Read/Write indicator. If routine type code (word 7 of this table) was a retrieve, this entry should be set to /0000. If word 7 indicates an update file, this entry should be /0000 for the retrieve and /0001 for the update. This word should be reset to /0000 before the next retrieve.
20	DC/0000	Prime data record number in process. This word must be reserved by the user.

Figure 16. Disk File Information for ISAM Sequential (Part 2 of 2)

Operation of the ISAM Sequential Routine. When the routine is entered at the open entry point ISEQO, it checks the validity of the DFI table entries, sets pointers and switches to be used internally by the routine and sets the return code in the DFI table to the code for file open.

If processing is not to start at the first record in the file, the routine is entered at ISETL to locate the starting record.

The routine is then entered at ISEQ to perform the processing functions.

When the routine is entered at ISEQC, the last record is processed and the return code is set to file closed.

The ISAM sequential routine returns to the statement immediately following the parameter that follows the LIBF to the routine for any of the four entry points.

ISAM Random

The ISAM random routine is used to retrieve and update records randomly on an ISAM file. The programmer first places the key field of the desired record in a user-defined area. ISAM random then searches the index to locate the cylinder containing the desired record and then searches that cylinder for the record. The sector containing the record is then read and that record is made available for processing.

The programmer can update each record immediately after it is processed by writing it back to the same location from which it was retrieved. This update is accomplished by specifying /1000 in the seventh word of the DFI table when the file is opened and modifying the nineteenth word of the DFI to /0001 before issuing the LIBF ISRD. The nineteenth word must be restored to /0000 prior to reading the next record. An update is not required if the records are not changed.

The sequence of events for an ISAM random operation is open the file, perform the function and close the file. To accomplish these objectives the ISAM random routine has three entry points.

ISRDO - open the file ISRD - process a record ISRDC - close the file

The ISAM random routine is a part of the System Library. It is called by a LIBF. One parameter must be passed to the routine on each call and that parameter is the address of the Disk File Information (DFI) table. This parameter must immediately follow the LIBF statement.

The coding required to retrieve and update records on an ISAM file using the random routine is as follows:

Label	Operation	FT	Г	Operands & Remarks
11	n 10	37 33	Į.	75 a0 45 50 55 a0 as 80
S.T.A.R.T				
			Γ	
	L.L.A.F		T	LS.R.D.O O.P.E.N. J.S.A.M. RANDOM
	DC.	-	r	DELAD
	1.0.		T	DELAD
	2 N.		t	ERROR . GO TO FRROR ROUTINE JE NEG
			t	
	LIAF		t	I.S.B.D
	D (+	t	DELAD . DEL ADDRESS (DCOULPED)
┝┅╍╺╺	1.0	-	t	
		-+-	t	SPROP GALLAND SPROP COULTING IS NEC
+	Div.	++-	t	ERISION L. L. LIND I HOLDERRICH RUCH LIND DI LINE ILIPI PIERA
	1.1.05	+	┝	
	L.L.B.F	-	┢	LISKILL I III ILLISEL IISAN KARDON IIIIIIII
		++	ł	DIFLIAD
	L.P.	H	⊢	DIFILADITIT.
		H	┞	ERKOR
<u> </u>	1 1 1 1 A	H	Ł	
* D.E.I	TABLE	F	P	R. I.S.A.M. R.A.M.DIO.M. R.O.UT.I.M.E.
DFIAD	D.S.A.	-	┢	FILLEA, I.I., U.S.E.R. FILLE I.I.I.I.I.I.I.I.I.I.I.I.I.I.I.I.I.I.I.
	-		⊢	<u></u>
ERROR	EGU		┡	F ERROR ROUTINE
ha shul			Ł	<u></u>
	E.N.D.	\vdash		START I I I I I I I I I I I I I I I I I I I
		4-	L	<u></u>
		1	1	<u> </u>
Lul	L L.		L	<u>La constante da la constante da constante</u>

Disk File Information (DFI) Table. The ISAM random routine requires a DFI table describing the file. The DFI table (which is 20 words long, numbered 0-19) has eighteen entries, eight of which must be filled in by the user. The remaining ten entries must be initialized to zero by the user and are filled in by the program during execution.

Figure 17 shows the DFI table for the ISAM random routine.

<u>Operation of the ISAM Random Routine</u>. When the routine is entered at the open entry point ISRDO, it checks the validity of the DFI table entries, sets pointers and switches to be used internally by the routine and sets the return code in the DFI table to the code for file open.

The file is then entered at ISRD to process a record.

When the routine is entered at ISRDC, the return code is set to file closed.

The ISAM random routine returns to the statement immediately following the parameter that follows the LIBF to the routine for any of the three entry points.

RPG OBJECT TIME SUBROUTINES

Included in the DM2 System Library is a group of subroutines that performs functions for the RPG object program. These subroutines are intended for system use only. Brief descriptions of the subroutines and their entry points are listed below.

RPG Decimal Arithmetic

Add, Subtract, and Numeric Compare. This subroutine performs the addition, subtraction and numeric comparison functions requested in the RPG calculation specification.

The entry points are:

RGADD - decimal addition routine

RGSUB - decimal subtraction routine

RGNCP - decimal numeric compare

<u>Multiply</u>. This subroutine multiplies two decimal fields defined in an RPG program.

The entry point is:

RGMLT - decimal multiply

<u>Divide</u>. The RPG object program calls this subroutine to divide one decimal field by another and store the quotient in a third field.

The entry point is:

RGDIV - decimal divide

<u>Move Remainder</u>. This subroutine is called by the RPG object program immediately following a divide operation. It places the remainder in a specified field.

The entry point is:

RGMVR - move remainder

Word	Entry	Meaning
0,1,2	DSA ,	The first entry in a DFI table is always a DSA statement. The DSA statement allows the programmer to refer symbolically to a disk stored data file without knowing its actual location. The label is defined as the current value of the Location Assignment Counter when the DSA statement is encountered. The operand is the name of the data file. Further information on DSA may be found in <u>IBM 1130/1800</u> <u>Assembler Language</u> .
3	DC /00XX	XX equals the key length in characters. Maximum is /0032. (50 characters). The key length must be the same as the key length in the file being accessed.
4	DC /0XXX	XXX equals the length of the record in words. The maximum entry is $/0140$ (one 320-word record). The record length must be the same as the record length on the file being accessed.
5	DC LABEL 1	The address of the index buffer. This address must be on an even-word boundary. The length of the index buffer is calculated by multiplying the number of index entries per sector by the index entry length and adding 2. The maximum length of this buffer is 322 words.
6	DC LABEL 2	The address of the data buffer. This address must be on an even-word boundary. The length of the data buffer is calculated by multiplying the number of records per sector by the record length and adding two. The maximum length of the data buffer is 322 words.
7	DC /XXXX	Routine type code. For ISAM random retrieve, XXXX = 0100. For ISAM random update, XXXX = 1000.
8	DC LABEL 3	Address of the key hold area containing the key of the record to be processed. The key must be in a special format: one character per word, occupying the rightmost eight bits of the word. See the coding for ISAM sequential that includes a LIBF to ISETL for an example of how the key can be placed in this format.
9	DC /0000	Return code. This entry must be reserved by the user. After each LIBF to any of the three entry points in the ISAM random routine it should be checked for the return code.
10	DC /0000	Address of record in process. This word must be reserved by the user.
1 Re	turn codes	for ISAM random are as follows:
Hex	adecimal	
Nun	ber Meaning	
5555	5 File is ope	en Iom retrieve or undate function
8051	Index buff	er not on even-word boundary
8052	2 Data buffe	r not on even-word boundary
8053	B Key lengt	h in DFI table not same as key length in label
8054	Record les	ngth in DFI table not same as record length in label
805	o rile is not Write befo	open ore read on update
8057	7 Record no	t on file
OFFI	F File is clo	sed
All 8 to retry	XXX return code the operation.	s except 8056 and 8057 are terminal errors. The file must be reopened to allow the program Processing will again start at the first record.

.

Figure 17. Disk File Information Table for ISAM Random (Part 1 of 2)

Word	Entry	Meaning
11	DC /0000	Address of the index entry used to locate the record. This word must be reserved by the user.
12	DC /0000	Number of index entries per sector. This word must be reserved by the user.
13	DC /0000	Index entry length in words. This word must be reserved by the user.
14	DC /0000	Number of records per sector. This word must be reserved by the user.
15	DC /0000	Prime data record number. This word must be reserved by the user.
16	DC /0000	Number of index entry in process. This word must be reserved by the user.
17	DC /0000	First-time switch. This switch is set off after one record has been processed.
18	DC /0000	Internal switch used to indicate the record found is in the overflow area. This word must be reserved by the user.
19	DC /XXXX	Read/Write indicator. If routine type code (word 7 of this table) was a retrieve, this word should be set to /0000. If word 7 indicates an update file, this word should be /0000 for the retrieve and /0001 for the update. This word should be reset to /0000 before the next retrieve.

Figure 17. Disk File Information Table for ISAM Random (Part 2 of 2)

<u>Binary/Decimal Conversion</u>. This subroutine converts a three-word binary number to a fourteen-digit decimal number and vice-versa.

The entry points are:

RGBTD - binary to decimal conversion

RGDTB - decimal to binary conversion

RPG Sterling and Edit

<u>Sterling Input Conversion</u>. This subroutine converts a field in the British sterling format of pounds, shillings, pence and decimal pence to a decimal format of pence and decimal pence.

The entry point is:

RGSTI - sterling input conversion

<u>Sterling Output Conversion</u>. This subroutine performs the reverse function of RGSTI.

The entry point is:

RGSTO - sterling output conversion

<u>Edit</u>. This subroutine edits a numeric field using a user-specified edit word or edit code and places the edited value in an output area.

The entry point is:

RGEDT - edit

RPG Move

The five subroutines that comprise this group are responsible for the movement of data and fields requested by the object program.

The entry points and functions of these subroutines are:

RGMV1, RGMV5 -	move data from I/O buffer to assigned core field
RG MV2 -	move data from assigned core field to I/O buffer
RGMV3 -	perform the RPG calc operation MOVE
RG MV4 -	perform the RPG calc operation MOVEL

RPG Compare

This subroutine is used to compare alphameric fields.

The entry point is:

RGCMP - alphameric compare

RPG Indicators

<u>Test</u>. The condition of indicators specified in columns 9-17 of an RPG calculation specification are tested. If the conditions are met, the calc operation is performed. If the conditions are not met the operation is skipped.

The entry point is:

RGSI1 - test indicators

Set Resulting Indicators On Conditionally. This subroutine sets on resulting indicators as required based on the results of an arithmetic operation, a compare operation, or a table lookup. The resulting indicators are specified in columns 54-59 of the calculation specification.

The entry point is:

RGSI2 - set resulting indicators conditionally

<u>Set Resulting Indicators On or Off</u>. This subroutine will set or reset from one to three resulting indicators.

The entry points are:

RGSI3 - set resulting indicators on unconditionally

RGSI4 - clear resulting indicators off unconditionally Zero or Blank Test. This subroutine tests for a zero or blank and returns an indication to the requesting program.

The entry point is:

RGSI5 - test a field for zero or blank

<u>RPG Miscellaneous</u>

<u>Test Zone</u>. Tests the zone of the leftmost position of an RPG alpha field and returns an indication to the requesting program.

The entry point is:

RGTSZ - perform TESTZ operation

<u>Convert Record ID</u>. Converts the record ID number supplied on a Record Address File (RAF) to a two-word binary number.

The entry point is:

RGCVB - convert record ID number to binary

Object-Time Error. Load accumulator with error number supplied by user and wait at \$PRET for operator action. This subroutine then interprets operator action and proceeds accordingly.

The entry point is:

RGERR - RPG object program error interface

<u>Blank After</u>. This subroutine performs the RPG blank-after function if specified on the RPG output specification.

The entry point is:

RGBLK - zero or blank a field

Subroutines Used by FORTRAN (C/PT System)

Many of the functions and capabilities available within the general I/O and conversion subroutines described in this manual are beyond specification by the FORTRAN language. For example, the feed function of the 1442 cannot be specified in FORTRAN. Therefore, a set of limited-function I/O and conversion subroutines is included in the subroutine library for use by FORTRAN-compiled programs. Any subroutines written in Assembler language that execute 1/0 operations, and that are intended to be used in conjunction with FORTRAN-compiled programs must employ these special I/O subroutines for any I/O device specified in a mainline *IOCS record or for any device on the same interrupt level.

These subroutines are intended to operate in an error-free environment and thus provide no preoperative parameter checking.

The subroutine library contains the following special routines:

- CARDZ 1442 I/O Subroutine
- TYPEZ Keyboard/Console Printer I/O Subroutine
- WRTYZ Console Printer Subroutine
- PRNTZ 1132 Printer Subroutine
- PAPTZ 1134/1055 Paper Tape I/O Subroutine
- PLOTX 1627 Plotter Subroutine (see PLOTX)
- HOLEZ IEM Card Code/EBCDIC Conversion Subroutine
- EBCTB EBCDIC/Console Printer Code Table
- HOLTB IBM Card Code Table
- GETAD Subroutine Used to Locate Start Address of EBCTB/HOLTB

GENERAL SPECIFICATIONS

Except for PLOTX, the FORTRAN I/O device subroutines operate in a nonoverlapped mode. Thus, the device subroutines do not return control to the calling program until the operation is completed. These subroutines are all LIBF's without parameters. The input/output buffer for the subroutines is a 121-word buffer starting at location /003C. The maximum amount of data transferable is listed in the description of each subroutine. Output data must be stored in unpacked (one character per word) EBCDIC format, /00XX. Data entered from an input device is converted to unpacked (one character per word) EBCDIC format, /00XX.

The EBCDIC character set recognized by the subroutine comprises digits 0-9, alphabetic characters A-Z, blank, and special characters \$-+. &= (), '/*<%#@. Any other character is recognized as a blank by all subroutines except HOLEZ. HOLEZ recognizes an invalid character as an asterisk.

The Accumulator, Extension, and Index Registers 1 and 2 are used by the FORTRAN device subroutines and must be saved, if required, before entry into any given FORTRAN subroutine.

The Accumulator must be set to zero for input operations. For output operations, the Accumulator must be set /0002, except for PRNT2 and WRTY2, in which output is the only valid operation. Index Registers 1 and 2 are set to the number of characters transmitted, except for PRNT2 (1132 Printer) in which Index Register 2 contains the number of characters printed plus an additional character for forms control.

ERROR HANDLING

Device errors, e.g., not-ready and read check, cause a WAIT in the subroutine itself. After the appropriate corrective action is taken by the operator, PROGRAM START is pressed to execute or reinitiate the operation.

DESCRIPTIONS OF I/O SUBROUTINES

The subroutines described in the sections that follow do not provide a check to determine validity of parameters (contents of Accumulator and Index Register 2). Invalid parameters cause indeterminate operation of the subroutines. TYPEZ - KEYBOARD/CONSOLE PRINTER I/O SUBROUTINE

<u>Buffer Size</u>. Maximum of 80 words input, 120 words output.

<u>Keyboard Input</u>. The subroutine returns the carrier, reads up to 80 characters from the Keyboard, and stores them in the I/O buffer in EBCDIC format. Upon recognition of the end-of-field character or reception of the 80th character, the subroutine returns control to the user (the remainder of the buffer is unchanged). Upon recognition of the erase field character or the backspace character, the carrier is returned and the subroutine is reinitialized for the reentry of the entire message. Characters are printed by the Console Printer during Keyboard input.

<u>Console Printer Output</u>. The subroutine returns the carrier and prints the number of characters indicated by Index Register 2 from the I/O buffer.

<u>Subroutines Required</u>. The following subroutines are required with TYPEZ:

HOLEZ, GETAD, EBCTB, HOLTB

WRTYZ - CONSOLE PRINTER OUTPUT SUBROUTINE

Buffer Size. Maximum of 120 words.

Operation. This subroutine returns the carrier and prints the number of characters indicated by Index Register 2 from the I/O buffer.

<u>Subroutines Required</u>. The following subroutines are required with WRTYZ:

GETAD, EBCTB

CARDZ - 1442 CARD READ PUNCH 1/O SUBROUTINE

Buffer Size. Maximum of 80 words.

<u>Card Input</u>. This subroutine reads 80 columns from a card and stores the information in the I/O buffer in EBCDIC format.

<u>Card Output</u>. This subroutine punches the number of characters indicated by Index Register 2 from the I/O buffer. Punching is done in IBM card code format. <u>Subroutines Required</u>. The following subroutines are required with CARDZ:

HOLEZ, GETAD, EBCTB, HOLTB

PAPTZ - 1134/1055 PAPER TAPE READER PUNCH I/O SUBROUTI NE

Buffer Size. Maximum of 80 characters.

<u>1134 Paper Tape Input</u>. This subroutine reads paper tape punched in PTTC/8 format. Paper tape is read until 80 characters have been stored or until a new-line character is read. If 80 characters have been stored and a new-line character has not been read, one more character, assumed to be a new-line character, is read from tape. (Delete and case-shift characters cause nothing to be stored.) If the first character read is not a case-shift character, it is assumed to be a lower case character. The input is converted to EBCDIC format.

1055 Paper Tape Output. The contents of the I/O buffer is converted from EBCDIC to PTTC/8, and the number of characters indicated by Index Register 2 is punched, in addition to the required case-shift characters.

PRNTZ - 1132 PRINTER OUTPUT SUBROUTINE

Buffer Size. Maximum of 121 characters.

<u>Index Register 2</u>. The value stored in Index Register 2 must be the number of characters to be printed plus an additional character for carriage control. Up to 120 characters can be printed in any one operation. The first character to be printed is stored in location /003D.

The carriage of the 1132 printer is controlled prior to the printing of a line. The following is a list of the carriage control characters and their related functions:

/00F1 Skip to channel 1 prior to
printing
/00F0 Double space prior to printing
/004E No skip or space prior to printing
Any other character - Single space prior
to printing.

Channel 12 Control. If a punch in channel 12 is encountered while a line is being printed, a skip-to-channel-1 is taken prior to the printing of the next line.

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Subroutines Used by FORTRAN (DM2 System)

Many of the I/O and conversion subroutines cannot be specified in FORTRAN. Therefore, the System Library includes a set of limited-function I/O and conversion subroutines for FORTRAN programs. Any Assembler language I/O subroutines used by FORTRAN programs must employ these special subroutines for any I/O device specified in a mainline *IOCS control record.

Of all the FORTRAN device subroutines, only DISKZ, PRNTZ, PRNZ, and PLOTX return control to the caller after initiating an operation (PLOTX is described with the basic ISSS).

These subroutines are intended for use in an error-free environment and thus provide no preoperative parameter checking.

The System Library contains the following ISS and conversion subroutines for FORTRAN programs:

CARDZ	-	1442 I/O Subroutine
PNCHZ	-	1442 Output Subroutine
REACZ	-	2501 Input Subroutine
TYPE Z	-	Keyboard/Console Printer 1/0
		Subroutine
WRTYZ	-	Console Printer Subroutine
PRNTZ	-	1132 Printer Subroutine
PRNZ	-	1403 Printer Subroutine
PAPTZ	-	1134/1055 Paper Tape I/O
		Subroutine
PLOTX	-	1627 Plotter Subroutine
DISKZ	-	Disk I/O Subroutine
HOL EZ	-	IBM Card Code/EBCDIC Conversion
		Subroutine
EBCTB		EBCDIC/Console Printer Code
		Table
HOLTB	-	IBM Card Code Table

GETAD - Subroutine to Locate Start Address of EBCTB/HOLTB

GENERAL SPECIFICATIONS (EXCEPT DISKZ)

The "Z" device subroutines are ISS subroutines. These subroutines are all LIBF's without parameters. They use a 121-word input/output buffer, contained in the nondisk FORTRAN I/O subroutine SFIO. The maximum amount of data transferable is listed in the description of each subroutine. Output data must be stored in unpacked right-justified (one character per word) EBCDIC format. Input data is converted to unpacked EBCDIC format. The EBCDIC character set recognized by the subroutines comprises digits 0-9, alphabetic characters A-Z, blank, and special characters \$-+. &= (), '/*<\$#0. Any other character is recognized as a blank by all subroutines except HOLEZ. HOLEZ recognizes an invalid character as an asterisk.

If a "2" subroutine is used by an Assembler language I/O subroutine, the user should be aware of the significant information carried by the different registers. The Accumulator, Extension, and Index Registers 1 and 2 are used by the FORTRAN device subroutines and must be saved, if required, before entry into the subroutines. The Accumulator must be set to zero for input operations.

For output operations, the Accumulator must be set to /0002, except for PRNZ, PRNIZ, PNCHZ, and WRTYZ, in which output is the only valid operation. Index Register 2 must be set to the number of characters to be transferred, except for PRNZ and PRNTZ. For these two subroutines, Index Register 2 must contain the number of characters to be printed plus an additional character for carriage control. Index Register 1 must contain the starting address of the input buffer.

ERROR HANDLING

Device errors, e.g., not ready and read check, result in a branch to \$PST1, \$PST2, \$PST3, and \$PST4 depending on the level to which the device is assigned. After the appropriate corrective action is taken by the operator, PROGRAM START is pressed to execute or reinitiate the operation.

If a monitor control record is encountered by CARDZ, REAEZ, or PAPTZ, the subroutine initiates a CALL EXIT. The control record itself will not be processed.

<u>DESCRIPTIONS OF I/O SUBROUTINES</u>

The subroutines described in the sections that follow do not provide a check to determine validity of parameters (contents of Accumulator and Index Register 2). Invalid parameters cause indeterminate operation of the subroutines.

<u>TYPEZ - KEYBOARD/CONSOLE PRINTER 1/0</u> <u>SUBROUTINE</u>

<u>Buffer Size</u>. Maximum of 80 words input, 120 words output.

<u>Keyboard Input</u>. The subroutine returns the carrier and reads up to 80 characters from the Keyboard and stores them in the I/O buffer in EBCDIC format. Upon recognition of the end-of-field character or reception of the 80th character, the subroutine returns control to the user (the remainder of the buffer is unchanged). Upon recognition of the erase field character or the backspace character, the carrier is returned and the subroutine is reinitialized for the reentry of the entire message. Characters are printed by the Console Printer during Keyboard input.

<u>Console Printer Output</u>. The subroutine returns the carrier and prints the number of characters indicated by Index Register 2 from the I/O buffer.

<u>Subroutines Required</u>. The following subroutines are required with TYPEZ:

HOLEZ, GETAD, EBCTB, HOLTB

WRTYZ - CONSOLE PRINTER OUTPUT SUBROUTINE

Buffer Size. Maximum of 120 words.

<u>Operation</u>. This subroutine returns the carrier and prints the number of characters indicated by Index Register 2 from the I/O buffer.

<u>Subroutines Required</u>. The following subroutines are required with WRTY2:

GETAD, EBCTB

CARDZ - 1442 CARD READ PUNCH I/O SUBROUTINE

Buffer Size. Maximum of 80 words.

<u>Card Input</u>. This subroutine reads 80 columns from a card and stores the information in the I/O buffer in EBCDIC format.

<u>Card Output</u>. This subroutine punches the number of characters indicated by Index Register 2 from the I/O buffer. Punching is done in IEM Card Code.

<u>Subroutines Required</u>. The following subroutines are required with CARD2:

HOLEZ, GETAD, EBCTB, HOLTB

PAPTZ - 1134/1055 PAPER TAPE READER PUNCH 1/0 SUBROUTINE

Buffer Size. Maximum of 120 characters.

1134 Paper Tape Input. This subroutine reads paper tape punched in PTTC/8 format. The subroutine reads paper tape until 120 characters have been stored or until a new-line character is read. If 120 characters have been stored and a new-line character has not been read, one more character, assumed to be a new-line character, is read from tape. (Lelete and case-shift characters cause nothing to be If the first character read is stored.) not a case-shift character, it is assumed to be a lower case character. Subsequent reads assume the same case as the last character read until the case is changed by another case-shift character. The input is converted to EBCDIC format.

1055 Paper Tape Output. The contents of the I/O buffer is converted from FBCDIC to PTTC/8, and the number of characters indicated by Index Register 2 is punched, in addition to the required case-shift characters.

PRNTZ - 1132 PRINTER OUTPUT SUBROUTINE

Buffer Size. Maximum of 121 characters.

<u>Index Register 2</u>. The value stored in Index Register 2 must be the number of characters to be printed, plus an additional character for carriage control. Up to 120 characters can be printed in any one operation. If PRNTZ is user-called by a LIBF PRNTZ, only an even number of characters are printed. To print an odd number of characters add one additional blank.

The carriage of the 1132 Printer is controlled prior to the printing of a line. The following is a list of the carriage control characters and their related functions:

/00F1 Skip to channel 1 prior to
printing
/00F0 Double space prior to printing
/004E No skip or space prior to printing
Any other character - Single space prior
to printing.

<u>Channel 12 Control</u>. If a punch in channel 12 is encountered while a line is being printed, a skip-to-channel-1 is taken prior to the printing of the next line provided the next function is not /004E (no skip or space prior to printing).

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PNCHZ - 1442 OUTPUT SUBROUTINE

Buffer Size. Maximum of 80 words.

<u>Card Output</u>. This subroutine punches from the I/O buffer the number of characters indicated in the location preceding the buffer. Punching is done in IBM Card Code.

<u>Subroutines Required</u>. The following subroutines are required with PNCH2:

HOLEZ, GETAD, EBCTB, HOLTB

<u>READZ - 2501 INPUT SUBROUTINE</u>

Buffer Size. Maximum of 80 words.

<u>Card input</u>. This subroutine reads 80 columns from a card and stores the information in the I/O buffer in EBCDIC format.

<u>Subroutines Required</u>. The following subroutines are required with READZ:

HOLEZ, GETAD, EBCTB, HOLTB

PRNZ - 1403 PRINTER SUBROUTINE

Buffer Size. Must be 121 words.

Index Register 2. The first character in the I/O buffer is the carriage control character, followed by up to 120 characters to be printed. If less than 120 characters are to be printed, the remainder of the buffer must be cleared to blanks before PRNZ is called. A value of 1 in Index Register 2 indicates that the I/O buffer contains only a carriage control character. A value of greater than 1 in Index Register 2 indicates that a line is to be printed.

The carriage is controlled prior to the printing of a line; no "after-print" carriage control is performed. The following is a list of the carriage control characters and their related functions:

/00F1 Skip to channel 1 prior to
printing
/00F0 Double space prior to printing
/004E No skip or space prior to printing
Any other character - Single space prior
to printing.

<u>Channel 12 Control</u>. If a punch in channel 12 is encountered while a line is being printed, a skip to channel 1 is executed prior to printing the next line provided the next function is not /004E (no skip or space prior to printing).

Data Code Conversion Subroutines

The basic unit of information within the 1130 computing system is the 16-bit binary word. This information can be interpreted in a variety of ways, depending on the circumstances. For example, in internal computer operations, words may be interpreted as instructions, as addresses, as binary integers, or as real (floating point) numbers (see "Arithmetic and Functional Subroutines").

A variety of data codes exists for the following reasons:

- 1. The programmer needs a compact notation to represent externally the bit configuration of each computer word. This representation is provided in the hexadecimal notation.
- A code is required for representing alphameric (mixed alphabetic and numeric) data within the computer. This code is provided by the Extended Binary Coded Decimal Interchange Code (EBCDIC).
- 3. The design and operation of the input/output devices is such that many of them impose a unique correspondence between character representations in the external medium and the associated bit configurations within the computer. Subroutines are needed to convert input data from these devices to a form on which the computer can operate and to prepare computed results for output on the devices.

This and following sections of the manual describe the data codes used and the subroutines provided for converting data representations among these codes.

A detailed description of the binary, hexadecimal, and decimal number systems is contained in the publication, <u>IBM Number</u> <u>Systems</u>, FC20-1618.

Descriptions of Data Codes

In addition to the internal 16-bit binary representation, the conversion subroutines handle the following codes:

- Hexadecimal Notation.
- IBM Card Code.

- Perforated Tape and Transmission Code (PTTC/8).
- Console Printer (1053) Code.
- 1403 Printer Code (DM2 System only).
- Extended Binary Coded Decimal Interchange Code (EBCDIC).

A list of these codes is contained in Appendix D.

HEXADECIMAL NOTATION

Although binary numbers facilitate the operations of computers, they are awkward for the programmer to handle. A long string of 1's and 0's cannot be effectively transmitted from one individual to another. For this reason, the hexadecimal number system is often used as a shorthand method of communicating binary numbers. Because of the simple relationship of hexadecimal to binary, numbers can easily be converted from one system to another.

In hexadecimal notation a single digit is used to represent a 4-bit binary value as shown in Figure 18. Thus, a 16-bit word in the 1130 System can be expressed as four hexadecimal digits. For example, the binary value

1101001110111011

can be separated into four sections as follows:

Binary 1101 0011 1011 1011 Hexadecimal D 3 B B

Another advantage of hexadecimal notation is that fewer positions are required for output data printed, punched in cards, or punched in paper tape. In the example above, only four card columns are required to represent a 16-bit binary word.

BINARY	DECI MAL	HEXADECIMAL	
0000	0	0	
0001	1	1	
00 10	2	2	
0011	- 3	3	
0 100	4	4	
0101	5	5	
0 1 10	6	6	
0111	7	. 7	
1000	8	8	
1001	9	9	
10 10	10	A	
10 1 1	11	B	
1100	12	с	
1101	13	D	
1110	14	Е	
1111	15	F	

Figure 18. Hexadecimal Notation

IBM CARD CODE

IBM Card Code can be used as an input/output code with the 1442 Card Read Punch, 1442 Card Punch, and 2501 Card Reader, and as an input code on the Keyboard.

This code defines a character by a combination of punches in a card column. Card code data is taken from or placed into the leftmost twelve bits of a computer word as shown below:

 Card Row
 12
 11
 0
 1
 2
 3
 4
 5
 6
 7
 8
 9

For example, a plus sign, which has a card code of 12, 6, 8, is placed into core storage in the binary configuration illustrated in the following diagram:



PERFORATED TAPE AND TRANSMISSION CODE (PTTC/8)

The PTTC/8 code is an 8-bit code used with IBM 1134/1055 Paper Tape units. This code represents a character by a stop position, a check position, and six positions representing the 6-bit code, BA8421. PTTC/8 characters can be packed two per computer word as shown below:

				1	st A							2r	nd .				
PTTC/8 Characters	S	B	A	с	8	4	2	<u>ر</u>	s	B	A	С	8	4	2		
Computer Word	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	l

The graphic character is defined by a combination of binary code and case; a control character is defined by a binary code and has the same meaning in upper or lower case. This implies that upper and lower case characters must appear in a PTTC/8 message wherever necessary to establish or change the case.

The binary and PTTC/8 codes for 1/(lower case) and =? (upper case) are shown in Figure 19.

The delete and stop characters have a special meaning (in check mode only) when encountered by the paper tape subroutines.



Figure 19. PTTC/8 Code for the Characters 1/ (if lower case) or the Characters =? (if upper case)

CONSOLE PRINTER CODE

The Console Printer uses an 8-bit code that can be packed two characters per 16-bit word.

The following control characters have special meanings when used with the Console Printer.

Character Control Operation

HT	Tabulate
RES	Shift to black ribbon
NL	Carrier return to new line
BS	Backspace
LF	Line feed without carrier
	return
RS	Shift to red ribbon

EXTENDED BINARY CODED DECIMAL INTERCHANGE CODE (EBCDIC)

EBCDIC is the standard code for internal representation of alphameric and special characters and for the 1132 Printer. This code uses eight binary bits for each character, thus making it possible to store either one or two characters in a 16-bit word. Combinations of the eight bits allow 256 possible codes. (At present, not all of these combinations represent characters.) The complete EBCDIC code is listed in Appendix D.

For reasons of efficiency, most of the conversion subroutines do not recognize all 256 codes. The asterisked codes in Appendix D constitute the subset recognized by most of the conversion subroutines.

1403 PRINTER CODE

The 1403 Printer uses a 6-bit binary code with one parity bit. Data format is two 7-bit characters per word, as follows:

Bit	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	
Value	*	P	32	16	8	4	2	1	*	P	32	16	8	4	2	1	
	Ìsi	da	ita' c	han	acte	er			2nd data character								
* = No P = Par	t U ity	sed Bit			•												

Parity bits are not assigned by the hardware. The conversion subroutine must assign the parity bits and arrange the characters in the form in which they are to be printed.

Conversion Subroutines

These subroutines convert data to and from 16-bit binary words and I/O device codes.

- BINDC Binary value to IBM Card Code decimal value.
- DCBIN IBM Card Code decimal value to binary value.

BINHX Binary value to IBM Card Code hexadecimal value.

- HXBIN IBM Card Code hexadecimal value to binary value.
- HOLEB IBM Card Code subset to EBCDIC subset; EBCDIC subset to IBM Card Code subset.
- SPEED IBM Card Code characters to EBCDIC; EBCDIC to IBM Card Code characters.
- PAPEB PTTC/8 subset to EBCDIC subset; EBCDIC subset to PTTC/8 subset.
- PAPHL PTTC/8 subset to IBM Card Code subset; IBM Card Code subset to PTTC/8 subset.
- PAPPR PTTC/8 subset to Console Printer or 1403 Printer code.
- HOLPR IBM Card Code subset to Console Printer or 1403 Printer code.
- EBPRT EBCDIC subset to Console Printer or 1403 Printer code.

The following conversion tables are used by some of the conversion subroutines.

- PRTY Console Printer and 1403 Printer code.
- EBPA EBCDIC and PTTC/8 subsets.
- HOLL IBM Card Code subset.

The following conversion subroutines are used by the DM2 system only.

- BIDEC 32-bit binary value to IBM Card Code decimal value.
- DECBI IBM Card Code decimal value to 32-bit binary value.
- ZIPCO Supplement to all standard conversions except those involving PTTC/8.

The first four listed subroutines and the DM2 subroutines BIDEC and DECBI change numeric data from its input form to a binary form, or from a binary form to an appropriate output data code. The last eight (including ZIPCO) convert entire messages, one character at a time, from one input/output code to another. The types of conversions accomplished by these subroutines are illustrated in Figure 20.

Except where specified, these subroutines do not alter the Accumulator, Extension, Carry and Overflow indicators, or any index register.

					с	ONVERTED TO	5		•	
CONVERTED FROM	Binary	IBM Card Code (256)	IBM Card Code (Subset)	PTTC/8 (Subset)	EBCDIC (256)	EBCDIC (Subset) 1132 Printer	Console Printer	Hex Equivalent (Card Code)	Decimal Equivalent (Card Code)	1403 Printer Code
Binary								BINHX	BINDC BIDEC	
IBM Card Code (256)					SPEED ZIPCO*		ZIPCO*			ZIPCO*
IBM Card Code (Subset)				PAPHL		HOLEB	HOLPR			HOLPR
PTTC/8 (Subset)		•	PAPHL			PAPEB	PAPPR			PAPPR
EBCDIC (256)		SPEED					ZIPCO*			ZIPCO*
EBCDIC (Subset) 1132 Printer			HOLEB	PAPEB			EBPRT			EBPRT
Hex Equivalent (Card Code)	HXBIN	•		PAPHL		HOLEB	HOLPR			HOLPR
Decimal Equivalent (Card Code)	DCBIN DECBI			PAPHL		HOLEB	HOLPR			HOLPR
1403 Printer Code		ZIPCO*			ZIPCO*		ZIPCO*			
Console Printer Code		ZIPCO*			ZIPCO*					ZIPCO*

* In conjunction with appropriate conversion table.

Figure 20. Types of Conversion

- <u>Note 1</u>. All mention of 1403 Printer Code applies to the DM2 system only.
- <u>Note 2</u>. The conversion subroutines and conversion tables for the Communications Adapter are described in the publication <u>IBM</u> <u>1130 Synchronous Communications</u> <u>Adapter Subroutines</u>. The subroutines are EBC48, HOL48, and HXCV. The adapter subroutine conversion table is STRTB.

Error Checking

All code conversion subroutines (except SPEED and ZIPCO) accept <u>only</u> the codes marked with an asterisk in Appendix D. An input character that does not conform to a specified code is an error.

BINHX and BINDC subroutines do not detect errors. HXBIN and DCBIN terminate conversion at the point of error detection; they do not replace the character in error. The contents of the Accumulator are meaningless when conversion is terminated because of an error. The remainder of the conversion subroutines replace the character in error with a space character, stored in the output area in output code. Conversion is not terminated when an error is detected.

When a conversion subroutine detects an error it turns the Carry indicator <u>off</u> and turns the Overflow indicator <u>on</u> before returning control to the user. Otherwise, the settings of the Carry and Overflow indicators are not changed by the conversion subroutines.

BINDC

This subroutine converts a 16-bit binary value to its decimal equivalent in five IBM Card Code numeric characters and one sign character. The five characters and the sign are placed in six computer words as illustrated below.

I/O Locations	Conversion Data	Bits in Core Storage														15	
Accumulator	+01538	0	0	0	0	0	1	1	0	0	0	0	0	0	0	1	0
OUTPT	+	1	0	0	0	0	0	0	Ò	Ì	0	1	0	G	0	D	, te
OUTPT + 1	0	0	0	1	0	0	0	0	0	0	0	0	0	0	Q	Ø	38
OUTPT + 2	1	0	0	0	1	0	0	0	0	0	0	0	0	đ.	0	Ø	×.
OUTPT + 3	5	0	0	0	0	0	0	0	1	0	0	0	0	Û.	0	0	Ø
OUTPT + 4	3	0	0	0	0	.0	1	0	0	0	0	0	0	Q	Ø	Ð	Ø
OUTPT + 5	8	0	0	0	0	0	0	. O	0	0	0	1	0	Ø	Ø	١ T	Ø

Calling Sequence

Label	Operation	F	T	Operands & Remarks				
21 25	27 30	32	33	35 40 45 50				
L	LIBE			DIC BIT N				
	DC.			INPUT				
	•							
I.N.P.U.T	BSS			6				
		ŀ						

Calling Sequence

Label Operation			F	Ŧ	Operands & Remarks
21 25	27	30	32	33	35 40 45 50
Indent 1	LLB	Ē			BIINDC
لنبيل	DC I			ŀ	$O_1U_1T_1P_1T_1$
	•				••••••••••••••••••••••••••••••••••••••
OUTPT	BSS			Ц	6

Input

Input is a 16-bit binary value in the Accumulator.

Output

Output is an IBM Card Code sign character (plus or minus) in location OUTPT, and five IBM Card Code numeric characters in OUTPT+1 through OUTPT+5.

Errors Detected

The BINDC subroutine does not detect errors.

DCBIN

This subroutine converts a decimal value in five IBM Card Code numeric characters and a sign character to a 16-bit binary word. The conversion is the opposite of the BINDC subroutine conversion. Input is an IBM Card Code sign character in location INPUT and five IBM Card Code decimal characters in INPUT+1 through INPUT+5.

Output

Input

Output is a 16-bit binary value displayed in the Accumulator.

Errors Detected

Any sign other than an IBM Card Code plus, ampersand, space, or minus, or any decimal digits other than a space or 0 through 9 is an error. Any converted value greater than +32767 or less than -32768 is an error.

BINHX

This subroutine converts a 16-bit binary word into hexadecimal notation in four IBM Card Code characters as illustrated below.

1/O Locations	Conversion Data	Bits in Core Storage 0				
Accumulator	A59E	1010 0101 1001 1110				
OUTPT	A	1001 0000 0000 0000				
OUTPT + 1	5	0000 0001 0000 0000				
OUTPT + 2	9	0000 0000 0001 00000				
OUTPT + 3	E	1000 0001 0000 0000				

Calling Sequence

Label	Operation		1	1		Operands & Kemarks
	LIBE					BINHX.
ور المراجع المراجع الم	DC.		_			OUTPT.
L. L. L. L.	· · · · · ·		Į			• • • • • • • • • • • • • • • • • • • •
OUT P.T	BISISI	ļ			_	4
المرافعة الأرام		Į.				: 4.

Input

Input is a 16-bit binary word in the Accumulator.

Output

Output is four IBM Card Code hexadecimal digits in location OUTPT through OUTPT+3.

Errors Detected

The BINHX subroutine does not detect errors.

Input

Input is four IBM Card Code hexadecimal digits in INPUT through INPUT+3.

Output

Output is a 16-bit binary word in the Accumulator.

Errors Detected

Any input character other than an IBM Card Code 0 through 9 or A through F is an error.

HOLEB

This subroutine converts IBM Card Code subset to the EBCDIC subset or converts the EBCDIC subset to IBM Card Code subset. Code conversion is illustrated below.

1/O Locations	Conversion Data	Bits in Core Storage
	<u>در</u> در	1101 0001 1110 0010 0101 0000 0000 0000
OUTPT + 1	S	0010 1000 0000 0000

HXBIN

This subroutine converts four IBM Card Code hexadecimal characters into one 16-bit binary word. The conversion is the opposite of the BINHX subroutine conversion illustrated above.

Calling Sequence

Linei	Cperation	F	T		Operands & Remarks
	L.I.B.F				HIX.B.I.N.
المسلمية الم	 D.C	 	L		LN.P.U.T.
а.н. нн		 		ļ.,	* 1 - 1 1 1 1 <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1</u> <u>1-</u>
I N P U T	 BSS			İ	4
1 II.	 			L	1. J

<u>Calling_Seguence</u>

Lube 1	Operation	1 1	1		Operand	s & Remarks		
		$\sim e$	r			×.	÷.	
	LIBF		HOLEB	L. CAL	LSU	S.S.E.T. C	ONYER	S.I.O.N
	DC.		1 dade .	L. CON	TROL	PARAM	ETER	
	D.C.		INPUT	I.N.P.	U.T. A	REA AD	DRESS	
11	pc .	11	OUTPUT	UUUT	PUT	REALA	DDRES	5
	0c		FLILL	AHID .	RACT	ER COU	ALL I	
[المتنا			والمتعادية المراجع				
a. a. a. a			+ 1		لللحيات			
المتناسس	44	· + + +.		يتعتب		ب الرب ب	بنبي	
L'N'B'N'L	BSS	-	9	I.N.P.	U.T. A.	REAL .	ببتبت	
		-+ +		ين المناحب			حبابت	
	╶┼┅╺╼┿┥		÷ 1-4-3-4-4-4-4-	· · · · · · · · · · · · · · · · · · ·			بنقيت	
· • • • • • •		-++-	ᢤ┻ᠴᠴᠴ	ب بيانين	. .			
U.T.P.T	855		factor and the second s	OUT	PUT	REA		لللمحد
		- 11.		ب ال حد				

where

e indicates the direction of conversion,

f is the number of characters to be converted,

g is the length of the input area. g must be equal to or greater than f if e is 0. If e is 1, g must be equal to f/2, or (f+1)/2 if f is odd.

h is the length of the output area. If e is 0, h must be equal to or greater than f/2, or (f+1)/2, if f is odd. If e is 1, h must be equal to or greater than f.

Control Parameter

The control parameter consists of four hexadecimal digits. Digits 1-3 are not used. The fourth digit specifies the direction of conversion:

0 - IBM Card Code to EBCDIC 1 - EBCDIC to IBM Card Code

Input

Input is either IBM Card Code or EBCDIC characters, (as specified by the control parameter) starting in location INPUT. EBCDIC characters must be packed two characters per binary word. IBM Card Code characters are stored one character to each binary word.

Output

Output is either IBM Card Code or EBCDIC characters starting in location OUTPT. Characters are packed as described above.

If the direction of the conversion is IEM Card Code input to EBCDIC output, the input area can overlap the output area if the address INPUT is equal to or greater than the address OUTPT. If the direction of the conversion is EBCDIC input to IEM Card Code output, the input area can overlap the output area if the address INPUT+n/2 is equal to or greater than the address OUTPT+n, where n is the character count specified. The subroutine starts processing at location INPUT.

Character Count

This number specifies the number of characters to be converted; it is not equal to the number of binary words used for the EBCDIC characters because those characters are packed two per binary word. If an odd count is specified for EBCDIC output, bits 8 through 15 of the last word in the output area are not altered.

Errors Detected

Any input character not asterisked in Appendix D is an error.

SPEED

This subroutine converts IBM Card Code to EBCDIC or EBCDIC to IBM Card Code. SPEED accepts all 256 characters defined in Appendix D.

If the input is IBM Card Code, the conversion time is much faster than that of HOLEB because a different conversion method is used when all 256 EBCDIC characters are accepted. If the SPEED subroutine is called before a card reading operation is completed, the SPEED subroutine synchronizes with a CARD subroutine read operation by checking bit 15 of the word to be processed before converting the word. If bit 15 is a one, the SPEED subroutine waits in a loop until the CARD0 or CARD1 subroutine sets the bit to a zero.

Note: SPEED should not be used with READO or READ1 since the 2501 subroutines do not prestore 1 bits in each word of the I/O area. Use HOLEB or ZIPCO for 2501 operations.

Calling Sequence

Lobel	Operation F		Operands & Remarks
		• •	
	LIBE	SPEED . CAL	L EBCOLC CONVERSION
1.1.1.1.1.1.1	0.0.	10.0.d.e	TROL PARAMETER
		INPUT INP	UT AREA ADDRESS
	0.0.	OUT.P.U.T. OUT	PUT AREA ADDRESS
		FL C.H.A	RACITER
			<u> </u>
	<u> </u>	<u> </u>	
SNPUT	B.S.S.	S	U.T. AREAL LINE LINE
	1		I I I HALAMA I LAND I I A A A I I
A		<u> </u>	
	<u> </u>		
DUT.P.T	B.S.S.	bi	PUT AREA
2	1		

where

d indicates whether the EBCDIC characters are packed or unpacked,

e indicates the direction of conversion,

f indicates the character count,

g is the length of the input area,

h is the length of the output area,

g and h are defined as follows:

IBM Card Code to packed EBCDIC

 $g \ge f$ h \ge f/2, or (f+1)/2, if f is odd.

IBM Card Code to unpacked EBCDIC

g≥f h≥f

Packed EBCDIC to IBM Card Code

 $g \ge f/2$, or (f+1)/2, if f is odd. $h \ge f$

Unpacked EBCDIC to IBM Card Code

g≥f h≥f

Control Parameter

This parameter consists of four hexadecimal digits. Digits 1 and 2 must be zero. The third digit indicates whether the EBCDIC code is packed or unpacked.

0 - Packed, two EBCDIC characters per binary word

1 - Unpacked, one EBCDIC character per binary word (left-justified)

The fourth digit indicates the direction of conversion:

0 - IBM Card Code to EBCDIC 1 - EBCDIC to IBM Card Code

Input

Input is either IBM Card Code or EBCDIC characters (as specified by the control parameter) starting in location INPUT. EBCDIC characters can be packed or unpacked. IBM Card Code characters are stored one character to each binary word.

Output

Output is EBCDIC or IBM Card Code characters starting in location OUTPT. EBCDIC characters can be packed or unpacked; IBM Card Code characters are not packed.

The input area should not overlap the output area because of restart problems that can result from card feed errors.

Character Count

This parameter specifies the number of EBCDIC or IBM Card Code characters to be converted. If the character count is odd and the output code is packed EBCDIC, bits 8 through 15 of the last word are unaltered.

Errors Detected

Any input character code not listed in Appendix D is an error. All IBM Card Code punch combinations, except multiple punches in rows 1-7, are legal.

PAPEB

This subroutine converts PTTC/8 subset to EBCDIC subset or EBCDIC subset to PTTC/8 subset. PAPEB conversion of FBCCIC to PTTC/8 with the initialize case option selected is illustrated below.

I/O Locations	Conversion Data	Bits in Core Storage				
INPUT	2L	1101	0001	1110	0010	
OUTPT +0 +1	UC J S DEL	0000 0011	1110 0010	0101 0111	0001 1111	

Calling Sequence

Lobel Operation		1.	Operands & Remarks
a	4	L.L	· · · · · · · · · · · · · · · · · · ·
	LIBE	LΤ	PAPER CALL BITCHA CONVERSION.
			V. & & d.e CONTROL PARAMETER
	DC.	Π	INPUT INPUT. AREA ADDRESS
	D.C.		OUTPUT OUTPUT AREA ADDRESS
		11	FLANDER CHARACITER COUNTLY
	<u> </u>		
	<u> </u>	11	
<u></u>			
INPUT	B.S.S.		S
		Ш	
	·		
OUTPT	B.S.S.	LΤ	h OUTPUT AREA
		Ш	<u>}</u>

where

d is the case initialization digit,

e indicates the direction of conversion,

f indicates the character count,

g is the length of the input area. g must be equal to or greater than f/2 or (f+1)/2, if f is odd.

h is the length of the output area. h must be equal to or greater than f/2 or (f+1)/2, if f is odd.

Control Parameter

This parameter consists of four hexadecimal digits. Digits 1 and 2 are not used. The third digit indicates whether or not the case is to be initialized before conversion begins:

0 - Initialize case

1 - Do not alter case

The fourth digit indicates the direction of conversion:

0 - PTTC/8 to EBCDIC

1 - EBCDIC to PTTC/8

Input

Input (either PTTC/8 or EBCDIC characters, as specified by the control parameter) starts in location INPUT. Characters are packed two per 16-bit computer word in both codes.

Output

Output is either EBCDIC or PTTC/8 characters starting in OUTPT. Characters in either code are in packed format. The subroutine starts processing at location INPUT.

If the output is in EBCDIC, overlap of the input and output areas is possible if the address INPUT is equal to or greater than the address OUTPT.

If the output is in PTTC/8, overlap of the input and output areas is not recommended because the number of output characters might be greater than the number of input characters.

Character Count

This parameter specifies the number of PTIC/8 or EBCDIC characters in the input area. The count must include case-shift characters even though they will not appear in the output. Because the input is packed, the character count will not be equal to the number of binary words in the input area. If an odd number of output characters is produced, bits 8-15 of the last word used in the output area are set to a space character if the output is EBCDIC, or to a delete character if the output is PTTC/8.

There is no danger of overflowing the output area if the number of words in a PTTC/8 output area is equal to the number of characters in the input area.

Errors Detected

Any input character that is not marked with an asterisk in Appendix D is an error.

Subroutine Operation

If the input is in PTTC/8 code, all control characters (except case-shift (LC or UC) characters) are converted to output. Case-shift characters only define the case mode of the graphic characters that follow.

If the initialize option is selected, the case is set to lower. All characters are interpreted as lower case characters until an upper case shift (UC) character is encountered. If the do-not-alter option is selected, the case remains set according to the last case-shift character encountered in the previous LIBF message.

If the input is in EBCDIC, all data and control characters are converted to output. The user should <u>not</u> specify case shifting in his input message; this is handled automatically by the PAPEB subroutine.

Case-shift characters are inserted in a PTTC/8 output message where needed to define certain graphic characters that have the same binary value and are differentiated only by a case-mode character. For example, the binary value 0101 1011 (5B), is interpreted as a \$ in lower case and an ! in upper case (see Appendix D).

If the initialize option is selected, the case-shift character needed to interpret the first graphic character is inserted in the output message and the case mode is initialized for that mode. If the do-not-alter option is selected, the case mode remains set according to the last case-shift character required in the previous LIBF message, i.e., no case shift is forced.

If a case-shift character appears in the input message, it is output but does not affect the case mode. If it is an upper case shift (UC) and the next input character requires an upper case shift, the subroutine still inserts an upper case shift into the message, i.e., two UC characters will appear in the output message.

The conversion is halted when the character count is decremented to zero or when a new-line (NL) control character is read.

PAPHL

This subroutine converts PTTC/8 subset to IBM Card Code subset or IBM Card Code subset to PTTC/8 subset. The relationship of the two codes for converting PTTC/8 to IBM Card Code is illustrated below:

I/O Locations	Conversion Data	Bits in Core Storage 0 → 15					
INPUT	UC J S T	0000 1110 0101 0001 0011 0010 0010 0011					
OUTPT OUTPT +1 OUTPT +2	J S T	0101 0000 0000 0000 0010 1000 0000 0000 0010 0100 0000 0000					

<u>Calling sequence</u>

Label	Operation	T	F	IT			Operan	ds & Remarks		
н в	r 10	ι,	a 🔟		35 40	6	¥	55	60	•3
	L.I.B.F		Γ		P.A.P.H.L.	CAL	L IB.		N.V.E.R.S.	I.O.N.
	DC.	П			1 ddde		TROL	PARAA	ETER .	
	ac.				LNP.U.T.	L.J.M.P.	ILT. A	REA: AL	DRESS	
	0.0				OUTPUT.	OULT	PUT	AREA A	DDBES	ടപ
	OC .	П		11		C.H.A.	RACIT	ER COL	LAGT	است
	·	П	1	11						ال ا
	·		Τ	П	1					ه بب
		П			L.L.L.					
I.N.P.U.T	BSS			П	9	TIN P	UT A	REA L+		
	Б		Ι	Γ			باير ب			ه ب ب
	T				بنبيني	بالتند				ال ا
	•			П		ب ب ب	ىلىت	بالمساليات	<u></u>	
OUTPT	85.5				h	O.U.T	PUT	AREA		
	1	П	Т				بالمراب الم			ليبي

where

d is the case initialization digit,

e indicates the direction of conversion,

f indicates the character count,

g is the length of the input area. g must be equal to or greater than f if e is 0. If e is 1, g must be equal to f/2, or (f+1)/2 if f is odd.

h is the length of the output area. If e is 0, h must be equal to or greater than f/2, or (f+1)/2, if f is odd. If e is 1, h must be equal to or greater than f.

Control Parameter

This parameter consists of four hexadecimal digits. Digits 1 and 2 are not used. The third digit indicates whether or not the case is to be initialized before conversion begins:

0 - Initialize case

1 - Do not alter case

The fourth digit indicates the type of conversion:

0 - PTTC/8 to IBM Card Code

1 - IBM Card Code to PTTC/8

Input

Input is either PTTC/8 or IBM Card Code characters (as specified by the control parameter) starting in location INPUT. PTTC/8 characters are packed two per binary word; IBM Card Code characters are not packed.

Output

Output is either IBM Card Code or PTTC/8 code characters starting in location OUTPT. PTTC/8 codes are packed two per binary word; IBM Card Code characters are not packed.

If the conversion is IBM Card Code input to PTIC/8 output, the input area may overlap the output area if the address INPUT is equal to or greater than the address OUTPT. Case-shift characters are inserted in the output message where needed to define certain graphic characters (see "PAPEB").

If the conversion is PTTC/8 input to IBM Card Code output, the input area may overlap the output area if the address INPUT+n/2 is equal to or greater than the address OUTPT+n, where n is the character count. The subroutine starts processing at location INPUT.

Character Count

This parameter specifies the number of PTTC/8 or EBCDIC characters in the input area. The count must include case-shift characters, even though they will not appear in the output. Because the input may be packed, the character count may not be equal to the number of binary words in the input area.

There is no danger of overflowing the output area if the number of words in the output area is equal to the number of characters in the input area.

Errors Detected

Any input character not marked by an asterisk in Appendix D is an error.

Subroutine Operation

Case- and shift-character handling is described under "PAPEB".

If an odd number of PTTC/8 output characters is produced, bits 8-15 of the last used word in the output area are set to a delete character.

The conversion is halted when the character count is decremented to zero or when a new-line (NL) control character is read.

PAPPR

This subroutine converts PTTC/8 subset to either Console Printer or 1403 Printer code. The conversion to 1403 Printer code is illustrated below:

1/O Locations	Conve Dat	rsion ta	Bits in Core Storage 0
	UC	J	0000 1110 0101 0001
OUTPT	J	\$	0101 1000 0110 0010

Calling Sequence

Label	Operation		F	Ы				Operands	& Remarks		
п в	v 3		22	n	35	4	45	50	ы	مه	45
	LIBA	-	Γ	П	PAP.	P.R	. CAL	L. PITI	C/18. C	ONVER	SION
	DC.				1.00	die	COM	TROIL	PARAM	ETER	. آب ب
	D.C.		L		TNP	U.T	, I INP	UT AR	EA AD	DRESS	
	00		Г		OUT.	PIT	, OUT	PUT A	REAL A	DDRES	SUL
1.1.1.1	DC				f1		CH.A	RACITIE	RICOU	ANT I	
				П						للبلب	ىلى ب
أسبب	1									ليت كما الم	بيني
INPUT	8.5.5				9		INIP	UT AA	EA.	ببعبت	بالبب
	·		L		بيعل	يبالبني			<u></u>	خير	
	Juin				1	ي ال					ليناسبون
					1 1						. السادي
OUTPT	855.				h			PUTLA	REAL		سلسب
		1	Г	IT	1						

where

<u>Character Count</u>

d is the case initialization digit,

e is the output printer code digit,

f is the number of characters in the input area to be converted,

g is the length of the input area. g must be equal to or greater than f/2 if the character count is even, (f+1)/2 if the character count is odd.

h is the length of the output area. h must be equal to or greater than f/2, minus the number of paper tape control characters in the input area, plus 1 if the result is odd.

Control Parameter

This parameter consists of four hexadecimal digits. Digits 1 and 2 are not used. The third digit indicates whether or not the case is to be initialized before conversion begins:

0 - Initialize case 1 - Do not alter case

The fourth digit determines the output printer code.

- 0 Console Printer code
- 1 1403 Printer code

Input

Input consists of PTTC/8 characters starting in location INPUT. PTTC/8 characters are packed two per binary word. All control characters except case-shift (LC or UC) characters are converted to output. Case-shift characters are used only to define the case mode of the graphic characters that follow.

Output

Output consists of either Console Printer or 1403 Printer characters starting in location OUTPT. This code is packed two characters per binary word. If overlap of the input and output areas is desired, the address INPUT must be equal to or greater than the address OUTPT. This is necessary because the subroutine starts processing at location INPUT. This parameter specifies the number of PTTC/8 characters in the input area. The count must include case-shift characters, even though they do not appear in the output. Because the input is packed, the character count is not equal to the number of binary words in the input area.

If an odd number of output characters is produced, bits 8-15 of the last used word in the output area are set to a space character.

The conversion is halted when the character count is decremented to zero or when a new-line (NL) control character is read.

Errors Detected

Any input character not marked by an asterisk in Appendix D is an error.

HOLPR

This subroutine converts IBM Card Code subset to either Console Printer or 1403 Printer code. The conversion to 1403 Printer code is illustrated below.

I/O Locations	Conversion Data	Bits in Core Storage 0
INPUT	J	0101 0000 0000 0000
INPUT + I	,	0010 0100 0010 0000
OUTPT	J,	0101 1000 0001 0110

Calling Sequence

124		Caeronds & Kemarks
	LIBF HOLPR	CALL CARD CODE CONVERSION
	D.C	CONTROL PARAMETER
	D.C	LINPUT, AREA ADDRESS
1	D.C	OUTPUT AREA ADDRESS
	D.C	CHARACILER COUNT
		العاميية فيمام متقنفته فتدافعا فالاعار والالالان
		المتحصيل والمتحدة المتعلية ووالالمتحد والمتحد والمتحد والمتحد
		المحاجبة المتعاج المتعادية المتعادية المتعادية المتعادية المتعار والمراج
I.N.P.U.T	B.S.S	. INPUT AREAL
4		المتحجية والمتعاد والمتعام والمتعاد والمت
	· · · · · · · · · · · · · · · · · · ·	المحاف ستلب سنجت الالتحالة بفتقا بالانفاذ فتقابر يفارق الرابا الداع
	L.•	الالمتحاجية بالالتان فالقاعات والفارة والفاطية فتقتله فالالا فالمراج
D.U.T.P.T	BS.S. h.	OULT PUT AREAL
		الابار بالمتشيب الانساف فتحاط تعالف والالا فالعتبع بفاطا والع

e is the output printer code digit,

f is the number of characters in the input area to be converted,

g is the length of the input area. g must be equal to or greater than f.

h is the length of the output area. h must be equal to or greater than f/2.

Control Parameter

This parameter consists of four hexadecimal digits. Digits 1-3 are not used. The fourth digit determines the output printer code.

0 - Console Printer code

1 - 1403 Printer code

Input

Input consists of IBM Card Code characters, starting in location INPUT. The characters are not packed.

Output

Output consists of either Console Printer or 1403 Printer characters, starting in location OUTPT. The code is packed two characters per binary word.

The input area may overlap the output area if the address INPUT is equal to or greater than the address OUTPT. The subroutine starts processing at location INPUT.

Character Count

This number specifies the number of IBM Card Code characters to be converted and is equal to the number of words in the input area. If an odd count is specified, bits 8-15 of the last word used in the output area are not altered. Any input character not marked with an asterisk in Appendix D is an error.

EBPRT

This subroutine converts EBCDIC subset to either Console Printer or 1403 Printer Code. The conversion to 1403 Printer code is shown below.

I/O	Conversion	Core Storage Bits
Locations	Data	0
INPUT	LE	1101 0011 1100 0101
INPUT + 1	ES	1100 0101 1110 0010
OUTPUT	LE	0001 1010 0110 1000
OUTPT + 1	ES	0110 1000 0000 1101

Calling Sequence

Lovel	Currentine		Liperandi & femanis
	LIBE	11	EBPRT
	DC.	11	10000 CONTROL PARAMETER
	DC.		INPUT INPUT AREA ADDRESS
[]	D.C.	TI	OUTPT OUTPUT AREA ADDRESS
	bc.	T	FL
1	•••••		المراجع والمراجع
	$[\cdot, \cdot, \cdot]$		l Latar and a state of the stat
	l. l.		i
I.N.P.U.T	8.5.5		9. LN.P.U.T. A.R.E.A.
	·	TI	I CALLER AND A REPORT OF A DESCRIPTION O
	••••		I CALLER AND A REPORT AND A R
	· · · · · · · · · · · · · · · · · · ·	П	
O.U.T.P.T	ASS.	T	hu OUT.P.U.T. A.R.E.A.
	T_{1}, T_{1}	П	

where

e is the output printer code digit,

f is the number of characters in the input area to be converted,

g is the length of the input area. g must be equal to or greater than f/2.

h is the length of the output area. h must be equal to or greater than f/2.

Control Parameter

This parameter consists of four hexadecimal digits. Digits 1-3 are not used. The fourth digit determines the output printer code.

- 0 Console Printer code
- 1 1403 Printer code

Input

Input consists of EBCDIC characters starting in location INPUT. EBCDIC characters are packed two per word.

Output

Output consists of either Console Printer or 1403 Printer code starting in location OUTPT. The code is packed two characters per binary word.

The address INPUT must be equal to or greater than the address OUTPT if overlap of the input and output areas is desired. The subroutine starts processing at location INPUT.

I/O Locations	Conversion Data	Core Storage Bits 0
Accumulator	+001(777010	0000 0001 0000 0000
Extension	10010///218	0000 0000 0000 0010

1/O Locations	Conversion Data	Core Storage Bits 0
OUTPT	+	1000 0000 1010 0000
OUTPT + 1	0	0010 0000 0000 0000
OUTPT + 2	0	0010 0000 0000 0000
OUTPT +3	· 1	0001 0000 0000 0000
OUTPT+4	6	0000 0001 0000 0000
OUTPT +5	7	0000 0000 0100 0000
OUTPT +6	7	0000 0000 0100 0000
OUTPT +7	7	0000 0000 0100 0000
OUTPT +8	2	0000 1000 0000 0000
OUTPT +9	1	0001 0000 0000 0000
OUTPT + 10	8	0000 0000 0010 0000

Calling Sequence

Character Count

This parameter specifies the number of EBCDIC characters to be converted. This count is not equal to the number of words in the input area. If an odd count is specified, bits 8-15 of the last word used in the output area are not altered; however, these bits may cause print checks if they comprise an illegal character.

Errors Detected

Any input character not marked with an asterisk in Appendix D is an error.

BIDEC

This subroutine converts a 32-bit binary value to its decimal equivalent in ten IBM Card Code numeric characters and one sign character. The conversion is illustrated below:

Latel	Γ	Operation		F	,	Γ	Operandis & Remarks
				5	:		a ab ab a a a a a a a a a a a a a a a a
	L	LIBF			Γ		BIDEC CALL AINARY COMMERSION
		D.C	L				OUTPUT, OUTPUT, AREA, ADDRESS,
		·		Ļ	1	L	
		••••		L	L		
	L	·		L	Ĺ	L	<u> </u>
0.0.7.9.1		B.S.S.	L				1.1
	Ĺ			1	L	L	

Input

Input is a 32-bit binary value in the Accumulator and Extension.

Output

Output is an IBM Card Code sign character (+ or -) in location OUTPT, and ten IBM Card Code numeric characters in OUTPT+1 through OUTPT+10.

Errors Detected

The BIDEC subroutine does not detect errors.

DECBI

This subroutine converts a decimal value consisting of ten IBM Card Code numeric characters and a sign character to a 32-bit binary word. This subroutine is the opposite of the BIDEC subroutine (see above) except that fewer than ten characters may be specified.

Calling Sequence

اجتما		Operation	Γ	ş	1	Γ									0	pera	nda ð	. Rem	arks							_		_		
n 15		7 F		n	33		13		-			-				*			33				*		_			65		
		LIBE					DEC.	BZ.				6	ŝ	سانيا		DE	4	Z A	1	Ē.		: 0	24			.8	s	I	0	AL.
المسلط		00.	L	L		L	IN.P.	UT.			-	I	NJ	2.14	L	A	8.	E.A		A.	2	2.6		s	s					_
		0.0	L				MOC	ALT.			<u> </u>	.in	0	R.D.		<u>.</u>	u.	M.T	1	A.	D.I	20	Æ		s	<u>د</u>				_
		·	L	L		L	L									1.			<u>ب</u>								•		L	-
لسبب	_	· · · · ·	L	L	L	L									_				-					ė	-	-	•			_
		·	L	L	L	L													L.											
NDCNT		0.0.	L	L	L	L				••••	-	.15	6	20		C-0	ų,	AL Z	ι.	L								_		ĥ
ببيب		<u></u>	L	L	L	┢	fiere								Ŀ		Ĺ.	-	. 1			ı							_	
		••••	L	L	L		.									1												<u> </u>	L.	_
	_	·	L	L	L	L	ـــــا								ı	ł								L	_					-+
LARUT		BS.S.	L	L	L	L						1	Alu	2.0	L	<u>م</u> د.	.B.	E.A						i.						
لي يوس			L	L	L	L								ت.	_							-					•		L.	-+

Errors Detected

Any of the following conditions causes the Overflow indicator to be turned on, the Carry indicator to be turned off, and an immediate exit to be made back to the caller:

- 1. Any sign other than a plus, minus, blank, or ampersand.
- Any character other than a space or 0 2. through 9.
- 3. Any converted value greater than +2,147,483,647 or less than -2,147,483,648.

ZIPCO

This subroutine supplements all standard conversions except those involving PTTC/8 code. It offers the user the option of supplying his own conversion tables and codes. ZIPCO uses direct table access and is considerably faster than the other conversion subroutines.

Operands & Remarks

L.I.B.F. ZIPCO ... CALL SPECIAL COMPERSION DC. //b.Cd. ... COMTROL PARAMETER DC. INPUT INPUT AREA ADDRESS

 DC.
 LIMIPUTT.
 UMAPUTT.
 <th

Calling Sequence

where

a is the number of characters to be converted not including the sign character.

b is the length of the input area. b must be equal to at least a plus 1.

Input

Input is an IBM Card Code sign character in location INPUT, the address (WDCNT) of the number of characters (1 to 10) to be converted, and specified number of characters in IBM Card Code in locations INPUT+1 through INPUT+N (where N = 1, 2,...10).

Output

Output is a 32-bit binary word, containing the converted value, in the Accumulator and Extension.

where

<u>|</u>

OUTPT B.S.S.

b is the input code digit,

c is the packed input digit,

d is the output code digit,

e is the packed output digit,

f is the number of characters to be converted,

g is the length of the input area,

h is the length of the output area,

j is the name of the conversion table to be used. This CALL is not executed; however, it is required following the character count parameter to cause the loading of the desired conversion table, provide the address of that table to ZIPCO, and provide information required by ZIPCO for the return to the calling program.

Control Parameter

r

This parameter consists of four hexadecimal digits as follows:

Diwit		1 for 12-bit IBM Card Code input
Digit	I	0 for all other types of input
Digit	2	1 for unpacked input
Digie	2	0 for packed input
Digit	3	1 for 12-bit IBM Card Code output
	5	0 for 8-bit IBM Card Code and all other types of output
Diait	ц.	f for unpacked output
Digit	-	0 for packed output

Input

Input consists of packed or unpacked characters in the code specified by the conversion table and starting at location INPUT.

Output

Output consists of packed or unpacked characters in the code specified by the conversion table and starting at location OUTPT.

Character Count

This parameter specifies the number of input characters to be converted. If an odd count is specified with packed input, bits 8-15 of the last word used in the output area are not altered.

<u>Table</u>

The type of conversion is determined by the table called with ZIPCO. The user may call one of the IBM-supplied conversion tables or he may supply his own.

The following IBM-supplied System Library tables may be called with ZIPCO.

- EBCCP EBCDIC to Console Printer Code.
- EBHOL EBCDIC to IBM Card Code.
- EBPT3 EBCDIC to 1403 Printer code.
- CPEBC Console Printer code to EECCIC.
- CPHOL Console Printer code to IBM Card Code.
- CPPT3 Console Printer code to 1403 Printer code.
- HLEBC IBM Card Code to EBCDIC.
- HOLCP IBM Card Code to Console Printer code.
- HLPT3 IBN Card Code to 1403 Printer code.
- PT3EB 1403 Printer code to EECCIC.
- PT3CP 1403 Printer code to Console Printer Code.
- PTHOL 1403 Printer code to IBM Card Code.

Each conversion table consists of 256 characters-- 128 words with two 8-bit characters per word. The seven low-order bits of the character to be converted (input character) are used as an address. The address designates the position in the table of the corresponding conversion character. The high-order bit (bit 0) of the input character designates which half of the table word is to be used. When bit 0 is 1, the left half of the word is used. When bit 0 is 0, the right half of the word is used. All dummy entries of the IEM-supplied tables contain the code for a blank.

The following is an example of the conversion performed by ZIPCO. The tables show (1) the input EBCDIC values, (2) the table EBPT3 used for the conversion, and (3) the output characters in 1403 Printer code.

input Location	Value	
INPUT	1111 0010 0111 0010	
INPUT+1	0000 0000 1000 0000	
INPUT +2	0111 1111 1111 1111	
Table Location	Value	
EBPT3	0111 1111 0111 1111	
EBPT3 + 1	0111 1111 0111 1111	
•	•	
	•	
EBPT3 + 114	0000 0001 0111 1111	
	•	
	•	
EBPT3 + 127	0111 1111 0111 1111	
F		
Output Location	Value	1403 Print Character
OUTPT	0000 0001 0111 1111	2, b
OUTPT + 1	0111 1111 0111 1111	Ь, Ь
OUTPT +2	0111 1111 0111 1111	6,6

When 12-bit IBM Card Code is specified as input (or output), ZIPCO performs a packing (or unpacking) of the character to 8-bits (or 12 bits). The 1-7 row punches

EBCDIC	TO 1403 C	ONV TABLE	FOR Z	IPCO			SOURCE
ADDR RE	L OBJECT	ST.NO.	LABEL	OPCD I	FT OPERANDS		
0000	050978F3	0020		ENT	EBPT3		
0000 0	7F7F	0021	EBPT3	DC	/7F7F	NO GRAPHIC	NUL
0001 0	7F7F	0022		DC	/7878	NO GRAPHIC	NO GRAPHIC
0002 0	7F7F	0023		DC	/7F7F	NO GRAPHIC	NO GRAPHIC
0003 0	7F7F	0024		DC	/7F7F	NO GRAPHIC	NO GRAPHIC
0004 0	7F7F	0025		DC	/7F7F	NO GRAPHIC	PF
0005 0	7F7F	0026		DC	/7F7F	NO GRAPHIC	HT
0006 0	7F7F	0027		DC	/7F7F	NO GRAPHIC	LC
0007 0	7F7F	0028		DĊ	/7F7F	NO GRAPHIC	DEL
0008 0	7F7F	0029		DC	/7F7F	NO GRAPHIC	NO GRAPHIC
:	:	:		1	•	:	:
ē	ě.	÷		i	:	•	:
0072 0	017F	0135		DC .	/017F	2	NO CRADUTC
0073 0	027F	0136		00	/027F	2	NO CRAPHIC
0074 0	437F	0137		DC	/437F	å	NO GRAPHIC
0075 0	047F	0138		ñC.	/047F	5	NO GRAPHIC
0076 0	457F	0139		nč	/457F	6	NO GRAPHIC
0077 0	467F	0140		ñč.	/467F	7	NO GRAPHIC
0078 0	077F	0141		DC	/077F	8	NO GRAPHIC
0079 0	,087F	0142		DC	/087F	9	
007A 0	7F7F	0143		DC	/7F7F	NO GRAPHIC	NO GRAPHIC
007B 0	7F7F	0144		ĎČ	/7F7F	NO GRAPHIC	NO GRAPHIC
007C O	7F7F	0145		DC	/7F7F	NO GRAPHIC	NO GRAPHIC
007D 0	7F08	0146		DC	/7F08	NO GRAPHIC	1
007E 0	7F4A	0147		DC	/7F4A	NO GRAPHIC	+
807F 0	7F7F	0148		80	/7275	NO COMPUTE	NO COADUIC
		0110		UÇ.	////	NU GRAFNIL	NU GRAPHIC

Figure 20.1 System Library EBPT3

on the card are expressed as a 3-bit hexadecimal number (there can never be more than one punch between the 1 and 7 row). In this format a 1 punch would be expressed as 001, a 7 punch as 111. The punches in the other card rows: 12, 11, 0, 8, and 9, are transferred directly.

For example, take the IBM Card Code character "+" which is a 12, 6, 8 punch.



Errors Detected

No errors are detected by ZIPCO.

Figure 20.1 is a sample of the System Library table EBPT3 (EBCDIC to 1403 Printer code) which may be called with ZIPCO.

Arithmetic and Functional Subroutines

The IBM 1130 Subroutine System Library includes the arithmetic and functional subroutines that are the most frequently required because of their general applicability. There are 44 subroutines, some of which have several entry points.

Figure 21 lists the arithmetic and functional subroutines that are included in the Subroutine System Library

REAL DATA FORMATS

Many of the IBM 1130 arithmetic and functional subroutines offer two ranges of precision: standard and extended. The standard precision provides 23 significant bits, and the extended precision provides up to 31 significant bits. The magnitude of a real number must not be greater than 2^{127} or less than 2^{-128} (approximately 10^{38} and 10^{-39}).

To achieve correct results from a particular subroutine, the input arguments must be in the proper format.

Standard-Precision Format

Standard-precision real numbers are stored in core storage as shown below:



Numbers can consist of up to 23 significant bits (mantissa) with a binary exponent ranging from -128 to +127. Two adjacent storage locations are required for each number. The first (lowest) location must be even-numbered. The sign of the mantissa is in bit zero of the first word. The next 23 bits represent the mantissa (2's complement if the number is negative) and the remaining 8 bits represent the characteristic. The mantissa is normalized to fractional form, i.e., the implied binary point is between bits zero and one. The characteristic is formed by adding +128 to the exponent. For example, an exponent of -32 is represented by a characteristic of 128-32, or 96. An exponent of +100 is represented by a characteristic of 100 + 128, or 228. Since 128_{40} =/80 the characteristic of a nonnegative exponent always has a 1-bit in position 1, while the characteristic of a negative exponent always produces a 0-bit in position 1. A normal zero consists of all zero bits in both the characteristic and the mantissa.

Extended-Precision Format

Extended-precision real numbers are stored in three adjacent core locations as shown below:



Numbers can consists of up to 31 significant bits with a binary exponent ranging from -128 to +127; however, normalization can, in some cases, cause the loss of 1 bit of significance.

Bits zero through seven of the first word are unused; bits eight through 15 of the first word represent the characteristic of the exponent (formed in the same manner as in the standard range format); bit zero of the second word contains the sign of the mantissa; and the remaining 31 bits represent the mantissa (2's complement if the number is negative).

Real Negative Number Representation

Real negative numbers differ from real positive numbers in only one respect; the mantissa is always the 2's complement of the equivalent positive value. Example:

+.53125 is represented in core as 44000080

-.53125 is represented in core as BC000080

+4.0 is represented in core as 40000083

-4.0 is represented in core as C0000083

Note that a real negative number is never represented by a value of 800000xx, where xx is any characteristic between 00 and FF. The mantissa value of 800000 is its own 2's complement and therefore lies outside the definition of a real negative number, i.e., the 2's complement of its absolute value.

Fixed-Point Format

Fractional numbers, as applied to the fixed-point subroutines, XSQR, XMDS, XMD, and XDD, are defined as binary fractions with implied binary points of zero. That is, the binary point is positioned between the sign (bit 0) and the most significant bit (bit 1).

The user can consider the binary point to be in any position in his fixed-point numbers. To correctly interpret the results the following rules must be observed:

- 1. Only numbers with binary points in equivalent positions can be correctly added or subtracted.
- 2. The binary point location in the product of two numbers is the sum of the binary point locations of the multiplier and the multiplicand.
- 3. The binary point location in the quotient of two numbers is the difference between the binary point locations of the dividend and the divisor.
- 4. The binary point location in a number that is input to the fixed-point square root subroutine (XSQR) must be an even number from 0-14. The binary point location in the output root is half the binary point location of the input number.

REAL NUMBER PSEUDO ACCUMULATOR

IBM 1130 real number subroutines sometimes require an accumulator that can accommodate numbers in real number format. Since all of the 1130 registers are only 16 bits in length, a pseudo accumulator must be set up to contain two- or three-word real numbers. The pseudo accumulator (designated FAC for floating accumulator) is a three-word register occupying the three highest locations of the Transfer Vector (see IBM <u>1130/1800 Assembler Language</u>). The user can refer to these words by using Index Register 3 plus a fixed displacement (XR3+125, 126, or 127). The format of the FAC is shown below.

Characteristic	Mantissa	Mantissa
 XR3 + 125	XR3 + 126	XR3 + 127
	FAC	

The effective address of the mantissa is always even. The eight rightmost bits of the FAC are zero when using standard precision.

Note: Arithmetic and functional subroutines do not save and restore the contents of the 1130 Accumulator or the Extension. The calling program should provide for this if the contents are significant. When execution of the user's program begins, all three words of FAC contain zeros. Results of arithmetic and functional subroutines are truncated.

CALLING SEQUENCES

The arithmetic and functional subroutines are called via a CALL or LIBF statement (whichever is required) followed, in some cases, by a DC statement containing the actual or symbolic address of an argument. In the descriptions that follow, the notations (ARG) and (FAC) refer to the contents of the operand rather than its address. The name FAC refers to the real number pseudo accumulator. The extended-precision subroutine names are prefixed with the letter E (subroutines that handle both precisions have the same name and do not have a prefix).

SUBROUTI NE	NAME			
Real (Floating Point)	Standard Precision	Extended Precision		
Add/Subtract Multiply Divide Load/Store FAC Trigonometric Sine/Cosine Trigonometric Arctangent Square Root Natural Logarithm Exponential (e) Hyperbolic Tangent Real Base to an Integer Exponent Real Base to a Real Exponent Real to Integer Integer to Real	<pre>*FADD/*FSUB *FMPY *FDIV *FLD/*FSTO FSINE/FCOSN, FSIN/FCOS FAIN, FAIAN FSQR, FSQRI FLN, FALOG FXPN, FEXP FINH/FIANH *FAXI FAXB IFIX FLOAT</pre>	* EADD/* ESU B * EMPY * EDIV * ELD/* ESTO ESINE/ECOSN, ESIN/ECOS EATN, EATAN ESQR, ESQRT ELN, EALOG EXPN, EEXP ETNH/ETANH * EAXI * EAXB IFIX FLOAT		
Normalize Real Binary to Decimal/Real Decimal to Binary Real Arithmetic Range Check Fixed-Point	NORM FBTD/FDTB FARC	NORM FBTD/FDTB FARC		
Integer Base to an Integer Exponent Fixed-Point Square Root Fixed-Point Fractional Multiply (short) Fixed-Point Double Word Multiply Fixed-Point Double Word Divide	*FIXI XSQR XMDS XMD XMD XDD	*FIXI XSQR XMD XDD		
Special Function				
Real Reverse Subtract Real Reverse Divide Real Reverse Sign Real Absolute Value Integer Absolute Value	*FSBR *FD VR SNR FAVL, FABS IABS	*ESBR *EDVR SNR EAVL, EABS IABS		
Miscellaneous				
Get Parameters	FGEIP	EGET P		
Note: By adding an X to those names prefixed with an asterisk, the user can cause the contents of Index Register 1 to be added to the address of the argument specified in the subroutine calling sequence to form the effective argument address. For example, FADDX would be the modified form of FADD.				

Figure 21. Arithmetic and Functional Subroutines

Note also that some of the functional subroutines can be called via two different calling sequences. One calling sequence assumes the argument is in FAC; the other specifies the location of the argument with a DC statement.

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In addition, some subroutines can have indexed linkage to the argument. The calling sequence is the same except for the subroutine name which contains an X suffix. Also, some subroutines perform more than one type of arithmetic or function. For example, FSIN and FCOS are different entry points to the same subroutine. Each subroutine is listed in Figure 21 with the corresponding entry points.

Real Add

LIBF	FADD, FADDX, EADD or EADDX
DC	ARG
Input	Real augend in FAC Real addend in location ARG
Result	(FAC) + (ARG) replaces (FAC)

Real Subtract

LIBF FSUB, FSUBX, ESUB or ESUBX DC ARG Page of GC26-5929-6 Revised May 21, 1971 By TNL GN33-8112

Input	Real	min	uend	in	FAC			
_	Real	sub	trahe	enđ	in	loca	ition	ARG
Result	(FAC)		(ARG)	re	epla	aces	(FAC)	

Real Multiply

LIBF	FMPY, FMPYX, EMPY or EMPYX
DC	ARG
Input	Real multiplicand in FAC
-	Real multiplier in location ARG
Result	(FAC) times (ARG) replaces (FAC)

Real Divide

LIBF	FDIV, FDIVX, EDIV or EDIVX
DC	ARG
Input	Real dividend in FAC
	Real divisor in location ARG
Result	(FAC) / (ARG) replaces (FAC)

Note: On a divide by zero, the divide check indicator is turned on, the dividend is not changed, and the dividend remains in FAC.

Load FAC

LIBF FLD, FLDX, ELD or ELDX DC ARG Real number in location ARG Input (ARG) replaces (FAC) Result

Store FAC

LIBF FSTO, FSTOX, ESTO or ESTOX DC ARG Real number in FAC Input Result (FAC) replaces (ARG)

Real Trigonometric Sine

CALL FSINE or ESINE Real argument (in radians) in FAC Input Sine of (FAC) replaces (FAC) Result

or

CALL FSIN or ESIN DC ARG Input Real argument (in radians) in location ARG Result Sine of (ARG) replaces (FAC)

Real Trigonometric Cosine

CALL	FCOSN or ECOSN
Input	Real argument (in radians) in FAG
Result	Cosine of (FAC) replaces (FAC)

or

CALL	FCOS or ECOS
EC	ARG
Input	Real argument (in radians) in
	location ARG
Result	Cosine of (ARG) replaces (FAC)

Real Trigonometric Arctangent

CALL DC Input Result	FATN or EATN ARG Real argument in FAC Arctangent of (FAC) replaces (FAC); the result lies within the range $\pm \frac{\pi}{2}$ radians (±90 degrees)			
	or			
CALL DC	FATAN Or EATAN ARG			

Input Real argument in location ARG Arctangent of (ARG) replaces Result (FAC); the result lies within the range ± $\frac{\pi}{2}$ radians (± 90 degrees)

Real Square Root

CALL Input Result	FSQR or ESQR Real argument in FAC Square root of (FAC) replaces (FAC)			
	or			
CALL DC Input	FSQRT or ESQRT ARG Real argument in location ARG			

Result Square root of (ARG) replaces (FAC)

Real Natural Logarithm

CALL	FLN or ELN	
Input	Real argument in FA	C
Result	Log _e (FAC) replace	es (FAC)

or

CALL	FALOG OF EALOG
DC	ARG
Input	Real argument in location ARG
Result	Log _n (ARG) replaces (FAC)

Real Exponential

CALL FXPN or EXPN Real argument in FAC = nInput Result eⁿ replaces (FAC)

or

CALL	FEXP or EEXP
DC	ARG
Input	Real argument in location ARG = n
Result	e ⁿ replaces (FAC)

Real Hyperbolic Tangent

CALL	FTNH or ETNH
Input	Real argument in FAC
Result	TANH (FAC) replaces (FAC)
CALL DC	or FTANH or ETANH ARG

Input Real argument in location ARG Result TANH (ARG) replaces (FAC)

Real Base to an Integer Exponent

LIBF	FAXI, FAXIX, EAXI, or EAXIX Arg
Tanut	Bool bago in FAC
Input	Integer exponent in location ARG
Result	(FAC) raised to the exponent
	(ARG) replaces (FAC)

Real Base to a Real Exponent

CALL	FAXB, FAXBX, EAXB or EAXBX
DC	ARG
Input	Real base in FAC Real exponent in location ARG
Result	(FAC) raised to the exponent (ARG) replaces (FAC)

Real to Integer

LIBF	IFIX
Input	Real number in FAC
Result	Integer in the Accumulator

Integer to Real

LIBF	FLOAT	
Input	Integer in the	Accumulator
Result	Real number in	FAC

<u>Normalize</u>

LIBF NORM

Input Real unnormalized number in FAC Result The mantissa portion of FAC is shifted until the most significant bit resides in bit position 1. The characteristic is changed to reflect the number of bit positions shifted.

Real Binary to Decimal

CALL FBTD DC LDEC

Input Real number in FAC

Result A string of EBCDIC-coded data starting at location LDEC. Each EBCDIC character occupies the rightmost 8 bits of a word. The last character of the string is a blank.

The output format is exactly as follows:

sd.dddddddEsddb

where:

s represents a sign (plus or minus) d represents one of the decimal digits 0-9 b represents a blank Real Decimal to Binary

- CALL FDTB
- DC LDEC
- Input A string of EBCDIC coded data at location LDEC. Each EBCDIC character occupies the rightmost 8 bits of a word. The leftmost 8 bits <u>must</u> be zeros. The first character of the input must be the sign (plus or minus). Following the sign, one to nine decimal digits (0-9) may be specified. The decimal point may appear before, within, or after the decimal digits. Immediately Immediately after the last decimal digit (or decimal point), the exponent is specified as follows.

Esddb

where:

s represents the sign of the exponent (plus or minus) d represents one of the decimal digits (0-9) b represents a blank (the blank is required to indicate the end of the string)

No embedded blanks may appear in the input string as the first blank is interpreted as the end of the data.

Result Real number in FAC

Real Arithmetic Range Check

LIBF FARC

Result This subroutine checks for real number overflow or underflow, and sets programmed indicators for interrogation by a FORTRAN program.

Integer Base to an Integer Exponent

LIBF	FIXI	or	FIXIX

DC ARG Input Integer base in the Accumulator Integer exponent in location ARG Result (Accumulator) raised to the exponent contained in ARG replaces (Accumulator)

Fixed-Point Square Root

CALL XSQR

Input	Fixed-point fractional argument
_	(16 bits only) in the Accumulator.
Result	Square root of (Accumulator)

replaces (Accumulator). If the argument is negative the absolute value is used and the Overflow indicator is turned ON.

Fixed-Point Doubleword Multiply

LIBF	XMD
Input	Doubleword fractional multiplier
-	in FAC (addressed by XR3 + 126)
	Doubleword fractional multiplicand
	in the Accumulator and Extension
Result	Doubleword fractional product in
	the Accumulator and Extension

Fixed-Point Fractional Multiply

LIBF XMDS Input Doubleword fractional multiplier in the Accumulator and Extension Doubleword fractional multiplicand in FAC (addressed by YP3 + 126)

in FAC (addressed by XR3 + 126) Result Product in the Accumulator and Extension (XMDS is shorter and faster than XMD; however, the resulting precision is 24 bits).

Fixed-Point Doubleword Divide

LIBF XDD

- Input Doubleword fractional dividend in FAC (addressed by XR3 + 126) Doubleword fractional divisor in Accumulator and Extension
- Result Doubleword fractional quotient in the Accumulator and Extension. The double dividend in FAC is destroyed by the execution of the subroutine.

Real Reverse Subtract

LIBF	FSBR, FSBRX, ESBR or ESBRX
DC	ARG
Input	Real minuend in location ARG
	Real subtrahend in FAC
Result	(ARG) - (FAC) replaces (FAC)

Real Reverse Divide

LIBF FDVR, FDVRX, EDVR or EDVRX DC ARG Input Real dividend in location ARG Real divisor in FAC Result (ARG) / (FAC) replaces (FAC)

Note: On a divide by zero, the divide check indicator is turned on, the dividend is not changed, and the dividend remains in FAC.

<u>Real Reverse Sign</u>

LIBF SNR Input Real number in FAC Result - (FAC) replaces (FAC)

Real Absolute Value

CALL FAVL or EAVL Input Real number in FAC Result Absolute value of (FAC) replaces (FAC) CALL FABS OF EABS DC ARG Input Real number in location ARG Result Absolute value of (ARG) replaces (FAC)

Integer Absolute Value

or

CALL	IABS	
Input	An integer in ARG	
Result	Absolute value of (ARG) replac (Accumulator)	ces

Get Parameters (FGETP or EGETP)

Example:

MAIN	CALL	SUBR	
NEXT	DC etc.	ARG	
•	•		
•	•	•	
•	•	•	
SUBR	DC	0	
	LI BF	FGETP or	EGET P
SUBEX	DC	0	
	etc.		
•	•	•	
•	•	•	
•	•	•	
-	•	•	
	BSC I	SUBEX	

The FGETP subroutine performs two functions for a subroutine accessed by a CALL statement. It loads FAC with the contents of ARG; it sets SUBEX to return to NEXT in the calling program.

ARITHMETIC AND FUNCTIONAL SUBROUTINE ERROR INDICATORS

The highest three-word entry in the Transfer Vector is reserved for the real number pseudo accumulator (FAC). The next to highest three-word entry is reserved for the arithmetic and functional subroutine error indicators.

The first word (addressed XR3 + 122) of the second entry is used for real number arithmetic overflow and underflow indicators. The second word (XR3 + 123) is used for a divide check indicator, and the third word (XR3 + 124) is used for functional subroutine indicators. When execution begins, all three words contain zeros.

Word One

Each real number subroutine checks for exponent underflow and overflow. If either occurs, word one and FAC are set as follows.

- if overflow has occurred (FAC = ± maximum), word one is set to 1.
- if underflow has occurred (FAC = zero), word one is set to 3.

Word Two

The real number divide subroutines check for division by zero. If this occurs, word two is set to 1. The dividend is not changed and remains in FAC.

Word Three

The functional subroutines check for the following error conditions and set word three as described. All error conditions detected by the functional subroutines are indicated in word three.

<u>Real Natural Logarithm</u>. When the argument is zero, FAC is set to the largest negative value and a bit is moved into position 15 of word three with an OR instruction. When the argument is negative, the absolute value of the argument is used and a bit is moved into position 15 of word three with an OR instruction.

<u>Real Trigonometric Sine and Cosine</u>. When the absolute value of the argument is equal to or greater than 2^{24} , FAC is set to zero and a bit is moved into position 14 of word three with an OR instruction.

<u>Real Square Root</u>. When the argument is negative, the square root of the argument's absolute value is returned, and a bit is moved into position 13 of word three with an OR instruction.

<u>Real to Integer</u>. When the absolute value of the argument is greater than $2^{45}-1$, the largest possible signed result is placed in the accumulator and a bit is moved into position 12 of word three with an OR instruction.

Integer Base to an Integer Exponent. When the base is zero and the exponent is zero or negative, a zero result is returned and a bit is moved into position 11 of word three with an OR instruction.

<u>Real Base to an Integer Exponent</u>. When the base is zero and the exponent is zero or negative, a zero result is returned and a bit is moved into position 10 of word three with an OR instruction.

<u>Real Base Raised to a Real Exponent</u>. When the base is zero and the exponent is zero or negative, a zero result is returned and a bit is moved into position 9 of word three with an OR instruction. When the base is negative and the exponent is not zero, the absolute value of the base is used and a bit is moved into position 15 of word three with an OR instruction.

End of File (DM2 System Only). When the end-of-file record in the unformatted I/O area is read, a bit is moved into position 2 of word three with an OR instruction.

Functional Subroutine Accuracy

Given:

е		Maximum error
f (x)	=	True value of the function
f* (x)		Value generated by subroutine
(<+∞)	=	≤Largest valid real number
(>-∞)	=	2Most negative real number

$$e \equiv \left| \frac{\operatorname{atn}(x) - \operatorname{atn}^{*}(x)}{\operatorname{atn}(x)} \right| < 2.0 \times 10^{-9}$$

for the range

$$-3.88336148 \times 10^{37} \le x \le 3.88336148 \times 10^{37}$$

EEXP

The following statements of accuracy apply to extended precision subroutines.

$$\frac{ESIN}{e} = \frac{\sin(x) - \sin^{*}(x)}{x} < 3.0 \times 10^{-9}$$

for the range

$$-1.0 \times 10^{6} \le x < 0$$
$$1.0 \times 10^{6} \ge x > 0$$

for $x = 0 \sin(x) \equiv 0$

$$\frac{\text{ECOS}}{\text{e}} \equiv \left| \frac{\cos(x) - \cos^*(x)}{|x| + \frac{\pi}{2}} \right| < 3.0 \times 10^{-9}$$

for the range

$$-1.0 \times 10^{6} \le x \le 1.0 \times 10^{6}$$

$$e \equiv \left| \frac{e^{x} - (e^{x})^{*}}{e^{x}} \right| < \begin{cases} 2.0 \times 10^{-9} |x| \\ \text{or} \\ 2.0 \times 10^{-9} \end{cases} \text{ whichever} \\ \text{is greater} \end{cases}$$

for the range

$$-\ln(\infty) < x < \ln(\infty)$$

i.e., $0 < e^{X} < \infty$

<u>ELN</u>

$$e \equiv \left| \frac{\ln(x) - \ln^*(x)}{\ln(x)} \right| < 3.0 \times 10^{-9}$$

for the range

0 < x <∞
<u>ETANH</u>

$$e = |tanh(x) - tanh^*(x)| < 3.0 \times 10^{-9}$$

for the range

- « < X < ∞

<u>ESQRT</u>

 $e \equiv \left| \frac{\sqrt{x} - \sqrt{x} *}{\sqrt{x}} \right| < 1.0 \times 10^{-9}$

for the range

0 < x <∞

STANDARD PRECISION SUBROUTINES

The following statements of accuracy apply to the standard precision subroutines.

FSIN

$$e \equiv \left| \frac{\sin(x) - \sin^*(x)}{x} \right| < 2.5 \times 10^{-7}$$

for the range

$$-1.0 \times 10^{6} \le x < 0$$

1.0 x 10⁶ \ge x > 0
for x = 0 sin (x) = 0

FCOS

$$e \equiv \frac{\cos(x) - \cos^{*}(x)}{|x| + \frac{\pi}{2}} < 2.5 \times 10^{-7}$$

for the range

$$-1.0 \times 10^{6} \le x \le 1.0 \times 10^{6}$$

<u>FATAN</u>

$$e \equiv \left| \frac{\operatorname{atn}(x) - \operatorname{atn}^*(x)}{\operatorname{atn}(x)} \right| < 5.0 \times 10^{-7}$$

for the range

$$-3.883361 \times 10^{37} \le x \le 3.883361 \times 10^{37}$$

FEXP

 $e \equiv \left| \frac{e^{x} - (e^{x})^{*}}{e^{x}} \right| < \begin{cases} 2.5 \times 10^{-7} |x| \\ or \\ 2.5 \times 10^{-7} \end{cases}$ whichever is greater

for the range

$$-\ln(\infty) < x < \ln(\infty)$$
 i.e., $0 < e^{x} < \infty$

<u>Fln</u>

$$e = \left| \frac{\ln(x) - \ln^*(x)}{\ln(x)} \right| < 4.0 \times 10^{-7}$$

for the range

0 < x < 1 $1 < x < \infty$ for x = 1 ln (x) $\equiv 0$

<u>FTANH</u>

$$e \equiv |tanh(x) - tanh^*(x)| < 2.5 \times 10^{-7}$$

for the range

-∞ < x < +∞

FSORT

$$e = \left| \frac{\sqrt{x} - \sqrt{x^*}}{\sqrt{x}} \right| < 2.5 \times 10^{-7}$$

for the range

0 < x < ∞

Elementary Function Algorithms

The choice of an approximating algorithm for a given function depends on such considerations as expected execution time, storage requirements, and accuracy. For a given accuracy, and within reasonable limits, storage requirements vary inversely as the execution time. Polynomial approximating is used to evaluate the elementary functions to effect the desired balance between storage requirements and efficiency.

SINE-COSINE

Polynomial Approximation

Given a real number, x, n, and y are defined such that

 $\frac{x}{2\pi} = n + y$

where n is an integer and $0 \le y \le 1$. Thus, x = $2\pi n + 2\pi y$, and the identities are

 $\sin x = \sin 2\pi y$ and $\cos x = \cos 2\pi y$.

The polynomial approximation, F(z), for the function (sin $2\pi z$)/z is used where -1/4≤z≤1/4.

The properties of sines and cosines are used to compute these functions as follows:

 $\cos 2\pi y = F(z)$

where

z = 1/4-y in the range $0 \le y \le 1/2$ z = y-3/4 in the range $1/2 \le y \le 1$

 $\sin 2\pi y = F(z)$

where

z = y in the range $0 \le y \le 1/4$ z = 1/2-y in the range $1/4 \le y \le 3/4$

z = y-1 in the range $3/4 \le y \le 1$

Extended Precision

F (z) =
$$a_1 z + a_2 z^3 + a_3 z^5 + a_4 z^7 + a_5 z^9 + a_6 z^{11}$$

where

a 1	=	6.2831853071
a,	= -	-41.341702117
a,	=	81.605226206
ຊັ	= :	-76.704281321
ຊື	=	42.009805726
a ₆	= •	-14.394135365

Standard Precision

$$F(z) = a_1^{z} + a_2^{z}^{3} + a_3^{z}^{5} + a_4^{z}^{7} + a_5^{z}^{9}$$

where

$$\begin{array}{rcl} a_1 &=& 6.2831853\\ a_2 &=& -41.341681\\ a_3 &=& 81.602481\\ a_4 &=& -76.581285\\ a_5 &=& 39.760722 \end{array}$$

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ARCTANGENT

Polynomial Approximation

The subroutine for arctangent is built around a polymonial, F (x), that approximates Arctan (z) in the range $-.23 \le z \le .23$. The Arctan (z) for z outside this range is found by using the identities

 $\operatorname{Arctan}(-z) = -\operatorname{Arctan}(z)$

a nd

Arctan(z) =
$$a_k + Arctan \left(\frac{z - b_k}{zb_k + 1} \right)$$

where

$$a_k = \frac{k\pi}{7}, b_k = \tan a_k$$

and k is determined so that

$$\tan \frac{(2k-1)\pi}{14} \le |z| < \tan \frac{(2k+1)\pi}{14}$$
 k = 1, 2, 3.

Having determined the value of k appropriate to z, the transformation $x=(z-b_k)/(zb_k+1)$ puts x in the range $-\tan \pi/14 \le x \le \tan \pi/14$. The polynomial F(x) was chosen to be good over a range slightly larger (i.e., $.23 \le \tan \pi/14$) so that the comparisons to determine the interval in which z lies need be only standard precision accuracy.

Arctan (z) =
$$a_{k}^{a} + F(x) z \ge 0$$

 $a_{k}^{a} - F(x) z < 0$

Extended Precision

F (x) = x (1.0 -
$$a_1 x^2 + a_2 x^4 - a_3 x^6 + a_4 x^8)$$

where

 $\begin{array}{rcl} a_1 &=& .33333327142\\ a_1 &=& .19999056792\\ a_2 &=& .14235177463\\ a_3 &=& .09992331248 \end{array}$

Standard Precision

F (x) = x(1.0 -
$$a_1 x^2 + a_2 x^4 - a_3 x^6)$$

where

$$a_1 = .333329573$$

 $a_2 = .199641035$
 $a_2 = .131779888$

SQUARE ROOT

Square Root (x)

Let
$$x = 2^{2b}F$$
 when $25 \le F < 1$
then $\sqrt{X} = 2^{b}\sqrt{F}$

where $\sqrt{F} = P_i$

i = number of approximation

P₁ = AF + B as a first approximation followed by 2 Newton iterations

where

$$A = .875, B = .27863$$
 when $.25 \le F < .5$

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A = .578125, B = .421875 when
$$.5 \le F < 1$$

 $P_{2} = \begin{pmatrix} P_{1} + \frac{F}{P_{1}} \\ 2 \end{pmatrix}$ $P_{3} = \begin{pmatrix} P_{2} + \frac{F}{P_{2}} \\ 2 \end{pmatrix}$

Polynomial Approximation

Given a normalized real number

$$x = 2^{k} x f$$

where the range of f is $1/2 \le f < 1$, and j and g are found such that $x=2^{j}g$ where $(\sqrt{2}/2 \le g < \sqrt{2})$. This is done by setting j=k-1, g=2f if $f < \sqrt{2}/2$ and j=k, g=fotherwise.

Thus:

 $\ln(x) = j \cdot \ln(2) + \ln(g)$.

The approximation for $\ln(g)$, $\sqrt{2}/2 \le g < \sqrt{2}$, is based on the series

$$\ln \frac{v + x}{v - x} = 2 \left[(x/v) + (x^3/3v^3) + (x^5/5v^5) + \dots \right]$$

which converges for (-v<x<v). With the transformation

 $x = v \frac{g-1}{g+1}$, $v = (\sqrt{2} + 1)^2$

so that $-1 \le x < 1$ for $\sqrt{2}/2 \le g < \sqrt{2}$. Substituting

$$\ln (g) = 2 (z + z^3/3 + z^5/5 + \dots)$$

where $z = x/v = \frac{g-1}{g+1}$.

The approximation used is G(z) for $\ln(g)/z$ in the range $\sqrt{2}/2 \le g < \sqrt{2}$.

Then for both extended and standard precision,

 $z = \frac{g-1}{g+1}$ $\sqrt{2/2} = .7071067811865$ $\ln (2) = .6931471805599$

Thus, the required calculation is

 $\ln(x) = j \cdot \ln(2) + zG(z)$

Extended Precision

$$G(z) = b_0^{+} b_2^{2} z^{2} + b_4^{2} z^{4} + b_6^{2} z^{6} + b_8^{2} z^{8}$$

where

$$b_0 = 2.0$$

 $b_2 = .666666564181$
 $b_4 = .400018840613$
 $b_6 = .28453572660$
 $b_8 = .125$

Standard Precision

$$G(z) = b_0^{+} b_2^{-} z^2 + b_4^{-} z^4 + b_6^{-} z^6$$

where

$$b_0 = 2.0$$

 $b_2 = .666664413786$
 $b_4 = .4019234697$
 $b_6 = .25$

EXPONENTIAL

Polynomial Approximation

To find e^x, the following identity is used.

To reduce the range, we let

 $x \log_2 e = n + d + z$

where

- n is the integral portion of the real number,
- d is a discrete fraction (1/8, 3/8, 5/8, or 7/8) of the real number, and
- z is the remainder which is in the range $-1/8 \le 2 \le 1/8$.

Thus,

 $\mathbf{e}^{\mathbf{x}} = 2^{\mathbf{n}} \mathbf{x} 2^{\mathbf{d}} \mathbf{x} 2^{\mathbf{z}}$

and it is necessary to only approximate 2^z for $-1/8 \le z \le 1/8$ by using the polynomial F(z).

Extended Precision

$$F(z) = a_0 + a_1 z + a_2 z^2 + a_3 z^3 + a_4 z^4 + a_5 z^5$$

where

$$a_0 = 1.0$$

 $a_1 = .69314718057$
 $a_2 = .24022648580$
 $a_3 = .055504105406$
 $a_4 = .0096217398747$
 $a_5 = .0013337729375$

$$a_0 = 1.0$$

$$a_1 = .693147079$$

$$a_2 = .240226486$$

$$a_3 = .0555301557$$

$$a_4 = .00962173985$$

HYPERBOLIC TANGENT

Tanh (x) =
$$\frac{e^{2x}-1}{e^{2x}+1}$$

for

x ≥ 32	Tanh(x) = 1
x ≤ -32	Tanh (x) = -1

REAL BASE TO REAL EXPONENT

$$A = e^{\ln A}$$

therefore:

$$A^{B} = (e^{\ln A})^{B} = e^{B\ln A}$$

<u>Standard Precision</u>

 $F(z) = a_0 + a_1 z + a_2 z^2 + a_3 z^3 + a_4 z^4$

Selective Dump Subroutines

The IBM 1130 Subroutine Library and the System Library include three dump subroutines: Dump Selected Data on the Console Printer, Dump Selected Data on the 1132 Printer, and Dump Status Area. These subroutines allow the user to dump selected portions of core storage during the execution of a user's program.

Dump Selected Data on Console Printer or 1132 Printer

Two subroutines are available to select an area of core storage and dump it on the Console Printer or the 1132 Printer. Each of these subroutines has two entry points, one for hexadecimal output and one for decimal output. The entry points for the various configurations are shown below:

- DMTX0 Dump on Console Printer in hexadecimal format, using the WRTY0 subroutine
- DMTD0 Dump on Console Printer in decimal format, using the WRTY0 subroutine
- DMPX1 Dump on 1132 Printer in hexadecimal format, using the PRNT1 subroutine
- DMPD1 Dump on 1132 Printer in decimal format, using the PRNT1 subroutine

Calling Sequence

The calling sequence for any of the above functions is as follows:

CALL	ENTRY	PO INT
DC	START	
DC	EN D	

START and END represent the starting and ending addresses of the portion of core storage to be dumped. A starting address greater than the ending address results in the error message, ERROR IN ADDRESS, and a return to the calling program.

Format

Before the actual dump appears on the selected output levice, the user is given

one line of status information. This line indicates the status of the Overflow and Carry indicators (ON or OFF), the contents of the Accumulator and Extension, and the contents of the three index registers. The index register contents are given in both hexadecimal and decimal form, regardless of which type of output was requested. The format of the status information is shown below:

OFF ON HHHH (±DDDDD) HHHH (±DDDDD) Overflow Carry Accumulator Extension

HHHH (±DDDDD) HHHH (±DDDDD) HHHH (±DDDDD) Index Reg 1 Index Reg 2 Index Reg 3

All other data is dumped eight words to a line; the address of the first word in each line is printed to the left of the line. Hexadecimal data is printed four characters per word; decimal data is printed five digits per word, preceded by a plus or minus sign.

Page numbers are not printed for either subroutine. However, the 1132 Printer subroutine does provide for automatic page overflow upon the sensing of a channel 12 punch in the carriage tape.

Dump Status Area

This subroutine provides a relatively easy and efficient means of dumping the first 80 words of core storage. These words contain status information relating to index registers, interrupt addresses, etc., which may be required frequently during the testing of a program. It may also be desirable to dump these words before loading because pressing PROGRAM LOAD destroys the data in the first 80 words of core storage.

The Dump Status Area subroutine is called via the following statement:

CALL DMP80

The Console Printer prints the first 80 words of core storage in hexadecimal form; the printing format provides spacing between words. After typing the last word, the subroutine returns control to the calling program.

Special Monitor Subroutines

The DM2 System Library contains a group of subroutines that perform various system utility functions. These subroutines, with the exception of SYSUP which can be called by the user, are intended for system use only. Under normal circumstances, they should not be deleted from the System Library.

The subroutines in the group are:

FLIPR - LOCAL/SOCAL overlay subroutine
RDREC - Read *ID record
CALPR - Call system print
FSLEN - Fetch phase IDs
FSYSU - Fetch system subroutine (FSYSU is an alternate entry point to FSLEN)
SYSUP - DCOM update

FLIPR (LOCAL/SOCAL OVERLAY)

The System Library contains a flipper subroutine (FLIPR) which is used to call LOCAL (load on call) and SOCAL (system load on call) subroutines into core storage. FLIPR is used with DISKZ, DISK1, or DISKN.

FLIPR passes the total word count to DISKZ, DISK1, or DISKN to fetch the LOCAL. When a LOCAL subroutine is called, control is passed to the flipper, which reads the LOCAL into core storage if it is not already in core and transfers control to it. All LOCALs in a given core load are executed from the same core storage locations; each LOCAL overlays the previous one. FLIPR fetches SOCALs in the same manner as LOCALs.

RDREC (READ *ID RECORD)

This subroutine is called by Disk Maintenance Programs to read the *ID (disk label) record. This subroutine is intended for system use only.

CALPR (CALL SYSTEM PRINT)

This subroutine calls FSLEN to bring the system print subroutine into core storage

for the purpose of printing one or more lines on the principal printer. This subroutine is intended for system use only.

FSLEN (FETCH PHASE IDS AND FETCH SYSTEM SUBROUTINE)

This subroutine has two entry points. They are FSLEN and FSYSU.

FSLEN (Fetch Phase IDs from SLET)

This entry point obtains the requested phase ID headers from SLET.

FSYSU (Fetch System Subroutine)

Fetches the requested system subroutine into core storage.

This subroutine is intended for system use only.

SYSUP (DCOM UPDATE)

Whenever a core load requires changing disk cartridges during the job, SYSUP must be called to update DCOM on the master cartridge (logical drive 0) with the IDs and DCOM information from all satellite cartridges mounted on the system. The cartridges are specified in the list or array in the SYSUP calling sequence. The list or array must be exactly five words long or be ended by a zero (not both).

The Assembler language calling sequence for SYSUP is:

Label		Operation		F	т		· · ·		Operands & Re	marks	
21 25		<i>11 w</i>		32	33		35 40	 45	50	**	м
		CALL					SIXS.U.P.	 CALL	N.C.O.M	UPDA	T.E.
	L	D.C	_	L	L	L	LILST	 <u> </u>			
		••••		L	L			 	1 - 1 - 1 - 1 - 1		
	L	• • •			L			 			
		• • • •	L		L	L	Ature	 	1-1-1-1-1-1		
LIST.		ט בי ט					ann	 			
		D.C					6	 1.1.1.1			
		י גאמ					G	 1			
		DC I					d	 			
		DC .		Γ	Γ		e	 1.1.1.1			1
							1	 1.1.1.1			

where

LIST is the address of the table of requested cartridge IDs,

a is the ID of the master cartridge on the system,

b is the ID of the first satellite cartridge on the system,

c is the ID of the second satellite cartridge on the system,

d is the ID of the third satellite cartridge on the system,

e is the ID of the fourth satellite cartridge on the system.

If a is 0, the master cartridge remains unchanged.

The FORTRAN calling sequence for SYSUP is:

 $\frac{1}{1} + \frac{1}{2} + \frac{1}{2} + \frac{1}{2} + \frac{1}{1} + \frac{1}$

where

a is the name of the last item in an array containing the IDs of the satellite cartridges on the system. The last entry in the array may be 0, in which case the master cartridge remains unchanged.

For example:

CALL SYSUP (K (5))

The array is stored in reverse order.

K (5)	DC
K (4)	DC
K (3)	DC
K (2)	DC
K (1)	DC

Thus K (5) is the entry for logical 0, the master cartridge.

SYSUP messages are listed in the publication IBM 1130 Disk Monitor System, Version 2, Programmer's and Operator's Guide. SYSUP execution is terminated if an error printout occurs. The IBM 1130 DM2 System Library mainline programs provide the user with the ability to perform disk maintenance and paper tape utility functions by requesting execution of the appropriate program directly through the job stream.

The calling sequences for the System Library mainline programs are listed below. The operating procedures and error messages are contained in the <u>IBM 1130 Disk Monitor</u> <u>System, Version 2, Programmer's and</u> <u>Operator's Guide</u>.

Disk Maintenance Programs

The disk maintenance programs are mainline programs and subroutines that are a part of the System Library and that initialize and modify disk cartridge IDs, addresses, and tables required by the DM2 system. Normally, they should never be deleted from the System Library.

The disk maintenance programs are:

IDENT -	Print Cartridge ID
DISC	Satellite Disk Initialization ¹
ID -	Change Cartridge ID
COPY	Disk Copy
ACRWS -	Write Sector Addresses in Working
	Storage
DFCNV -	Disk Data File Conversion
DLCIB -	Delete CIB
DSLET -	Dump System Location Equivalence
	Table
MODIF -	System Maintenance Program

MODSF - Library Maintenance Program

'All new cartridges are initialized using the standalone program DCIP (see <u>IBM 1130</u> <u>Disk Monitor System, Version 2,</u> <u>Programmer's and Operator's Guide</u>).

<u>IDENT (Print Cartridge ID)</u>

This program prints the ID and physical drive number of each cartridge mounted on the system.

IDENT prints all cartridge IDs regardless of validity (JOB card processing only recognizes valid IDs). The calling sequence for IDENT is:

1	1	3	4	\$		1	7	i	,	10	11	12	13	14	15	16	17	. Ц	: 19	20	21	77	23	24	29	26	27	28	29	30	31	22	33
/	Ū		Y	E	6	h		7	0	£	N	17	;				1	,						1								1	
•																																	
				_	L_	1					1						1		1		ь,	_						<u> </u>		•	L.,	L.,	1 1

DISC (Satellite Disk Initialization)

This program reinitializes up to four satellite cartridges -- all but the master cartridge. It writes the sector addresses, defective cylinder addresses, a cartridge ID, a LET, a DCOM, and a CIB on each cartridge being reinitialized.

DISC overrides all cartridge IDS specified on the JOB card except the master cartridge ID.

The calling sequence for DISC is:

1 2 3	4 5	6 7	1	9	10	11	12	13	14_1	15. 1	6 17	18	19 3	10 2	1 22	23	24	25	26	27	28	29	30	31	32	33	34	15 34
11	V.F	ል	.n		۰.	r																						
<u> </u>	A L-	Ψ1	-	/	5	-	_	_	-	_		1	Ц	_	 _	1_							. 1	. 1				_
		_			-	-			-	_			_	_						_	-	-			-	-	-	
₩, <i>I</i> ,0	FII	Dı1	1.1	7	J.	Dı	L.	• P	Fu	Ζıć	02	1.	ı7ı.	Z, I	0,2		۰.	•	•	F	Z.	D,	ħ,	• 1	T	Z	D	
<u>*,<i>I</i>,0</u>	FII	D11	وا	T	J	D	1	3 ¹	Fil	ζı	02	<u>ډا</u>	<u>, 7</u> ,	<i>I</i> /	0,2	J.	•	•	•	F	Z.	<u>D</u>	1 1	ş 1	T	I	D	

where

FID1 through FIDn are the IDs currently on the satellite cartridges to be reinitialized (identified by IDENT or a JOB record),

TID1 through TIDn are the IDs to be written on the satellite cartridges by this program. A valid cartridge ID is a number between /0001 and /7FFF.

ID (Change Cartridge ID)

This program changes the ID on up to four satellite cartridges.

		-			_		_	-	_	_	_								_	_	_	_	_	_			_		_	_	
1 2 3	4 5	.67		. 9	10	11	12	13	14	15	16	17	- 12	1 19	20	- 21	- 22	23	- 24	25	26	27	28	29	- 30	31	- 32	33	- 34	35	36
			_		_			_		-				_								_	_		-	-			-	-	
				-																											
17.7.	. Y.F.	<i>(</i> 1																													
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													_											_	_	_	_	_	_	-	_
		~ 4		-		-					-										-	-					-				
HDLL		\boldsymbol{u}			. / 1	Lh	21	- 1		1	L	12	۰.	. /	14	14	ΑZ					.,						. I	. 6		
					_	_		-		-	-		- 42		1.00	-			-	_					- 44		<u>.</u>			- 4	_
																									-	-					
					4 8																		•								

FID1 through FIDn are the IDs currently on the satellite cartridges being changed (these IDs must be in the same logical order as the entries on the JOB card),

TID1 through TIDn are the new IDs to be written on the selected satellite cartridges.

COPY (Disk Copy)

This program copies the contents (except the defective cylinder table and the cartridge ID) of one cartridge onto another. The copy ID (word 5 of sector 0) is incremented by one prior to being written on the new cartridge.

The calling sequence for COPY is:

L.	,	,				٨	,			10	11	19	11	14	15	14	. 17	- M	1 19	20	21	22	23	24	25	24	77	20	90	20	11	71	11	14	-	14
<u> </u>	·•	~	-		<u> </u>	•	ć		-								,									-			47	30	31	*	-33		-45	30
17	۰,	/	•			-		^	~		194	•																								
v.	1	1	14		- 11	2		5	υ	2	17	1	_	L	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1 1
Г .				_																			_	_			_		_		_	_	_		_	_
11	. 1	r./	ער	- 1	۲.,	n	1		.7	. 7	1	11	·	.F	. 7	•1	29) .	.7	- 7	.1	12	P				. 2			۰.		. 7	. 7	. /	۰.	
12	14	14				-		_ و ا		14	1.4	14	4	1			10	. و	12	14	14	41.5	цg	1.	1-		1	14	114	44	19			14	14	ш
L								-					-					-																		
L		1	1		1	- 1		1	t I		•	۱.	•				1	1		4			1	4	1					4 É	1					
	-		-		-		-	_	· ·		-	•			-	-		-	-	_		-	-				-	•							-	_

where

FID1 through FIDn are the IDs of the cartridges to be copied.

TID1 through TIDn are the IDs of the cartridge onto which the copies are to be made.

If multiple copies are to be made from a single master, FID1 through FIDn will all contain the same ID.

<u>ADRWS (Write Sector Addresses in Working</u> <u>Storage)</u>

This program, linked to from DUP on detection of the DUP control record DWADR, writes sector addresses on all sectors of Working Storage on the disk cartridge specified by the DWADR control record (see DUP in the <u>IBM 1130 Disk Monitor System</u>, Version 2, Programmer's and Operator's Guide).

DFCNV (Disk Data File Conversion)

This program converts 1130 FORTRAN or Commercial Subroutine Package disk data files to disk files acceptable to 1130 RPG. The calling sequence for DFCNV is:

2	4	6	8	10	12	14	16	18	20	22	24	26	28	30	32	34
77	X	Q	D	FCI	W.		1.1		1.							
										11						

DLCIB (Delete Core Image Buffer)

This program deletes the CIB from a nonsystem cartridge. If a user area is defined, the user area is moved two cylinders closer to cylinder 0. The new addresses of disk areas moved as the result of the deletion of the CIB are reflected in DCOM on the master cartridge, on the nonsystem cartridge from which the CIB is deleted, and in COMMA.

The calling sequence for ELCIE is:

1 2 3 4 5 6 7	8 9 19 11	12 13 14 19 16	17 18 19 20 21 22 23	34 25 34 27 28 29 39 31 22 39 34 58
11. YED	DI C.T	R		
	10000 100 100			
# <i>.I.</i> 06,4,87	<u> </u>			

where

CART is the ID of the non-system cartridge from which the CIB is to be deleted.

DSLET (Dump System Location Equivalence Table)

This program dumps the contents of SLFT on the principal printer. Each entry printed consists of a symbolic name, a phase ID, a core address, a word count, and a disk sector address. A SLFT dump is listed in the publication <u>IBM 1130 Disk Monitor</u> <u>System, Version 2, Programmer's and</u> Operator's Guide.

The calling sequence for LSLFT is:

Ŀ	2	3	4	5	4	7	1	,	10	11	12	13	14	19	16	17	10	19	20	21	22	23	24	25	26	27	28	29	30	. 31	32	33	34	35	7
L.	1	1	v	-	· ~			~	• •			•																							-
Ľ	1	1	LĀ	" C	ιų.	_	<u>_</u>		12	1C	1	١	L.,	1	1	1	1	1	1	1	1	1	1	1	L	1	1	1	L	L.	1	L.	1	ı.	ı.
1					•																												_	-	-
L	1	1	1				1	1	1	4	1	L	1	1	١	L	1	1	1	1	1	1	1	1	1	L.,	1	1	1	1	1	1	1	1	I.

MODIF (System Maintenance Program)

Included in the DM2 System Library is a system maintenance program, MOCIF, that provides the user with the ability to update the Monitor system on the master cartridge. This program makes changes to the version and modification level word in DCOM, the Supervisor, DUP, FORTRAN Compiler, Assembler, and/or System Library. A card deck or paper tape containing corrections to update the Monitor system to the latest version and modification level is supplied by IBM. Every modification must be run to update the version and modification level, even if the affected program has been deleted from the system. The calling sequence for MODIF is:

1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25 26 27 28 29	30 31 32 33 34 35 1
1/1 XEQ MODIF.	

and shake when we have been a set to be the set of the

a series de la companya de la companya de la companya de la companya de la companya de la companya de la compa A series de la companya de la companya de la companya de la companya de la companya de la companya de la company A series de la companya de la companya de la companya de la companya de la companya de la companya de la company . المركز بالمصحفة المساجل بالمطلقة المركزة المركز المركز المركز المركز المركز المركز المركز المركز المركز المركز

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MODSF (Library Maintenance Program)

The purpose of MODSF is to update a library program written in the disk system format (DSF) and located in the User Area of disk storage. (To modify or replace a <u>system</u> program, see "MODIF (System Maintenance Program) " described earlier in this section.)

A program is updated by either replacing existing code, inserting additional code at the end of the program, or both. Existing coding is replaced as the program resides in the User Area. Several programs may be updated in a MODSF run, but only the last program in a MODSF run may have code added to it. When additional code is inserted, MODSF moves the program to Working Storage and inserts it there and ends its run by invoking DUP. To move the updated program back to the Users Areas, the user must provide the necessary *DELETE and *STORE records.

To update a program with MODSF, the user must prepare a patch control record, one or more patch data records, and a patch terminator record.

The calling sequence for MODSF is:

1.1	1	4		_2		_		£٦	Ū	Ď.	11	14	И	Ľ	L Û	7_1	1	19_1	9	21	22	23	M	2	1 2	1 2	7 8	1 2	•	¥.	31	22	11	34	29
1.1		¥. 1	E.C.			./	•	n .	c.	c .																									
K-K-		2.12	-				~	.	·			-	-	-			-	-		_			-	•	-		-	-			-	-	_		
		1	1	1	1	1	1	1	1	_		_	1	1	L	1	_	1	1	_		L.,	L	L	L	1		1	1	_		1		L.,	1
1.																												,							
	_	-	-		-		-	-	-	_	_	-	•	-	-	-	-	-		-		-		-		-		-	_		-	-	-		

PAPER TAPE UTILITY (PTUTL)

This program accepts input from the keyboard or the 1134 paper tape reader and provides output on the console printer and/or the 1055 paper tape punch.

PTUTL allows changes and/or additions to FORTRAN and Assembler language source records as well as monitor control records.

The calling sequence for PTUTL is:

t	2	3	4	1	6	7.1		10	1_11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31	32	33	34	35 ;
1		1	14	_	~																												
2	Δ		1Å1	E,	64	/	17	14	47	4			L		1	1	L	۱	ட		1	L.	L	1	ı I	1	1	1.1		ı –	•		11
																										_			_				
	. 1		1 1			1	1	1	1	E I			L . I			ı I	ı.	ı.	1	ı I			1	ı		1	1		t I				

The section on Writing ISSs and ILSs for the DM2 system will be found in the <u>IBM</u> <u>1130 Disk Monitor System, Version 2,</u> <u>Programmer's and Operator's Guide.</u>

Interrupt Service Subroutines

The following rules must be adhered to when writing an ISS:

- 1. Precede the ISS statement with an LIBR statement if the subroutine is to be called by a LIBF rather than a CALL.
- 2. Precede the subroutine with an EPR (extended) or an SPR (standard) statement if precision specification is necessary.
- 3. Precede the subroutine with one ISS statement defining the entry point (one only), the ISS number, and the ILS subroutines required. The device interrupt level assignments and the ISS numbers used in the IBM-provided ISS and ILS subroutines are shown in Figure 22.
- 4. The entry points of an ISS are defined by the related ILS. This must be taken into consideration when a user-written ISS is used with an IBM-supplied ILS. The ILS executes a BSI to the ISS at the ISS entry point plus n (see Figure 22). The ISS must return to the ILS via a <u>BSC</u> instruction (not a BOSC).
- When assembling the ISS on the monitor system an *LEVEL n control card must be included for each interrupt level associated with the device.

Interrupt Level Subroutines

An ILS is included in a program only if requested by a loaded ISS. The following are rules for writing an ILS:

- 1. Precede the subroutine with an ILS statement.
- Precede all instructions by an ISS Branch Table and include one word per ILSW bit used. If the ILSW will not be scanned, (i.e., a single ISS subroutine to handle all interrupts on the level), then a one-word table is sufficient. The minimum table size is one word. Table words must be nonzero.

ILSW Bit 15 word ILSW Bit 14 word ILSW Bit 0 word

The ISS Branch Table identifies both the ISS subroutine and the point within the ISS which should be entered for each bit used in the ILSW. The actual linkage is generated by the Relocating Loader or Core Image Converter. Basic to this generation is the ISS number implied by bits 8-15 of the branch table word and specified in the ISS statement. This number identifies a core location in which the loader or converter has stored the address of the called entry point in the ISS. This entry point address is incremented by the value in bits 0-7 of the branch table word, producing the branch linkage. The loader or converter replaces the ISS branch

ISS Number	Device	Device Interrupt Level Assignments	n
r 1	1442 Carl Read Punch	0,4	+4,+7
2	Keyboard/Console Printer	4	+4
j 3	1134/1055 Paper Tape Reader/Punch	1 4	1 +4
j 4	Single Disk Storage	2	+4
j 6	1132 Printer	1 1	+4
7	1627 Plotter	3	+4

Figure 22. C/PT System ISS/ILS Correspondence

table word with the generated branch linkage.

During execution, the ISS Branch Table contains core addresses. It may be used with an indirect BSI instruction to reach the ISS corresponding to that ILSW bit position. The ILSW bit that is ON can be determined by the execution of a SLCA instruction. At the completion of this instruction, the index register specified contains a relative value equivalent to the bit position in the ISS branch table. An indirect, indexed BSI may then be used to reach the appropriate ILS.

Each word in the ISS branch table has the following format:

Bits 0-7: Increment added to the entry point named in the ISS statement to obtain the interrupt entry point in the ISS for this ILSW bit. (In IBM-written ISS subroutines, this increment is +4 for the primary interrupt level and +7 for the secondary interrupt level.)

Bits 8-15: ISS number +51 for the ISS subroutine for this ILSW bit. This

address should match word 13 of the compressed ISS header card.

- 3. The ILS entry point must immediately follow the ISS branch address table and must be a zero. The first zero word in the program is the end of the branch table and is also the entry point of the ILS. (The table must contain at least one entry.) The interrupt results in a ESI to the ILS entry point.
- To clear the level, a user-written ILS, used with an IEM-supplied ISS, should exit via the return linkage with a BOSC instruction.

ILSs supplied by IBM in the Card/Paper Tape System, except ILS01, pass word 2 of the Sense Device IOCC to all ISSs. The ISSs in the Card/Paper Tape System require that this word be passed in the Accumulator. Observe this convention when writing ILSs, and when writing ISSs to be used with IBM-supplied ILSs.

20010571								9100
26010511				0143				8100
16010511	#571 35N35 01 3301	00504		20	SNJS	0050	0	4100
05010571		0		20		0000	ŏ	9100
6010571	FOR SAVING ACC AND EXT.	Ś	Э	558	٥v	2000	-	V100
8+010571					*			
24010571	GO TO RESTORE AND EXIT	TAT2		XQM		70F7	0	£100
94010571	PROCESS 1132 INTERRUPT	59	1	158	۸IP	00000844	10	1100
5010571					*			
99010S71	TURN OFF LEVEL AND EXIT	10571	1	3208		400002	to	J0 00
E#010571					*	•		
24010211		*-*	רז	хaл	18X	00000059	00	0000
14010571		QA.		007		7080	0	2000
00010511	2401239	0		507	TAT2	2000	0	8000
62010571					*			
8010571	TAUAABTNI ABTGADA 22300A9	STAD	I	128		10000844	10	6000
75010211	PROCESS 1132 INTERRUPT	Z+ . GIV	٦	228		11008524	10	2000
950102.11	MS71 35N35	1-SN35		01X		7080	0	9000
62010271			_		*			_
+E010571	I NDEX I	1+18x	I.	XT2		8069	0	5000
11201033	SUTAT2	TAT2		ST2		2806	0	4000
SE0102-1	SAVE ACC AND EXTENSION	DA		ars		0180	0	£000
10010571	COOD NOITADOJ AIV				*			
02010511	IZA BAANGAAH YA GEGETUE	0		00	10511	0000	0	2000
62010571		8540/		00	5140	BEA0	0	1000
12010511		OF AON	10	50	60	05.40	0	0000
22010511			.0	211	-			
62010511	*****	********			-			
52010571	***************************************	******			*****			
62010571		110 +						
22010511		119 *						
12010571	* • • • • • • • • • • • • • • • • • • •	μ μ μ μ μ μ μ μ μ μ μ μ μ μ				*		
07010571	* NO 384 8319404 SNOTTADINUM	# COM						
61010571	* BHT ONA RETUIRS SELL	3HT *			*			
81010571	BOUTINE FOR LEVEL 1.	18AS #			*			
11010571	* IS THE INTERRUPT LEVEL *	1HL *			*			
91010571	# 10571 0E11	******			¥			
51010571	*			- SE	LON #			
¥1010S71	*	318ASU39 -	- 52	atue i s	ITTA #			
61010571	* IND	IN - 2439A	ж	DM/S3T	18AT *			
21010211	*	- NONE	905	-EGE	*			
11010571	INDIRECT THROUGH . ILSOI . *	2208 - JAN	NGORN	4 - ST	* EXI.			
01010571	* 3NO	IN - SENIT	ROL	JAN93	* EXT			
60010571	*	E	NON	- 10e	1100 *			
80010571	T		JNO	лод - Л - ТС	NGNI #			
200105-11		6 N	111		1 1 7 #			
20010211	ENTERED BY HADDWADF BCI +	-10211 - 3	271C					
20010511		TIME FOO	거리님!	13 13/	131 1			
					7790 1.21 👬			
C0010571		73437 30M	dun.		- 216			
£0010571	* 0	NGE FEAEF	CHN	- 501	.V15 *			
1001003 1001003	* 0	NOE FEAET	H)	- sni E - 1	* 21V			

:

Sample ISS (C/PT)

			***	HDNG		LIBF CARD	00	CRD00001
		02050120	1120	LIBR	01	CARDO	0 4	CRD00002
0000		03024130	******	133	****	*******	***********	CR000004
			*	THIS	113	O SUBROUT	INE OPERATES THE 1442 CARD *	CR000005
			*	READ	ERP	UNCH. IT	INITIATES REQUESTED OPERA- *	CRD00006
			*	TION	S. P	ROCESSES	ANY COLUMN OR OPERATION *	CRD00007
			*	COMP	LETE	INTERRU	TS. AND AUTOMATICALLY #	CRD00008
			*	INIT	IATE	S ERROR R	RECOVERY PROCEDURES.	CR000009
			*		TIEV	ING FEAT	PF - NO FROOR PARAMETER +	CRD00011
			*	1000	1161 ****			CR000012
			*			LOADER DE	FINED LOCATIONS +	CRD00013
			*****	****	****	*******	· * * * * * * * * * * * * * * * * * * *	CRD00014
0000 0).	695E	CARDO	STX	1	CA30+1	LIBF ENTRANCE (+0)	CRD00015
0001 0	00	65800000	LINK	LDX	Ι1	0	LOADER STORES TV ADDR (+2)	CRD00016
0003 0	2	7006		MDX		CAIO		CRD00017
0004 0	2	0000	INT1	DC		0	COLUMN INTERRUPT (+4)	CR000018
0005 0	1	4C00000A	INTO	BSC DC	-	0	OP CMPLTE INTERPURT (+7)	CRD00019
0007 0	ונ	4C00009F	TINE	BSC	L	NT10		CRD00021
		400000	*****	++++	****	******	*******************	CRD00022
			*			LIBE PROC	CESSING *	CRD00023
			*****	****	***;	*******	*************************	CRD00024
			¥	THIS	POF	TION STOP	RES CALLING SEQUENCE . INFO *	CRD00025
			*	AND	CHEC	KS THE DE	EVICE STATUS BEFORE ANY 1/0 *	CRD00026
			*	NOT		N 15 INII	ALCEC AN EDDOD EVIT TO	CR000027
			*		TION	J 41. IF 7	THE OPERATION WILL CAUSE *	CR000029
			*	INTE	RRUF	TS. THE F	ROUTINE IS SET BUSY AND THE *	CRD00030
			*	LOCS	COL	INTER IS	INCREMENTED TO INDICATE *	CRD00031
			*	INTE	RRUP	TIS PEND	DING• *	CRD00032
			****	****	***1	*******	************************	CRD00033
000A 0	C	C07B	CAIO	STO		TEMP	SAVE STATUS	CRD00034
0008	c	2856		STS	-	CA32		CRD00035
0000 0	2	6A54		STX	2	CA31+1	NI- 1000 05 CM 11	CRD00036
0000 0	2	1800		LD CDA	1	12	XIª ADDR OF CALL+I	CR000037
DODE C		40200015		BSC	1	CA14.7	NO	CRD00039
0011 0	5	C078		LD	-	BUSY	YES. IS ROUTINE BUSY	CRD00040
0012 0	С	4818		BSC		+-		CRD00041
0013 0	D	7101		MDX	1	+1	NO. EXIT TO CALL+3	CRD00042
0014 0	D	7046		MDX		CA28	YES. EXIT TO CALL+2	CRD00043
0015 (0	9077	CA14	s		D0004	IS FUNCTION LEGAL	CRD00044
0016 0	01	4C300070		BSC	Ļ	CA40.Z-	NO . ERROR	CRD00045
0018	0	8074		A		H7003		CR000046
0019 0	0	DOOR		STO	-	CA20		CR000047
0018	0	0007		STO		CAIR		
0010	ŏ	C06D	CA15	LD		BUSY	IS ROUTINE BUSY	CRD00050
OUID	01	4C20001C		BSC	L	CAISIZ	YES. WAIT TIL NOT	CRD00051
OC1F	0	0868	CA17	X10		SENSE-1	IS DEVICE READY	CRD00052
0020	01	4C040072		8SC	L	CA42.E	NO + ERROR	CRD00053
0022	0	C066		LD		SENSE	SETUP CONTROL LOCC	CRD00054
0023	0	9075	CAIB	5		SETUP		CRD00055
0024	0	7000	C 4 20	SIO				CRD00056
0025	o o	7003	CAZO	MDX		CA20+1	■ GFT	CPD00057
0027 0	D	7030		MDX		CA36	= 021 = PUT	CRD00059
0028	D	701B		MDX		CA25	= FEED	CRD00060
0029	0	702B		MDX		CA26	= STK	CRD00061
002A 0	D	9072	CA21	S		SETUP+4	GET FUNCTION	CRD00062
0028	5	D059	C	STO		COLM+1	SET UP READ 1/0	CRD00063
0020 (10	40080070	CA218	LU	11	1	= FODAD 15 7500 00 MEC	
0020 0	.	40080070		65C	L.	CA40++	= ERROR IF ZERO OR NEG	CRD00065
0031 0	Ď	8054		Δ Δ		00001	SAVE WORD COUNT +1	CR000067
0032	0	D061		STO		COUNT	BECAUSE DECREMENT IS	CRD00068
0033 (D	D063		STO		RSTRT	BEFORE COLUMN READ	CRD00069
0034	0	905B		s		D0081		CRD00070
0035	01	4C300070		BSC	L	CA40+Z-	= ERROR IF OVER +81	CR000071
0037	0.	C101		LD	1	1		CRD00072
0038	0	0004	(STO		CA23+1		CR000073
0039	00	C050	CAZZ		62	00001		CR000074
0030	00	0000000	CA23	STO	L2	0	STORE +1 IN DATA AREA	CRD00076
003E	0	72FF	_	MDX	2	-1	F NOT READ INDIC FOR	CRD00077
003F (0	70FC		MDX		CA23	SPEED CONVRT SBRT	CR000078
0040 0	0	C101	CA24	LD	1	1	SAVE DATA ADDRESS	CRD00079
0041 0	0	D042		STO		COLM		CRD00080
0042 0	0	7101		SID	,	H218141	SET YI TO SKID 2ND DADAM	CBD00081
0044	0	0843	C425	XIO	1	SENSE-1	SET AT TO SALE ZIND PARAM	CRDAAAA
0045	Ó	1003		SLA		3	IS LAST CARD IND ON	CRD00084
0046	01	40100050		BSC	L	CA258	NO	CRD00085
0048	D	CODC		LD		CA20	IS FUNCTION GET OR FEED	CRD00086
0049 0	01	40040050		BSC	L	CA25B+E	NO	CRD00087
0048 (D	1008		SLA		8	IS FUNCTION GET	CRD00088

Writing ISSs and ILSs (C/PT System) 115

004C	0	4808		BSC		+		CR00089
004D	0	71FF		MDX	1	-1	YES. SET XR1 = LIBF+1	CRD00090
004E	0	0838		x10	•	FFFD-1	FUECT CAPD	CPD00091
0045	0	7020		MDY		CA43	LUCCT CARD	CR000091
004	ň	74010033	CASER	MDM		CA43		CRD00092
0050	00	14010032	CA250	MUX	L	20++1	INCREMENT TUCS COUNTER	CRD00093
0052	0	1000		NOP				CRD00094
0053	0	C038		LD		D0001		CRD00095
0054	0	D035		STO		BUSY	SET ROUTINE BUSY	CRD00096
0055	0	CO3F	CA26	LD		ERROR		CRD00097
0056	01	4C20005A		BSC	L	CA27.Z		CRD00098
0058	0	0820		XIO		INIT-1	INITIATE 1/0	CRD00099
0050	ñ	7001		MDY		CA28		CDD00100
0055	Ň	0805	CA 37	~		CAZO		CROOOTOL
0054	~	0020	CMET	ATO		FEED-I		CROOOIDI
0058	0	101	CA28	MUX	1	+1		CRD00102
005C	0	C029		LD		TEMP		CRD00103
005D	0	6906	CA29	STX	1	CA34+1	SET EXIT TO SKIP IST PARAM	CR000104
005E	00	65000000	CA30	LDX	LI	0	RESTORE STATUS	CRD00105
0060	00	66000000	CA31	LDX	L2	0		CR000106
5000	0	2000	CA32	I DS		ō		CPD00107
0063	00	40000000	CA34			õ	EVIT	CDD00101
0005	~	0038	CAR	6 C	-	CETUDIE		CREDOULOE
0065	č	9036	CASO	5		SETUP+D		CRDOOLOS
0000	0	DOTE		510		COLM+1	SETUP PUNCH 170	CROOOTIC
0067	00	000001		LD	11	1		CRD00111
0069	01	4080070		BSC	L	CA40++	* ERROR IF ZERO OR NEG	CRD00112
006B	0	D028		STO		COUNT		CRD00113
006C	0	D02A		STO		RSTRT	SAVE WORD COUNT	CRD00114
0060	0	9021		s		D0080	DO NOT PUNCH OVER 80 COL	CRD00115
006E	0	4808		BSC		+		CR000116
006F	0	7000		MDX-		CA24		CR000117
0070	0	C021	CA40	LD		H1001	FRROR CODE - ILLEGAL CALL	CRD00118
0071	õ	7008		MDX		CAAA		CR000119
0071	~	1801	C	504			TE DEVICE BUEN	CDD00130
0072	<u>.</u> .	1801	CA42	SRA			IS DEVICE BUST	CR000120
0073	01	4004001		BSC	L	CALTHE	YES WALL HIL NUT	CRDOOLSI
0075	0	1003		SLA		3	IS DOW ERROR INDIC ON	CR000122
0076	01	4C10007C		BSC	L	CA43+-	NO	CRD00123
0078	0	COAC		LD		CA20	YES, IS FUNCT GET/FEED	CRD00124
00 79	01	4C04007C		BSC	L	CA43+E	NO	CR000125
007B	0	D019		STO		ERROR	YES. INDIC SKIP 1ST CD	CRD00126
0070	0	C014	CA43	LD		H1000	FRROR CODE - DVCE NOT RDY	CR000127
0070	ō	7155	CAAA	MDX	1	+1		CR000128
0075	00	6000029	0.144	STY		40	STORE CALL ADDR IN AD	CPD00129
0072	00	6000028		314		40	STORE CALL ADDR IN 40	CR000129
0080	0	6129		LUX	1	41	SET EXTLEPOR 41	CR000130
0081	0	7008		MDX		CA29		CR000131
			****	****	***1	**********	***********************	CRD00132
			*			CONSTANTS	*	CRD00133
			*****	****	****	**********	*********************	CRD00134
0082		0000		BSS	E	0		CRD00135
0082	1	0082	ADDR	DC		CHAR-1	ADDR TO REPLACE OZP AREA	CRD00136
0083	ò	0000	CHAR	DC DC		0	TEMPOPARY REGISTER	CR000137
0083	š	0000	COLM	50		0		CR000139
0084	0	0000	COLM	UC .		0	TUCC FOR COLUMN 170 E	CRUUUISE
0085	U U	0000		DC		U	0	CH000139
0C 86	0	0000	TEMP	DC		0	TEMPORARY STORAGE	CRD00140
0087	0	0400	INIT	DC		/0400	IOCC TO INITIATE 1/0 0	CRD00141
0088	0	2075	CONST	DC		SETUP-CA18-	-1+/2000	CRD00142
0089	0	1700	SENSE	DC		/1700	SENSE DSW WITHOUT RESET 0	CRD00143
0084	0	0000	BUSY	DC		0 .	ROUTINE BUSY INDICATOR	CR000144
0088	ō	1402	FEED	DC.		(1402	LOCC TO FEED 1 CARD	CRDOOLAS
0080	ň	0001	DODO	50		A1		CPD00144
0080	0	0001	00001	50		T1		CREGOULAT
0080	0	0004	00004			T-1		CRUUUL#/
008E	_	0008		nc		+8		
008F	0	0000	D0008					CR000140
0090	0 0	0050	D0008 D0080	DC		+80		CRD00149
0091	0 0 0	0050	D0008 D0080 D0081	DC DC		+80 +81		CRD00149 CRD00150
	0000	0050 0051 1000	D0008 D0080 D0081 H1000	DC DC DC		+80 +81 /1000		CRD00149 CRD00150 CRD00151
0092	000000	0050 0051 1000	D0008 D0080 D0081 H1000 H1001	DC DC DC DC		+80 +81 /1000 /1001		CRD00149 CRD00150 CRD00151 CRD00151
0092	000000	0050 0051 1000 1001 7003	D0008 D0080 D0081 H1000 H1001 H7003			+80 +81 /1000 /1001 /7003	INSTRUCTIONS = MDX X +3	CRD00149 CRD00150 CRD00151 CRD00152 CRD00153
0092 00 93	0000000	0050 0051 1000 1001 7003	D0008 D0080 D0081 H1000 H1001 H7003			+80 +81 /1000 /1001 /7003	INSTRUCTIONS = MDX X +3	CRD00149 CRD00150 CRD00151 CRD00152 CRD00153
0092 0093 0094	00000000	0050 0051 1000 1001 7003 0000	D0008 D0080 D0081 H1000 H1001 H7003 COUNT	DC DC DC DC DC DC DC		+80 +81 /1000 /1001 /7003 0	INSTRUCTIONS = MDX X +3 NO. WORDS TO XFER	CRD00149 CRD00150 CRD00151 CRD00152 CRD00153 CRD00154 CRD00154
0092 0093 0094 0095	00000000	0050 0051 1000 1001 7003 0000 0000	D0008 D0080 D0081 H1000 H1001 H7003 COUNT ERROR	DC DC DC DC DC DC DC DC DC		+80 +81 /1000 /1001 /7003 0	INSTRUCTIONS = MDX X +3 NO• WORDS TO XFER SKIP ONE CARD INDIC	CRD00149 CRD00150 CRD00151 CRD00152 CRD00153 CRD00153 CRD00155
0092 0093 0094 0095 0096	000000000	0050 0051 1000 1001 7003 0000 0000 0000	D0008 D0080 D0081 H1000 H1001 H7003 COUNT ERROR INDIC	DC DC DC DC DC DC DC DC DC DC		+80 +81 /1000 /1001 /7003 0 0	INSTRUCTIONS = MDX X +3 NO• WORDS TO XFER SKIP ONE CARD INDIC FEED CHK (RD STATION) IND	CRD00149 CRD00150 CRD00151 CRD00152 CRD00153 CRD00154 CRD00155 CRD00156
0092 0093 0094 0095 0096 0097	00000000000	0050 0051 1000 1001 7003 0000 0000 0000 0000	D0008 D0080 D0081 H1000 H1001 H7003 COUNT ERROR INDIC RSTRT			+80 +81 /1000 /1001 /7003 0 0 0	INSTRUCTIONS = MDX X +3 NO: WORDS TO XFER SKIP ONE CARD INDIC FEED CHK (RD STATION) IND RESTART INFO - WORD COUNT	CRD00149 CRD00150 CRD00151 CRD00152 CRD00153 CRD00153 CRD00154 CRD00155 CRD00156 CRD00157
0092 0093 0094 0095 0096 0097 0098	000000000000	0050 0051 1000 1001 7003 0000 0000 0000 0000	D0008 D0080 D0081 H1000 H1001 H7003 COUNT ERROR INDIC RSTRT			+80 +81 /1000 /1001 /7003 0 0 0 0 0	INSTRUCTIONS = MDX X +3 NO• WORDS TO XFER SKIP ONE CARD INDIC FEED CHK (RD STATION) IND RESTART INFO - WORD COUNT DATA ADDR	CRD00149 CRD00150 CRD00151 CRD00152 CRD00153 CRD00155 CRD00156 CRD00156 CRD00158
0092 0093 0094 0095 0096 0097 0098 0099	00000000000000	0050 0051 1001 7003 0000 0000 0000 0000 0000 000	D0008 D0080 D0081 H1000 H1001 H7003 COUNT ERROR INDIC RSTRT SETUP			+80 +81 /1000 /1001 /7003 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	INSTRUCTIONS = MDX X +3 NO• WORDS TO XFER SKIP ONE CARD INDIC FEED CHK (RD STATION) IND RESTART INFO - WORD COUNT DATA ADDR INITIATE IOCC SETUP - GET	CRD00149 CRD00150 CRD00151 CRD00152 CRD00153 CRD00154 CRD00155 CRD00156 CRD00155 CRD00155
0092 0093 0094 0095 0096 0097 0098 0099	000000000000000000000000000000000000000	0050 0051 1000 1001 7003 0000 0000 0000 0000 000	D0008 D0080 H1000 H1001 H7003 COUNT ERROR INDIC RSTRT SETUP			+80 +81 /1000 /1001 /7003 0 0 0 0 0 /02FC /02FF	INSTRUCTIONS = MDX X +3 NO. WORDS TO XFER SKIP ONE CARD INDIC FEED CHK (RD STATION) IND RESTART INFO - WORD COUNT DATA ADDR INITIATE IOCC SETUP - GET - PUT	CRD00149 CRD00150 CRD00151 CRD00152 CRD00153 CRD00156 CRD00156 CRD00156 CRD00156 CRD00159 CRD00150 CRD00150
0092 0093 0094 0095 0096 0097 0098 0099 0098	000000000000000000000000000000000000000	0050 0051 1000 1001 7003 0000 0000 0000 0000 000	D0008 D0080 D0081 H1000 H1001 H7003 COUNT ERROR INDIC RSTRT SETUP			+80 +81 /1000 /1001 /7003 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	INSTRUCTIONS = MDX X +3 NO. WORDS TO XFER SKIP ONE CARD INDIC FEED CHK (RD STATION) IND RESTART INFO - WORD COUNT DATA ADDR INITIATE IOCC SETUP - GET - PUT - FEFD	CRD00149 CRD00150 CRD00151 CRD00153 CRD00153 CRD00155 CRD00155 CRD00157 CRD00155 CRD00155 CRD00161
0092 0093 0095 0096 0097 0098 0099 0098		0050 0051 1000 1001 7003 0000 0000 0000 0000 000	D0008 D0080 D0081 H1000 H1001 H7003 COUNT ERROR INDIC RSTRT SETUP			+80 +81 /1000 /1001 /7003 0 0 0 0 0 /02FC /02FF /02FF /02FF	INSTRUCTIONS = MDX X +3 NO. WORDS TO XFER SKIP ONE CARD INDIC FEED CHK (RD STATION) IND RESTART INFO - WORD COUNT DATA ADDR INITIATE LOCC SETUP - GET - PUT - FEED	CRD00149 CRD00150 CRD00153 CRD00153 CRD00154 CRD00154 CRD00156 CRD00156 CRD00156 CRD00156 CRD00160 CRD00160
0092 0093 0094 0095 0096 0097 0098 0099 009A 009B 009C	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	0050 0051 1000 1001 7003 0000 0000 0000 0000 000	D0008 D0080 D0081 H1000 H1001 H7003 COUNT ERROR INDIC RSTRT SETUP			+80 +81 /1000 /1001 /7003 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	INSTRUCTIONS = MDX X +3 NO. WORDS TO XFER SKIP ONE CARD INDIC FEED CHK (RD STATION) IND RESTART INFO - WORD COUNT DATA ADDR INITIATE IOCC SETUP - GET - FEED - STK	CRD00149 CRD00150 CRD00152 CRD00152 CRD00153 CRD00154 CRD00156 CRD00156 CRD00156 CRD00160 CRD00160 CRD00160
0092 0093 0094 0095 0096 0097 0098 0099 009A 0098 0096 0090	00000000000000000	0050 0051 1001 7003 0000 0000 0000 0000 0000 000	D000B D0080 D0080 H1000 H1000 H1001 H7003 COUNT ERROR INDIC RSTRT SETUP			+80 +81 /1000 /1001 /7003 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	INSTRUCTIONS = MDX X +3 NO• WORDS TO XFER SKIP ONE CARD INDIC FEED CHK (RD STATION) IND RESTART INFO - WORD COUNT DATA ADDR INITIATE IOCC SETUP - GET - PUT - FEED - STK COLUMN IOCC SETUP - GET	CRD00149 CRD00153 CRD00153 CRD00153 CRD00153 CRD00154 CRD00155 CRD00155 CRD00155 CRD00165 CRD00161 CRD00163
0092 0093 0094 0095 0096 0097 0098 0099 0098 0098 0098 0098 0095	000000000000000000000000000000000000000	0050 0051 1000 1001 7003 0000 0000 0000 0000 000	D000B D0080 D0081 H1000 H1001 H7003 COUNT ERROR INDIC RSTRT SETUP			+80 +81 /1000 /1001 /7003 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	INSTRUCTIONS = MDX X +3 NO. WORDS TO XFER SKIP ONE CARD INDIC FEED CHK (RD STATION) IND RESTART INFO - WORD COUNT DATA ADDR INITIATE IOCC SETUP - GET - PUT - FEED - STK COLUMN IOCC SETUP - GET - PUT	CRD00149 CRD00150 CRD00152 CRD00152 CRD00154 CRD00154 CRD00155 CRD00155 CRD00155 CRD00155 CRD00161 CRD00163 CRD00163 CRD00164

	CRDOO251				GNΒ				SHOO
	02500090	EXIT	LINI	I	25 8	SSTN	40000834	to	3300
	CR000249	EXECUTE COLUMN 1/0	COLM		01X	OSTN	5680	0	3300
	84500080	SET ADDR FOR NEXT COLUMN	COLM+1	Г	XQM	BITN	A80010A7	10	00€C
	74500080	LOCATION	COLM		012		8600	0	8300
	90200000	PUNCH FROM TEMPORARY	8004		01		2605	0	VE00
	##2000HJ	M (MO	PO000		012		6600	0	6300
	64200GH2	100 NI (21 LIB) LIB	80000 wT00	r	00		C4900094	10	8300
	24200043	HONDA ADIS THON THE	1++100	÷.	YOW	Q L I NI	\$80008VJ	10	9300
	19200080	NEXT COL	221N	•	XOW		9000000	.0	5300
	CRD00240	AES' SEL LO SKID	1++10002	٦	XQM		7001000	10	1300
	CRD00239	IS THIS READ COL INTERR	-+911N	, r	SSB		#3000L2#	10	-000
	CRD00238	AES	811N		XOW		0004	ō	3000
	CED00237	ANY MORE COLS TO PROCESS	COUNT+-1	٦	XQM		74FF0094	10	CODC
	CRD00236	SENSE DSM MITH RESET	1-9AHC		01X		9480	о	9000
	CRD00235	COLUMN REQUEST INTERRUPT	RAHD		012	411N	640Q	0	A 000
	¢ES000934	******	*******	***	****	*****			
	CRD00233	COMPLETE INTERRUPT.	90 NA 3TAI	TIN	1 01	*			
	CR000232	ATION IS GIVEN TO THE 1442 *	DIGNI NA .	HED	PUNC	*			
	16500083	* SI COLUMN REQUESTED IS	MHEN IHE F	- H	0000	*			
	472000AD	SELA LINDNED DEE VOLDMA *	HAN ROA 210	FIGG	1NTE	*			
	822000HD	PER THE READESTED NO. OF *	1 NE 000 AF	100	HEAS	*			
	12200083			04	5141				
	222000H2		*********	***	****	******			
	CRD00225	LERRUPT PROCESSING *	NI NWOTOD			*			
	CR000224	******	********	***	****	*****			
	CRDOOSS3		SITN		XOM		8307	о	6000
	CRD00222	SET BIT 1 OF INDIC	RRROR		012		DOBC	ō	8000
	CRD00221		EOOTH		r٥	JEIIN	8800	0	1000
	CRD00220	NO ZKIH IL LO CHK (BD)	NT12+C	٦	228	8ET1N	4C0200C2	10	5000
	CR000219	EXIL	STNI	1	228		70000B3A	۱o	£000
	CR000218	NOITARERO ONI ETAITINI	1-11N1		01X		5880	0	SOD2
	CRD00217		COLM		OT2		2800	0	1000
	CR000216		1+19129		٦٥		2000	0	0000
	CRD00215	YATAA AOA 9UTA2	COUNT		OTS		DOC¢	0	300CF
	CR000214		19129		го	EITN	8202	0	3000
	C1200082		SINI	I	228		70000824	10	2200
	21200085	SKIP 121 CARD	FFED-1	_	01X		OBBE	0	8200
	CB000511	UN	-+•EIIN		228		AC1800CF	10	6000
	CE000510	IS CARD SKIP NECESSARY	EBBOB	_			5005		8200
	802000a0	WALL IN READER READY	3. SITU	1	228	71.1.1.1	2004074	10	9000
	10200080	CHARD IST HINE FOR		-	250	CITIN		10	8000
	20200083	Udda Tal alwa -ON	-T'JUIIN	1	729		1001	10	2000
	CR000206	NI KIYA TOO ANG SAM	1+14158		000		9003	ň	1200
	CR000204	AES	3+8611N	٦	258		5000000	ιŏ	1800
	CR000203	12 FUNCT FEED	1		ARA		1081	ō	3800
	CR000202	YES. DONT SKIP	AT12+E	٦	258		S200#02#	to	2800
	CR000201	IZ EUNCL PUNCH	LINI		۳D		8202	0	8800
	CRD00200	SAVE FD CHK (RD STA) 14D	· G		VTS	ETTN	5001	0	4800
	C6000136		3011N		XQM		70F3	0	6800
/	8610008C	YES. EJECT LAST CD	FEED-1		01X		1080	0	8800
	79100093	ON	~+		228		8184	0	7800
	96100080	15 FUNCTION PUT	COLM		ROR		FOCD	0	9800
	56100085		10000		4		9008	0	5800
	46100080		800V		201	XUIIN	(0000		7000
	26100083		STUI	•	720		20000830	10	1400
	161000103	X318 SHITIOD GV315	A3118 01		AHC		0101	~	0900
	06100080	DECKEWENI TOCE COONL			dON		0001	ŏ	1400
	68100080	NO. TERMINATE FUNCT	1-+05	٦	XOW	3011N	74FF0032	00	dyor.
	CED00188	YES. INITIATE FUNCT	NT12+C	٦	258		¢C0200C2	10	8400
	CR000187		RCROR		510		D0EA	0	4400
	CED00186		91		AAZ		0181	0	6400
	C6D00182	YES. WAS THIS SKIP OP	S		٩٦S		1002	0	8400
	CBD00184		ERROR		r۵		COED	0	VA00
	CRD00183	NO+ LAST CARD	+Z•X0IIN	٦	228	8011N	4C280084	10	2400
	CR000182	NO • EGROR	O+111N	Г	228		A8002024	10	EAOC
	CRD00181	IS OPERATION OK	E		A.12		£001	0	SAOC
	08100095	SENSE DOW WITH RESET	I-AAHO		OIX		OBEO	0	IAOC
	67100093		RAHD		012		DOES	0	0400
	ST 100093	OPER COMPLETE INTERRUPT	10000		· ¥	OILN	BOEC	0	3000
	LL 100083		*********	1220	****	******			
		202 ARE FOSTIUMED AND INC *	19-1	1 1	314#	*			
	PF 100000	* IA •YUAJY CAMUJAC SAPI 2 * 3ut nun finnitions 306 200		IBVE!	JINI	*			
	21000aJ	* SAH NOTANGHU JHI JIINU 240	SOUTINE LO	เลกร	THE	*			
	2/100082	# TSIMARKIO OTTATA	SSECTO	เกษษ	JINI	*			
	12100080	* EIADIGNI OT GETH	IS DECREM	яэт	клор	*			
	02100080	* SOOT AND THE LOCS *	S SI BNIL	ററപ	3HT	*			
	69100080	* • OBTOBTE DEEN DETECTED . *	E 04º IL N	IN I T	0092	*			
	691000H2	* TERED FROM INTERR LEVEL *	STION IS E	604	2 I H T	*			
	78000167	************************	********	***	****	*****			
	CBD00199	* INTERRUPT PROCESSING *	Ob COWERE.			*			
	CBD00192	*****	*********	***	****	*****			

Appendix A. Listing of Subroutines

Subroutine		Names	Subroutines Required
Loader Reinitialization (card only) Data Switch Sense Light On Overflow Test Divide Check Test Function Test Trace Start Trace Stop Integer Transfer of Sign Real Transfer of Sign (E) Real Transfer of Sign (S)	LC DA SL O' FC TS TS ISI ES FS	DAD NTSW ITE, SLITT VERF CCHK TST TRT TRT GN GN IGN	None None None None None TSET TSET TSET Sone ESUB, ELD FSUB, FLD
Called by LIBF (Card/Paper Tape) Real IF Trace (E) Real IF Trace (S) Integer IF Trace (E) Integer Arithmetic Trace (E) Integer Arithmetic Trace (S) Real Arithmetic Trace (S) Real Arithmetic Trace (S) Real Arithmetic Trace (E) Real Arithmetic Trace (E) Computed GO TO Trace (S) Computed GO TO Trace (S) Trace Test-Set Indicator Pause Stop Subscript Calculation Store Argument Address I/O Linkage (E) I/O Linkage (S) Card Input/Output Printer-Keyboard Output Printer-Keyboard Input/Output 1132 Printer Output Paper Tape Input/Output Card Code-EBCDIC Conversion Conde Printer Code Table Card-Keyboard Code Table Address Calculation	VII WI VI WI VI WI VA WI VA WI VI VI VI WI VI WI VI WI VI VI VI VI VI SUU SUU VI VI SUU SUU SUU SUU SUU SUU SUU SUU SUU SU	F F F F F AR, VIARX AR, WIARX AR, WIARX AR, WIARX AR, WIARX AR, WIARX AR, WIARX AR, WIARX SOTO	TTEST, VWRT, VIOF, VCOMP FSTO, TTEST, WWRT, WIOF, WCOMP TTEST, WWRT, VIOF, VCOMP TTEST, WWRT, WIOI, WCOMP TTEST, WWRT, WIOI, VCOMP FSTO, TTEST, WWRT, WIOF, WCOMP TTEST, WWRT, WIOI, WCOMP TTEST, WWRT, WIOI, WCOMP None None None None None ELOAT, ELD/ESTO, IFIX HOLEZ GETAD, EBCTB, HOLEZ None None None None None None None None
		• • •	

Figure 23 is a listing of the Card/Paper Tape System Subroutine Library. The Disk Monitor 2 System Library is listed in Figure 24.

Figure 23. C/PT System Subroutine Library (Part 1 of 3)

I

Subroutine	Names	Subroutines Required
ARITHMETIC AND FUNCTIONAL		
Called by CALL		
Real Hyperbolic Tangent (E) Real Hyperbolic Tangent (S) Real Base to Real Exponent (E) Real Base to Real Exponent (S) Real Natural Logarithm (E) Real Natural Logarithm (S) Real Exponential (E) Real Exponential (S) Real Square Root (E) Real Square Root (S) Real Trigonometric Sine/Cosine (S) Real Trigonometric Arctangent (E) Real Trigonometric Arctangent (S) Fixed-Point Square Root Real Absolute Value (E) Real Absolute Value (S) Integer Absolute Value Real Binary to Decimal/Real Decimal to Binary	ETNH, ETANH FTNH, FTANH EAXB, FAXBX FAXB, FAXBX ELN, EALOG FLN, FALOG EXPN, EEXP FXPN, FEXP ESQR, ESQRT FSQR, FSQRT FSIN, ESINE, ECOS, ECOSN EATN, EATAN FATN, FATAN XSQR EAVL, FABS FAVL, FABS FAUL, FABS FBID, FDTB	EEXP, ELD/ESTO, EADD, EDIV, EGETP FEXP, FLD/FSTO, FADD, FDIV, FGETP EEXP, ELN, EMPY FEXP, FLN, FMPY XMD, EADD, EMPY, EDIV, NORM, EGETP FSTO, XMDS, FADD, FMPY, FDIV, NORM, FGETP XMDS, FARC, EGETP ELD/ESTO, EADD, EMPY, EDIV, EGETP FLD/FSTO, FADD, FMPY, FDIV, EGETP EADD, EMPY, NORM, XMD, EGETP FADD, FMPY, NORM, XMD, EGETP FADD, FMPY, NORM, XMDS, FSTO, FGETP EADD, EMPY, FDIV, XMDS, FSTO, FGETP None EGETP FGETP None None
Called by LIBF		
Get Parameters (E) Get Parameters (S) Real Base to Integer Exponent (E) Real Base to Integer Exponent (S) Real Reverse Divide (E) Real Reverse Divide (S) Real Divide (S) Real Multiply (E) Real Multiply (E) Real Multiply (S) Real Reverse Subtract (E) Real Add/Subtract (S) Load/Subtract (S) Fixed Point Double Word Divide Fixed Point Double Word Divide Fixed Point Fractional Multiply Fixed Point Fractional Multiply Fixed Reverse Sign Integer Base to an Integer Exponent Normalize Real Arithmetic Range Check	EGETP FGETP EAXI, EAXIX FAXI, FAXIX EDVR, EDVRX FDVR, FDVRX EDIV, FDIVX FDIV, FDIVX FMPY, FMPYX ESBR, ESBRX FSBR, FSBRX FADD, EADDX, ESUB, ESUBX FADD, FADDX, FSUB, FSUBX ELD, ELDX, ESTO, ESTOX FLD, FLDX, FSTO, FSTOX XDD XMD XMDS SNR FLOAT IFIX FIXI, FIXIX NORM FARC	ELD FLD ELD/ESTO, EMPY, EDVR FLD/FSTO, FMPY, FDVR ELD/ESTO, EDIV XDD, FARC FARC XMD, FARC XMD, FARC XMDS, FARC EADD FARC, NORM NORM, FARC None None XMD None None None None None None None None
DUMP		
Called by CALL		
Dump Status Area Selective Dump on Console Printer Selective Dump on Printer	DMP80 DMTX0, DMTD0 DMPX1, DMPD1	None WRTYO PRNTI
INTERRUPT LEVEL* Level 0 Level 1 Level 2 Level 3 Level 4 *These subroutines are not identified by name in the	card and paper tape systems	None None None None None
CONVERSION		
Called by LIBF		
Binary to Decimal Binary to Hexadecimal Decimal to Binary EBCDIC to Console Printer Code IBM Card Code to or From EBCDIC IBM Card Code to Console Printer Code	BINDC BINHX DCBIN EBPRT HOLEB HOLPR	None None EBPA, PRTY EBPA, HOLL HOLL, PRTY

| Figure 23. C/PT System Subroutine Library (Part 2 of 3)

Subroutine	Names	Subroutines Required
Called by LIBF (Cont'd)		
Hexadecimal to Binary EBCDIC to or from PTTC/B IBM Card Code to or from PTTC/B PTTC/8 to Console Printer Code IBM Card Code to or from EBCDIC EBCDIC and PTTC/B Table IBM Card Code Table Console Printer Code Table	HXBIN PAPEB PAPHL PAPPR SPEED EBPA HOLL PRTY	None EBPA EBPA, HOLL None None None None None
DISK SUBROUTINE INITIALIZE		
Called by CALL		
Set Pack Initialization Subroutine	SPIRO, SPIR1, SPIRN	DISKO, DISKI, DISKN
NTERRUPT SERVICE		
Called by LIBF		
Card Disk Paper Tape Plotter 1132 Printer Karb Land (Canada Brinter	CARDO, CARDI DISKO, DISKI, DISKN PAPTI, PAPTN PLOTI PRNTI TVPEO, WETVO	1LS00, 1LS04 1LS02 1LS03 1LS03 1LS01 RO11 RRTX, 1LS04
	11120, WK110	
Standard Plot Calls Standard Precision Character	FCHAR	ESIN ECOS EPIOT ECHRX FID ESTOX ESTO
Standard Precision Scale Standard Precision Grid Standard Precision Plot	SCALF FGRID FPLOT	FRULE FPLOT, POINT, FADD, FLD, FSTO, SNR FMOVE, YPLT, PLOTI
Extended Plot Calls		
Extended Precision Character Extended Precision Scale Extended Precision Grid Extended Precision Plot	ECHAR SCALE EGRID EPLOT	ESIN, ECOS, EPLOT, ECHRX, ELD, ESTO, ESTOX ERULE EPLOT, POINT, EADD, ELD, ESTO, SNR EMOVE, XYPLT, PLOTI
Common Plot Call		
Point Characters	POINT	PLOTI
Standard Plot LIBFs		
Standard Precision Annotation Standard Precision Plot Scaler	FCHRX, FCHRI, WCHRI FRULE, FMOVE, FINC	FLCAT, FMPY, IFIX, FADD, FLDX, FINC, XYPLI, PLOTI, FSTOX, FLD FLDX, FSUBX, FMPYX, FLD, FSTOX, FMPY, IFIX, FADD
Extended Plot LIBFs		
Extended Precision Annotation	ECHRX, ECHRI, VCHRI	FLOAT, EMPY, IFIX, EADD, ELDX, EINC, XYPLT,
Extended Precision Plot Scaler	ERULE, EMOVE, EINC	ELDX, ESUBX, EMPYX, ELD, ESTOX, EMPY, IFIX, EADD, ESTO
Common Plot LIBFs		
Pen Mover Interfere	XYPLT	PLOTI
Interrupt Service	PLOTX	
Synchronous Communications Adaptor Subroutines		
Synchronous Communications Adaptor (SCA)	SCATI	IOLOG/CPLOG, ILSOI
SCA (BSC, Point-to-Point Mode)	SCAT2	IOLOG/CPLOG, ILS01
1132-SCA Print with Overlap	PRNT2	ILSOI
4 of 8 Code to EBCDIC, EBCDIC to 4 of 8 Code 4 of 8 Code to IBM Card, IBM Card Code	EBC48 HOL48	HACV, SIRIB HXCV, HOLCA, STRTB
to 4 of 8 Code 4 of 8 Code to Table of Displacements Table of IBM Card Codes Table of A of 8 and FBCDIC Codes	HXCV HOLCA STRTB	None None None

Figure 23. C/PT System Subroutine Library (Part 3 of 3)

n Library Programs Names Type Subty			(13-13)
	· -		
DISC	2 -	SYSUP, REREC, LISKZ	U6 C
I DENT	2 -	CALPR, DISKZ	U6F U6C
10	2	EISKZ	
COPY ADRWS (cannot be called)	2 - 2 -	RDREC, DISKZ Linked from DUP DWADR	U6 B U6 A
DICIB	2 -	REREC, EISKZ	U6 C
LSL ET	2 -	FSLEN, DISKZ	U6E
MODSF	2 -	LISKZ	U6 I
MOEIF	2 -		U6H 1-1-1
DECIN	2 -	FLE, NORM	WIT
PTUTL	2 -	PAPHL, PAPPR, PAPT1, 1YPE0	U6J
DMIDO, DMIXO	4,0	WR1Y0	U5 P
DMPD1, DMFX1	4,0	PRNT1	U5 C
DEF80	4,0	NONE FOITN FOVEII	USA USE
CALPR	4.0	FSLEN, FSLSO	117A
RDREC	4,0	FSLEN	U7C
FSLEN, FSYSU	4,0	CISKZ	U 7 E
IOLOG/CPLOG	4,0	Ncne	
DAISW	4,8	None	T 3A
DVCHK	4,8	None	13 B
FCIST	4,8	None	13C
OVERF	4,8	None	T3E
PDURP	4,0	SFID, SIDAI, SIDAF, SWRT, SCOMP	T3F
	DISC IDENT ID COPY ADRWS (cannot be called) DICIB CSLET MODSF MODIF DFCNV PTUTL DMTD0, DMTX0 DMTD0, DMTX0 MPD1, DMFX1 DMP80 SYSUP CALPR RDREC FSLEN, FSYSU IOLOG/CPLOG DATSW DVCHK FCTST OVERF PDUMP	DISC 2 - IDENT 2 - ID 2 - COPY 2 - ADRWS (cannot 2 - be called) DLCIB 2 - CSLET 2 - MODSF 2 - MODIF 2 - DFCNV 2 - PTUTL 2 - PTUTL 2 - PTUTL 2 - DMIDO, DMTXO 4,0 DMPD1, DMFX1 4,0 DMP80 4,0 SYSUP 4,0 CALPR 4,0 RDREC 40 FSLEN, FSYSU 4,0 IOLOG/CPLOG 4,0 DATSW 4,8 PDUMP 4,0	DISC2SYSUP, REREC, LISKZID2CALPR, DISKZID2RIREC, CALPR LISKZCOPY2RDREC, DISKZADRWS (cannot be called)2Linked from DUP DWADRDLCIB2RIREC, LISKZCSLET2FSLEN, DISKZMODSF2LISKZMODSF2DISK2DFCNV2DISK1, ELD FLL, NORMPTUTL2PAPHL, PAPPR, PAPT1, 1YPEODMTDO, DMTXO4,0WRTYO FLL, NORMPTUTL2PAPHL, FAPPR, FSLEN, FSYSUCALPR4,0FSLEN, FSYSU CALPR4,0FSLEN, FSYSUIOLOG/CPLOG4,0NoneDATSW4,8 None PDUMPNone YOURFDATSW4,8 NONE PDUMPNone YOURFDATSW4,8 NONE PDUMPNone YOURF, SWRT, SCOMP

Figure 24. 1130 Disk Monitor Version 2 System Library (Part 1 of 9)

System Library Programs	Names	Type and Subtype	Subroutines Required	ID Field (73-75)
Common FORTRAN Calls (continued)				,
Sense Light Control and Test FORTRAN Trace Stop FORTRAN Trace Start Integer Transfer of Sign	SLITE, SLITT TSTOP TSTRT ISIGN	4,8 4,8 4,8 4,8	None ISET TSET None	T3G T3 H T3 I T3 D
Extended Arith/Funct Calls				
Extended Precision Hyperbolic Tangen	it ETANH, ETNH	4,8	EEXP, EADD, ELIV, EGETP,	S 21
Extended Precision A**B Function	EAX B, EAX BX	4,8	EEL/ESIO EEXP, ELN, FMPY	S2C
Extended Precision Natural Logarithm	ELN, EALOG	4,8	XMD, EADD, EMPY, EDIV,	S2E
Extended Precision Exponential	EEXP, EXPN	4,8	XMC, FARC,	S2D
Extended Precision Square Root	esor, esort	4,8	EADD, EMPY, ECIV, EGETP,	S2H
Extended Precision Sine-Cosine	ESIN, ESINE, ECCS, ECCSN	4,8	ELC/ESTO EADD, EMPY, NORM, XMD, EGETP	S2G
Extended Frecision Arctangent	EATN, EATAN	4,8	EADD, EMPY, EDIV, XMD,	S2E
Extended Precision Absolute Value Function	EABS, EAVL	4,8	EGETP	S2A
FORTRAN Sign Transfer Calls				
Extended Precision Transfer cf Sign Standard Precision Transfer cf Sign	ESIGN FSIGN	4,8 4,8	ESUB, ELD FSUB, FLD	S2 F R2 F
Standard Arith/Funct Calls				
Standard Precision Hyperbolic Tangen	t FIANH, FINH	4,8	FEXP, FADD, FCIV, FGETP,	R 21
Standard Precision A**B Function Standard Precision Natural Logarithm	FAXB, FAXBX FIN, FALOG	4,8 4,8	FEXP, FLN, FMPY FSIO, XMDS FADD, FMPY, FLIV, NCRM ECETR	R2C R2 E
Standard Precision Exponential	FEXP, FXPN	4,8	XMDS, FARC,	R2D
Standard Precision Square Root	FSQR, F <i>S</i> QR1	4,8	FADD, FMPY, FDIV, FGETP, FLI/FSTC	R 2H
Standard Precision Sine-Cosine	FSIN, FSINE FCCS, FCCSN	4,8	FADD, FMPY, NORM, XMDS, ESTO ECETP	R2G
Standard Precision Arctangent	FAIN, FAIAN	4,8	FAID, FMPY, FEIV, XMDS,	R2 B
Standard Precision Absolute Value Function	FABS, FAVL	4,8	FGETP	R2A

Figure 24. 1130 Disk Monitor Version 2 System Library (Part 2 of 9)

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System Library Programs	Names	Type and Subtype	Subroutines Required	IE Field (73-75)
Common Arith/Funct Calls Fixed Point (Fractional) Square Roct Integer Absolute Function Floating Binary/EBC Decimal Conversions	XSQR IAES FBID (BIN. IO DEC.) FDIB (DEC. TO BIN.)	4,8 4,8 4,0	None Ncne None	T 1C T1B T1A
Flipper for LOCAL/SOCAL Subprograms	FL 1PR	4,0	DISKŹ, DISK1, or DISKN	U5D
FORTRAN Trace Subroutines				
Extended Floating Variable Trace	SEAR, SEARX	3,0	ESTO, TIEST, SWRI, SIOF,	S2J
Fixed Variable Trace	SIAR, SIARX	3,0	TTEST, SWRT,	16B
Standard Floating IF Trace	SFIF	3,0	FSTO, TTEST, SWRT, SIOF, SCOMP	R2K
Extended Floating IF Trace	SEIF	3,0	FSIO, TIEST, SWRT, SIOF, SCOMP	S2 K
Fixed IF Trace	SIIF	3,0	TIEST, SWRT, SIOI, SCOMP	16C
Standard Floating Variable Trace	SFAR, SFARX	3,0	FSTO, TTEST, SWRT, SIOF, SCOMP	R2J
GO TO Trace	SGOTO	3,0	ITEST, SWRI, SIGI, SCOMP	16A
Nondisk FORTRAN Format 1/0				
FORTRAN Format Subroutine	SFIO, SIOI, SIOAI, SICF, SICAF, SICFX, SCCMP, SWR1, SRED, SICIX	3,3	FLOAT, IFIX, ELE/FSTC or FLE/FSTC, PAUSE	т 4С
FORTRAN Find Subroutine	SEFNE	3,1	DISKZ, DISK1, Cr DISKN	T 4B
<u>Eisk FORTRAN I/O</u>	SDFIO, SDRED, SEWRT, SDCOM, SDAF, SDF, SDI SDIX, SDFX, SDAI	3,1	DISKZ, DISK1, or DISKN, FAUSE	14A
Unformatted FORTRAN Disk 1/0	UFIO, URED, UWRT, UIOI, UICF, UICAI, UICAF, UICFX, UICIX, UCCMF, BCKSP, EOF, REWND	3,1	DISKZ, DISK1, or DISKN, PAUSE	T4C

Figure 24. 1130 Disk Monitor Version 2 System Library (Part 3 of 9)

System Library Programs	Names	Type and Subtype	Subroutines Required	ID Field (73-75)
FORTRAN Common LIBFS				
FORTRAN Pause FORTRAN Stop FORTRAN Subscript Displacement	PA USE STOP SUB SC	3,0 3,2 3,0	None None None	T2 A T2 B T2 D
FORTRAN Subroutine Initialization FORTRAN Trace Test and Set	SUBIN TIEST, TSET	3,0 3,0	None None	T 2C 12 E
FORTRAN 1/O and Conversion Subroutines				
FORTRAN 1442 Input/Output Subroutine	CARDZ	5,3	HOLEZ, GETAD, EPCTB, HOLTB, ILSO0 ILSO4	т5а
FORTRAN 1442 Output Subroutine	PNCH Z	5,3	HOLEZ, GFTAD, EFCTE, HOLTB, ILSO0, ILSO4	T5G
FORTRAN 2501 Input Subroutine	READ Z	5,3	HOLEZ, GETAD, FECTE, HOLTB, ILSO4	т5Ј
Disk I/O Routine (Part of Supervisor)	DISKZ	-	ILS02	
FORTRAN Paper Tape Subroutine	PAFIZ	5,3	ILSO4	T5 F
FORTRAN 1132 Printer Subroutine	PRNIZ	5,3	ILS01	15 H
Call to PRNTZ to Call to PRNT2 Conversion	PR 122	5,3	PRNT2, ILSO1	WIK
FORTRAN 1403 Printer Subroutine	PRNZ	5,3	ILSO4	T51
FORTRAN Keyboard-Typewriter Subroutine	IYPEZ	5,3	GETAD, EECTB, Holez, Ilso4	15 K
FORTRAN Typewriter Subroutine	WRIYZ	5,3	GETAD, FECTB, ILS04	T5L
FORTRAN 1627 Plotter Subroutine	PLOTX	5,0	ILSOJ	V 1L
FORTRAN Hollerith to EBCDIC Conversion	HOLEZ	3,3	GETAD, EBCTB, HOLTB, PAUSE	15D
FORTRAN Get Address Routine	GEIAD	3,3	None	T5C
FORTRAN EBCDIC Table	EBCIB	3,3	None	T5 B
FORTRAN Hollerith Table	HOLIB	3,3	None	T5 E
FORTRAN Multiple Terminal Communications Adapter (MICA) Call Interface	M1CA Z	4,0	МТСАО	W5 C
Extended Arith/Funct LIBFs				
Extended Precision Get Parameter	EGFTP	3,2	ELD	s 1e
Extended Precision A**I Function	EAXI, EAXIX	3,2	ELD/ESTO EMPY, EDVR	S 1B
Extended Precision Divide Reverse	EDVR, FDVRX	3,2	ELD/FSTO, EDIV	SID
Extended Precision Float Divide	ECIV, EDIVX	3,2	XDD, FARC	S1C
Extended Precision Float Multiply	ЕМРҮ, ЕМРҮХ	3,2	XMD, FARC	S1G
Extended Precision Subtract Reverse	ESBR, EXBRX	3,2	EADD	S1 H
Extended Add-Subtract	EADD, ESUB, EADDX, ESUBX	3,2	FARC, NORM	STA
Extended Load-Store	ELD, ELDX, ESIO, ESICX	3,0	None	S1 F

Figure 24. 1130 Disk Monitor Version 2 System Library (Part 4 of 9)

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System Library Programs	Names	Type and Subtype	Subroutines Required	ID Field (73-75)
Standard Arith/Funct LIBFs				
Standard Precision Get Parameter	FGETP	3,2	FLD	R1E
Standard Precision A**I Function	FAXI, FAXIX	3,2	FLD/FSIO, FMPY, FDVR	R 1B
Standard Precision Divide Reverse	FDVR, FDVRX	3,2	FLD/FSTO, FDIV	R1D
Standard Precision Float Divide	FDIV, FDIVX	3,2	FARC	R1C
Standard Precision Float Multiply	FMPY, FMPYX	3,2	XMDS, FARC	R1G
Standard Precision Subtract Reverse	FSBR, FSBRX	3,2	FADD	R 1 H
Standard Add-Subtract	FADD, FSUB, FADDX, FSUBX	3,2	NORM, FARC	R1A
Standard Load-Store	FLD, FLDX, FSIO, FSTCX	3,0	None	R1F
Standard Precision Fractional Multiply	XMDS	3,2	None	S3 I
Common Arith/Funct LIBFs				
Final (Franting), Double Divide	145.5	2.2		
Fixed (Fractional) Double Divide	XDD XMD	3,2	XMD	SJG CZ H
Multiply	AD	5,2	NOTE	1 66
Sign Reversal Function	SNR	3,2	None	S 3F
Integer to Floating Point Function	FLOAI	3,0	NORM	S3C
Floating Point to Integer Function	IFIX	3,0	None	S3 D
I**J Integer Function	FIXI, FIXIX	3,2	None	S3 B
Normalize Subroutine	NORM	3,0	None	S3 E
Subroutine	FARC	3,2	None	S3 A
Interrupt Service Subroutines				
1442 Card Read Punch Input/Output (No error Parameter)	CARD 0	5,0	ILSOO, ILSO4	U2 A
1442 Card Read Punch Input/Output (Error Parameter)	CARD1	5,0	1LS00, 1LS04	U 2B
2501 Card Read Input (Nc Error Parameter)	READ0	5,0	ILSO4	U21
2501 Card Read Input (Error Parameter)	READ1	5,0	ILSO4	U 2 M
1442 Card Punch Output (No Error Parameter)	PNCH0	5,0	ILS00, ILS04	U 2H
1442 Card Punch Output (Error Parameter)	PNCH1	5,0	ILS00, 1LS04	U 2I
Multiple Sector Disk Input/Output (Part of Supervisor)	DISK 1	-	1LS02	
High Speed Multiple Sector Disk Input/Output (Part of Supervisor)	DISKN		1LS02	
Synchronous Communications Adapter (SCA) STR Mode	SCAT 1	5,0	IOLOG/CPLOG,	W1F
SCA (BSC, Point-to-Point Mode)	SCAT2	5,0	IOLOG/CPLOG, ILS01	W 1H
SCA (BSC, Multi-Point Mode	SCAT3	5,0	ICLOG/CPLOG, ILS01	W11
Paper Tape Input/Output	PAPT1	5,0	ILS04	U2D
Simultaneous Paper Tape Input/Output	PAPTN	5,0	ILS04	U2E
Character/Word Count Paper Tape Input/Output	PAPTX	5,0	ILS04	U2 F

Figure 24. 1130 Disk Monitor Version 2 System Library (Part 5 of 9)

System Library Programs	Names	Type and Subtype	Subroutines Required	IE Field (73-75)
Interrupt Service Subroutines (continued)				<u> </u>
Plotter Output Subroutine	PT.OT 1	5.0	11.003	1120
Plotter Output Subroutine	PLOTX	5.0	11.503	02G V1T
1132 Printer Output Subroutine	PRNT1	5 0	11201	
1132-SCA Print With Overlan	DDNT2	50	11001	
1403 Printer Output Subroutine	DDNT3	5,0		WIE
Keyboard/Console Printer Input/Output		5,0		02 K
Console Printer Input/Output	TIPEO	5,0	ILS04	U2N
Console Printer Output Subroutine	WRTY0	5,0	ILSO4	U 20
Subroutine	OMPR1	5,0	ILSO4	U2C
MICA Base Section	MT CAO	5,0	ILSO3, TSM41, TSTTY	W5B
MTCA 2741 Terminal Select	TSM41	4.0	Ncne	W5D
MTCA Teletype Select	TSTTY	4,0	None	W5E
Conversion Subroutines				
Binary Word to 6 Decimal Characters	EINCC	3,0	Ncne	U4B
Binary Word to 4 Hexadecimal Characters (Card Code)	BINHX	3,0	None	U 4C
6 Lecimal Characters (Card Code) to Binary Word	DCBIN	3,0	None	U4G
BCDIC to Console Printer Output Code	FBPRT	3.0	EBPA . PRTY	11 32
Card Code to EBCDIC-EBCDIC to Card	HOLEB	3,0	EEPA, HOLL	U3 B
Card Code to Console Printer Output	HOLPR	3,0	HOLL, PRTY	U 3C
4 Hexadecimal Characters (Card Code) to Binary Word	HXEIN	3,0	None	U 3 D
PTIC/8 to EBCDIC-EBCDIC to PTTC/8	PAPFB	3.0	EBPA	11 3E
PTTC/8 to Card Code-Card Code to	PAPHL	3,0	EEPA, HOLL	U3 F
PITC/8 to Console Printer Output Code	DADDR	3.0	FRDA DDTV	11.30
Card Code to FBCDIC-FBCDIC to Card	SPEED	3 0	None	113 11
	Dreed	5,0	None	05 H
of 8 Code to EBCDIC-EBCDIC to 4 of	EEC48	3,0	HXCV, SIRTB	W 1A
+ of 8 Code to IBM Card Code-	HOL48	3,0	HXCV, HOLCA,	W 1B
L of 8 Code to Table of Displacements.		3 0	Nono	11: 1 1)
32-Bit Binary Value to IBM Card Code	BIDEC	3,0	None	U4 A
BM Card Code Decimal Value to 32-Bit	DECBI	3,0	Ncne	U4H
Supplement to All Standard Conversions Except Those Involving	Z IPCO	3,0	Any ZIPCO Conversion	U 3I
PITC/8			Table	
MTCA Code Conversion	FEB41, BEB41, F41EB, B41EB, QEB41, Ç41EB	4,0	None	W5 A
Conversion Tables				
BCDIC and PTTC/8	EBPA	3,0	Ncne	U4 K
Card Code Table	HOLL	3.0	None	U4P
		2.0	None	11/10
Console Printer Output Code Table	PRTY	3,0	None	UNY
Console Printer Output Code Table Table of IBM Card Codes	HOLCA	3,0	None	W1C

Figure 24. 1130 Disk Monitor Version 2 System Library (Part 6 of 9)

System Library Programs	Names	Type and Subtype	Subroutines Required	IE Field (73-75)
ZIFCO Conversion Tables				
EPCDIC to Console Printer Code	EBCCP	4,0	None	U41
EECDIC to IEM Card Code	EBHOL	4.0	None	U4J
EECDIC to 1403 Printer Code	EBPT3	4.0	None	U4L
Console Printer Code to EBCDIC	CPEBC	4.0	None	U4D
Console Printer Code to IBM Card Code	CPHOL	4.0	None	114F
Console Printer Code to 1403 Printer	CPPT3	4.0	None	U4F
Code		-		
IBM Card Code to EBCDIC	HLEBC	4,0	None	U4M
IEM Card Code to Console Printer Code	HOLCP	4,0	None	U4O
IEM Card Code to 1403 Printer Code	HLPT3	4,0	None	U4N
1403 Printer Code to EBCDIC	PIJEB	4,0	None	U4S
1403 Printer Code to Console Printer	PI 3CP	4,0	None	• U4R
Code				
1403 Printer Code to IBM Card Code	PTHOL	4,0	None	U4 T
Log Subroutine				
Durmy Log Subrouting				
called by SCAT1, SCAT2, SCAT3	IOLOG, CPLOG	4,0	None	W1J
Interrupt Level Subroutines				
Interrupt Level Zero Subroutine	ILS00	7.0	None	11 1 A
Interrupt Level One Subroutine	ILSO1	7.0	None	U1B
Interrupt Level Two Subroutine (Part	ILS02	7,1	None	U1C
OF Supervisor)	TT CO 3	7 0		
Interrupt Level Intee Subroutine	TT203	7,0	None	UID
of Supervisor)	10204	/,1	None	UIE
Special Interrupt Level Subroutines				
(Restores Index Register 3)				
Interruct Level Zero Subroutino	TTCVA	7 0		
Interrupt Level One Subroutine	TLOAU TLOAU	7,0	None	UIF
Interrupt Level Two Subroutine	TLCV2	7,0	None	01G
Interrupt Level Three Subroutine	TTON2	7,0	None	UIH
Interrupt Level Four Subroutine	TTOVI	7,0	None	011
instrupt lever four suffortine	III SX4	7,0	None	U 1 J
Standard Plot Calls				·
Standard Precision Character	FCHAR	4,0	FSIN, FCOS, FPLOT, FCHRX, FLL, FSTOX,	V1F
Standard Precision Scale	SCALE		FSIC	
Standard Precision Grid	FGRID	4,0 4,0	FRULE FFLCT, FOINT, FADD, FLD,	V10 V1日
Standard Precision Plot	FPLOT	4,0	FSTC, SNR FMCVE, XYFLT, FLCT1	V1I
Extended Precision Character	ECHAR	4,0	ESIN, ECOS, EPLOT, ECHRX, ELD, ESTO, ESTCX	V1A

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Figure 24. 1130 Disk Monitor Version 2 System Library (Part 7 of 9)

System Library Programs	Nanes	Type and Subtype	Subroutines Required	ID Field (73-75)
<u>Standard Plct Calls</u> (continued)				
Extended Precision Scale Extended Precision Grid	SCALE EGRID	4,0 4,0	ERULE EPICT, FOINT, EADD, EID,	V1N V1C
Extended Precision Plot	EPLOI	4,0	ESIC, SAR EMCVE, XYPLT, PLCT1	V1C
Common_Plot_Call	•		·	
Foint Characters	POINT	4,0	PICTI	V 1M
Standard Plot LIBFs				
Standard Precision Annotation	FCHRX, FCHRI, WCHRI	3,0	FLOAT, FMPY IFIX, FADD, FLCX, FINC, XYFLT, FLOTI, FSTCX FLD	V1G
Standard Precision Plot Scaler	FRULE, FMOVE, FINC	3,0	FIDX, FSUBX, FMFYX, FLD, FSTCX, FMPY, IFIX, FADD	V IJ
Extended Plct LIBFs				·
Extended Precision Annotation	ECHRX, ECHRI, VCHRI	3,0	FLOAT, EMPY, IFIX, EACD, ELDX, EINC, XYFLT, FLOTI,	V1B
Extended Frecision Plot Scaler	ERULE, EMOVE, EINC	3,0	ESICX, ELD ELDX, ESUBX, EMPYX, ELD, ESICX, EMPY, IFIX, EADD, ESIC	V1E
Common_Plct_LIBFs				
Pen Mover Interface Interrupt Service	XYFLT PLOTI PLOTX	3,2 3,2 5,0	PLOTI PLOTX IIS03	V1P V1K V1L
Disk_1/0				
Sequential Address	SEÇOP, SEQIO, SEOCL	3,0	DISKZ	W2F
Direct Access	DAOPN, DAIO, DACLS	3,0	DISKZ	W3E
ISAM Load	ISLDO, ISLD, ISLDC	3,0	DISKZ	W3C
ISAM Add	ISADO, ISAD, ISADC	3,0	DISKZ	W3C
ISAM Sequential	ISEQO, ISEIL, ISEQ, ISEQC	3,0	DISKZ	W3E
ISAM Random	ISRDO, ISRD, ISRDC	3,0	DISKZ	W3A

Figure 24. 1130 Disk Monitor Version 2 System Library (Part 8 of 9)

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stem Library Programs	Names	Type and Subtype	Subroutines Required	ID Field (73-75)
G Lecimal Arithmetic				
ld, Subtract, and Numeric Compare ¹	RGADD, RGSUB, RGNCP	3,0	None	W2T
ltiply ¹	RGMLT	3.0	RGBTD, RGDTB	W2S
.vide 1	RGDIV	3,0	None	W2R
ve Remainder ¹	RGMVR	3,0	RGETL	W2Q
nary Conversion ¹	RGBTD, RGDTB	3,0	None	W2 P
G Sterling and Edit				
erling Input Conversion	RGSTI	3,0	RGBTD, RGDTB	W 4B
erling Output Conversion ⁴	RGSTO	3,0	RGETL, RGLIE	W4A
lit ¹	RGEDI	3,0	RGMV2, RGS15	W20
<u>G Move</u>				
com I/O Buffer to Core¹	RGMV1, RGMV5	3,0	Ncne	W 2N
om Core to I/O Buffer¶	RGMV2	3,0	None	W2M
WE Operation ⁴	RGMV3	3,0	None	W2L
WEL Operation [¶]	RGMV4	3,0	None	W2 K
G Compare				
.phameric ¹	RGCMP	3,0	Ncne	W 2J
G Indicators				
est ¹	RGS 11	3,0	None	W 2T
t Resulting On¶	RGS12	3,0	None	W2 H
t on, Set off	RGSI3, RGSI4	3,0	None	W2 G
st for 0 or Blank ¹	RG SI 5	3,0	None	W2 E
G Miscellaneous				
st Zone ¹	RGTSZ	4,0	None	W 2D
nvert to Binary ¹	RGC VB	3,0	None	W2C
ject Time Error ¹	RGERR	4,0	None	W2 B
ank After	RGBLK	3,0	None	W2 A
ternating sequence'	ALISE			
bt distributed to papertape users.				

Figure 24. 1130 Disk Monitor Version 2 System Library (Part 9 of 9)

Appendix B. Errors Detected by the ISS Subroutines

EDDOD	CONTENTS OF ACCUMULATOR	Contents of
	Binary Hexadeo	imal (if any)
1442 Card Read Punch or 1442 Card Punch	· · · · · · · · · · · · · · · · · · ·	
*Last card	000000000000000000000000000000000000000	0
*Feed check *Read check *Punch check Device not ready Lost card indicator on for Read	-[0001000000000000000000000000000000000	1 O
Illegal device (not 0 version) Device not in system Illegal function Word count over +80 Word count zero or negative	-[0001000000001 100	1
Keyboard/Console Printer		
Device not ready Device not in system Illegal function Word count zero or negative Keyboard wants input (TYPEZ only)	0 0 1 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	D 1 2
1134/1055 Paper Tape Reader/Punch		
*Reader not ready Device not ready Illegal device	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	5 0
Illegal function Word count zero or negative Illegal check digit	-[0011000000000001 300	1
2501 Card Reader		
*Last card	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	D
*Feed check }	-[0000000000000000000000000000000000000	1
Device not ready	0 1 0 0 0 0 0 0 0 0 0 0 0 0 0 4 0 0	0
Illegal function Word count over +80 Word count zero or negative	-[01000000000001 400	1
*Data error remaining after 16 attempts (DM2) or 10 attempts (C/PT) during read operation.		Bits 0-3 logical drive number, bits 4-15 working sector address.
*Data error remaining after 16 attempts (DM2) or 10 attempts (C/PT) during write operation.	-[0000000000000000000000000000000000000	Bits 0–3 logical 2 drive number, bits 4–15 working sector address.
*Seek failure remaining after 16 attempts (DM2) or 10 attempts (C/PT).	-[00000000000011 0003	Bits 0-3 logical drive number, bits 4-15 working sector address.
*Attempt to read or write above sector address 1599 (disk over- flow).	- <u>[</u> 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	
• Device not ready.	0101000000000000 500	0

130

. '

	EKKUK						. J'					- Extension								
				Binary Hexadecimal (if an				(if any)												
Disk (continued)																				
Illegal dev device noi to write ir word coun starting ac (DISK1 an	vice, invalid function, t on system, attempt i file protected area, t zero or negative or Idress over 1599 d DISKZ only).	-{0	1	0 1	0	0	0 (00	0 0	0	0	0 (0 0) 1		5	0	0 1	l	
Write sele	ct/Power unsafe.	0	1	0 1	0	0	0	0 0	0 0	0	0	0	0 1	0		5	0	0 2	2	
Data error operation 4 16 attempt attempts (attempt to above sect (disk overf during the Monitor ca DISKZ on	during read/write or seek failure after is (DM2) or 10 C/PT) or on an read or write or address 1599 Iow). Error occurred processing of a ill. (DISK1 and iy).	{0	1	0	1 0	0	0	0	0 0) 0	0	0	1	0 ()	5	0	0 :	3	Bits O–3 logical drive number, bits 4–15 working, sector address, except for disk overflow.
Data error attempts (remaining after 16	[0	1	0	10	0	0	0	0 0	0	0	0	1	0 (5	0	0	4	Bits 0-3 logical drive number, bits 4-15 working sector address + 1.
1132 Printer			_	_				_	_				•			_		_	_	
*Channe *Channe Device Illegal Word co	I 9 detected I 12 detected not ready or end of forms function	0	001	0 0 1,			0000	0000					0 1 0 0	000	0	0 0 6 6	00000	000	3 4 0	
Word c	ount zero or negative)				•															
Plotter																				
Plotter n	ot ready	0	1	1	1 (0 (0	0	0	0 (0 0	0	0	0	0	7	0	0	0	
illegal f Word co	unction unt zero or negative		1	1	1 (0 0	0	0	0	0 (0 0	0	0	0	١	7	0	0	1	
1403 Printer																				
*Ring ch *Sync cl	eck	10	0	0	0 0	0 0	0	0	0	0 (0 0) 0	0	0	1	0	0	0	1	
*Parity (*Channe	check) I 9 detected	0	0	0	0 0	0 0	0	0	0	0	0 0	0	0	1	1	0	0	0	3	
*Channe Device	not ready or end of forms	1	0	0	1 (0 0	0	0	Ő	0) 0	ò	ŏ	ŏ	9	0	0	0	
Word c	ount over +60	י]+-	0	0	1 (0 0	0	0	0	0	0 0) 0	0	0	1	9	0	0	1	
PRNZ only { Ring cl Sync cl Parity :	heck	- <u>-</u> [1	0	0	1 (0 0	0	0	0	0	0 0	0	0	1	0	9	0	0	2	
Optical Mark P	oge Reader			_										_			_	_		
Master n Timing n	nark nark error)	0	0	0	0	0 0 0 0	0	0	0.	0	00) () (1	0	0	0	0	2	
Read err Hopper (or ∮ empty		0	0	0	0 0	0	0	0	0	0 0	0	0	1	1	0	0	0	3	
Documer Device	nt selected not ready	0	0 0	0 1	0	0 0 0 0	0	0 0	0	0	0 (0 () () (10	0 0	0	0 A	0	0	4 0	
Illegal f Feed ch Feed ch proces	unction eck, last document processed eck, last document not sed		0 0 0	1 1 1	0	000000000000000000000000000000000000000) 0) 0) 0	0	0 0 0	0 0 0				0	1 0 1	A A A		0 0 0	1 2 3	
•····		ľ	2				-	-	~				5					-	-	1

of the I/O device. All other errors cause a branch to location /0029 on the C/PT system or to one of the traps in the DM2 system, at \$PRET, \$PST1, \$PST2, \$PST3, or \$PST4. If the error WAIT is in the preoperative Error Trap the address of the LIBF or the CALL is in location \$PRET.

Appendix C. Subroutine Action on Return from a **User's Error Subroutine**

Error Code	Condition	Subroutine Action
1442 Card Read Punch or 1442 Card Punch 0000 000 11	If function is PUNCH otherwise If Accumulator is 0 otherwise	Eject card and terminate Terminate immediately Terminate immediately C/PT System: Loop until 1442 is ready, then reinitiate operation DM2 System: WAIT at \$PST4, clear 1442 with NPRO key, assure that 2nd card run out is in correct pre-punched form (1st card con- tains punch error), replace cards in 1442 hopper, press 1442 start key, and press PROGRAM START.
2501 Card Reader 0000 00011	If Accumulator is 0 otherwise	Terminate Terminate immediately WAIT at \$PST4 until 2501 is readied and PROGRAM START pressed
1134/1055 Paper Tape Reader/Punch 0004,0005	If Accumulator is 0 otherwise	Terminate immediately Check again for device ready
Disk 0001, 0002, and 0003 0004	If Accumulator is 0 otherwise	Terminate immediately Retry 10 times (C/PT system), or 16 times (DM2 system) Terminate and go to exit
1132 Printer 0003, and 0004	If Accumulator is 0 otherwise	Terminate immediately Skip to channel 1 and then terminate
1403 Printer 0001 0003, and 0004	If Accumulator is 0 otherwise If Accumulator is 0 otherwise	Terminate immediately Check for device ready and reinitiate the operation Terminate immediately Skip to channel 1 and then terminate
1231 OMPR 0001 00021 0003	If Accumulator is 0 otherwise If Accumulator is 0 otherwise	Continue normal processing Use contents of Accumulator as new address of I/O area Terminate immediately Check device ready, then reinitiate operation Terminate
0004*	II ACCUMULATOR IS U otherwise ²	Check device ready, then reinitiate operation

1

Assumes operator intervention
User must provide a WAIT in his error subroutine to allow him to remove the sheet from the select stacker, place the sheet back in the hopper, and make the 1231 ready.

Appendix D. Character Code Chart

	EBCDIC		81	M Card C	Code				1132	PTTC/8	Console	1403
Ret No.	Binary	Hex		Rows	7 1	Hex	Graph	Names	EBCDIC	U-Upper Case L-Lower Case	Hex	Printer Hev
0 1 2 3 4 5* 6* 7* 8 9 10 11 12 13 14 15	0123 4387 0000 0000 0001 0010 0010 0101 0100 0101 0110 1001 1001 1010 1011 1100 1111 1110	00 01 02 03 04 05 06 07 08 09 0A 09 0A 0B 0C 0D 0E 0F	12 11 12 12 12 12 12 12 12 12 12 12 12 12 12 12 12 12 12 12 12 12 12 12 12 12 12 12 12 12 12 12 12 12 12 12	0 9 8 9 9 9 9 9 9 9 9 9 9 9 9 9 8 9 8 9 8	1 1 2 3 4 5 6 7 1 2 3 4 5 6 7	B030 9010 8810 8210 8090 8050 8030 8030 8830 8830 8830 8830 883	NUL PF HT LC DEL	Punch Off Horiz.Tab Lower Case Delete	Subset Hex	6 D (U/L) 6E (U/L) 7F (U/L)	41 ①	
16 17 18 19 20* 21* 22* 23 24 25 26 27 28 29 30 31	0001 0000 0001 0010 0011 0100 0101 0110 0111 1000 1001 1010 1011 1100 1111	10 11 12 13 14 15 16 17 18 19 1A 19 1A 1B 1C 1D 1E 1F	12 11 11 11 11 11 11 11 11 11 11 11 11 11	98 99 99 99 99 99 98 88 98 88 98 88 98 88 98 88 98 88 98 88 98 88 98 88 98 88 99 88 99	1 2 3 4 5 6 7 1 2 3 4 5 6 7	D030 5010 4810 4410 4210 4090 4050 4050 4030 5030 4830 4430 4230 4130 4080 4070	RES NL BS IDL	Restore New Line Backspace Idle		4C (U/L) DD(U/L) 5E (U/L)	05 Ø 81 ③ 11	
32 33 34 35 36 37* 38* 39 40 41 42 43 44 45 46 47	0010 0000 0001 0010 0011 0100 0101 0110 0111 1000 1001 1010 1001 1010 1011 1100 1101 1110	20 21 22 23 24 25 26 27 28 29 2A 29 2A 29 2A 2D 2E 2F	11	0 9 8 0 9 0 9 0 9 0 9 0 9 0 9 0 9 8 0 9 8 0 9 9 0 9 9 0 9 0 9 9 8 0 9 9 8 0 9 9 8 0 9 9 8 0 9 9 8 0 9 9 8 0 9 9 8 0 9 9 8 0 9 9 8 0 9 9 8 0 9 9 8 0 9 9 8 0 9 9 8 0 9 9 8 0 9 9 8 0 9 9 8 0 9 9 8 0 9 9 8 8 0 9 8 8 0 9 8 8 0 9 8 8 0 9 8 8 0 9 8 8 0 9 9 8 8 0 9 9 8 8 8 0 9 9 8 8 8 0 9 9 8 8 8 0 9 9 8 8 8 0 9 9 8 8 8 0 9 9 8 8 8 0 9 8 8 8 8 0 9 8 8 8 8 9 8 8 9 8 8 9 8 8 9 8 8 9 8 8 9 8 8 8 9 8 8 9 8 8 9 8 8 9 8 8 9 8 8 9 8 8 9 8 9	1 1 2 3 4 5 6 7 1 2 3 4 5 6 7	7030 3010 2810 2410 2210 2050 2050 2050 2030 3030 2830 2430 2230 2130 2080 2070	BYP LF EOB PRE	Bypass Line Feed End of Block Prefix	-	3 D (U/L) 3 E (U/L)	03	· ·
48 49 50 51 52 53* 54* 55 56 57 58 59 60 61 62 63	0011 0000 0001 0010 0010 0100 0101 0100 1001 1010 1011 1010 1101 1101	30 31 32 33 34 35 36 37 38 39 34 39 3A 3D 3E 3F	12 11	0 9 8 9 9 9 9 9 9 9 9 9 9 9 9 9 8 9 8 9 8	1 1 2 3 4 5 6 7 1 2 3 4 5 6 7	F030 1010 0810 0410 0110 0050 0050 0030 1030 0830 0430 0430 0230 0130 0080 0070	PN RS UC EOT	Punch On Reader Stop Upper Case End of Trans.		0 D (U/L) 0 E (U/L)	09 ④	

NOTES: Typewriter Output

1) Tabulate 2 Shift to black

Carrier Return
 Shift to red

* Recognized by all Conversion subroutines Codes that are not asterisked are recognized only by the SPEED subroutine

	EBCDIC			IBM Card Code						1132	PTTC/8	Console	1403
No.	Binary 0123 45	567 He	12	2 11	Row: 0	s 98	7-1	Hex	Graphics and Control Names	EBCDIC Subset Hex	Hex U-Upper Case L-Lower Case	Printer Hex	Printer Hex
64* 65 66 67 70 71 72 73 74* 75* 76* 75* 78* 79*	0100 00 000 000 01 01 01 01 01 01 01 01 01 01	000 40 001 41 010 42 011 43 100 44 101 42 111 43 100 44 111 45 110 46 111 47 000 48 001 49 010 40 011 48 100 40 110 46 111 47	12 12 12 12 12 12 12 12 12 12 12 12 12 1	nc	0 0 0 0 0 0 0 0 0	hes 7 7 7 7 7 7 7 7 7 7 7 7 7 7 7 7 7 7 7	1 2 3 4 5 6 7 1 2 3 4 5 6 7	0000 B010 A810 A410 A210 A050 A050 A050 A050 B020 B420 B420 B420 B120 B0A0 B060	blank ¢ . (period) < (+ I (logical OR)	40 * * 48 4D 4E	10 (U/L) 20 (U) 68 (L) 02 (U) 19 (U) 70 (U) 38 (U)	21 02 00 DE FE DA C6	7F 6E 57 6D
80* 81 82 83 84 85 86 87 88 89 90* 91* 92* 93* 94* 95*	0101 00 00 00 00 01 01 01 01 01 01 01 01 01 0	000 50 001 51 010 52 011 53 000 54 101 55 110 56 111 57 000 58 001 59 010 5A 011 5B 000 5C 001 5D 111 57 000 5C 010 5A 111 5B 100 5C 111 5F	12 12 12 12 12 12 12 12 12	11 11 11 11 11 11 11 11 11 11 11 11 11		88888888888888888888888888888888888888	1 2 3 4 5 6 7 1 2 3 4 5 6 7	8000 D010 C810 C210 C110 C090 C050 C030 5020 4820 4820 4820 4420 4120 40A0 4060	& ! \$; ; (logical NOT)	50 58 5C 5D	70 (L) 58 (U) 58 (L) 08 (U) 1A (U) 13 (U) 68 (U)	44 42 40 D6 F6 D2 F2	15 62 23 2F
96* 97* 98 99 100 101 102 103 104 105 106 107* 108* 109* 110* 111*	0110 00 00 00 01 01 01 01 01 01 01 01 01 01 10 10	000 600 101 61 1010 62 111 63 100 64 101 62 111 63 100 64 111 65 100 64 111 67 000 68 001 69 010 6A 011 6B 100 6C 101 6B 100 6C 111 6F	12	11 11 11 11 11 11 11 11		8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8	1 2 3 4 5 6 7 1 3 4 5 6 7	4000 3000 6810 6410 6210 6050 6050 6050 6050 6050 2000 2420 2220 2120 2220 2120 20A0 2060	- (dash) / (comma) • % (underscore) ?	60 61 6B	40 (L) 31 (L) 38 (L) 15 (U) 40 (U) 07 (U) 31 (U)	84 BC 80 06 BE 46 86	61 4C 16
112 113 114 115 116 117 118 119 120 121 122* 123* 124* 125* 126* 127*	0111 00 00 00 00 01 01 01 01 01 01 01 10 10 1	000 70 001 71 010 72 011 73 000 74 101 75 111 75 001 76 101 75 001 76 001 79 001 70 001 70 100 7C 100 7C 111 7F	12 12 12 12 12 12 12 12 12	11 11 11 11 11 11	000000000000000000000000000000000000000	8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8	1 2 3 4 5 6 7 1 2 3 4 5 6 7	E000 F010 E810 E410 E10 E100 E050 E050 E050 E050 E030 1020 0820 0420 0420 0420 0420 0420 0420 0	: # @ ' (apostrophe) = "	7D 7E	04 (U) 08 (L) 20 (L) 16 (U) 01 (U) 08 (U)	82 C0 04 E6 C2 E2	08 4A

** Any code other than those defined for the 1132 will be interpreted by the PRNT1 subroutine as a blank.
	EBCDIC		IBM Card Code			1132	PTTC/8	Console	1403
Ref No.	Binary 0123 4567	Hex	Rows 12 11 0 9 8 7-1	Hex	Hex Graphics and Control Names		Hex U-Upper Case L-Lower Case	Printer Hex	Printer Hex
128 129 130 131 132 133 134 135 136 137 138 139 140 141 142 143	1000 0000 0001 0010 0011 0100 0101 0110 0111 1000 1001 1011 1010 1011 1100 1011 1110	80 81 82 83 84 85 86 87 88 87 88 87 88 87 88 80 88 80 88 85 88 85 88 85 88 85 88 88 85 88 88	$\begin{array}{cccccccccccccccccccccccccccccccccccc$	8020 8000 A800 A400 A100 A080 A040 A020 A010 A820 A420 A420 A120 A0A0 A060	a b c d e f g h i				:
144 145 146 147 148 149 150 151 152 153 154 155 156 157 158 159	1001 0000 0001 0010 0010 0101 0100 0111 1000 1001 1010 1011 1100 1101 1110	90 91 92 93 94 95 96 97 98 97 98 97 98 97 98 97 98 97 98 97 98 97 98 97 98 97 98 97 98 97 98 97 98 97 97 97 97 97 97 97 97 97 97 97 97 97	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	D020 D000 C800 C400 C100 C080 C040 C020 C010 C820 C420 C420 C420 C120 C120 C0A0 C060	i κ Ε η ο β q r				
160 161 162 163 164 165 166 167 168 169 170 171 172 173 174 175	1010 0000 0001 0010 0010 0100 0101 0110 0111 1000 1001 1010 1011 1100 1101 1110	A0 A1 A2 A3 A4 A5 A6 A7 A8 A7 A8 A7 A8 AC AD AE AF	$\begin{array}{cccccccccccccccccccccccccccccccccccc$	7020 7000 6800 6400 6200 6100 6080 6040 6020 6040 6420 6420 6420 6420 642	s t v w x y z				
176 177 178 179 180 181 182 183 184 185 186 187 188 189 190 191	1011 0000 0001 0010 0010 0101 0101 0110 0111 1000 1001 1010 1011 1100 1101 1110	B0 B1 B2 B3 B4 B5 B6 B7 B8 B8 B8 B8 B8 BB BB BE BF	$\begin{array}{cccccccccccccccccccccccccccccccccccc$	F020 F000 E800 E400 E000 E080 E040 E020 E010 E820 E820 E820 E820 E820 E120 E040 E040 E040					

		EBCDIC				BM	Card	Code			1132	PTTC/8	Console	1403
Ref No.	Bin	ary	Hex			Rows	_		Hex	Graphics and Control Names	EBCDIC	U-Upper Case	Printer Hey	Printer
	0123	4567		12		0 9	8	7-1			Subset Hex	L-Lower Case		
192 193* 194* 195* 196* 197* 198* 199* 200* 201* 202 203 204 205 206 207	1100	0000 0001 0010 0101 0100 0101 0110 1000 1001 1010 1010 1100 1101 1110	C0 C1 C2 C3 C4 C5 C6 C7 C8 C7 C8 C7 C8 C7 C8 C7 C8 C7 C8 C7 C6 C7 C7 C7 C7 C7 C7 C7 C7 C7 C7 C7 C7 C7	12 12 12 12 12 12 12 12 12 12 12 12 12 1		0	8 9 8 9 8 9 8 9 8 9 8 9 8 9 8 9 8 9 8 9	1 2 3 4 5 6 7 2 3 4 5 6 7	A000 9000 8800 8400 8100 8080 8040 8020 8040 8020 8010 A830 A430 A430 A430 A430 A430 A430 A080 A070	(+ zero) A B C D E F G H I	C12 C24 C58 C8	61 (U) 62 (U) 73 (U) 64 (U) 75 (U) 67 (U) 67 (U) 67 (U) 79 (U)	3C or 3E 18 or 1A 1C or 1E 30 or 32 34 or 36 10 or 12 14 or 16 24 or 26 20 or 22	64 25 26 67 68 29 2A 68 2C
208 209* 210* 211* 212* 213* 214* 215* 216* 217* 218 219 220 221 222 223	1101	0000 0001 0010 0100 0101 0100 0101 0110 1001 1001 1011 1100 1101 1110 1111	D0 D1 D2 D3 D4 D5 D6 D7 D8 D7 D8 D7 D8 D7 D8 D7 D7 D7 D7 D7 D7 D7 D7 D7 D7 D7 D7 D7	12 12 12 12 12 12 12	11 11 11 11 11 11 11 11 11 11 11	°0 \$ \$ \$ \$ \$ \$ \$ \$ \$ \$ \$ \$ \$ \$ \$ \$ \$ \$	8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8	1234567 234567	6000 5000 4800 4400 4100 4080 4040 4020 4010 C830 C430 C230 C130 C0B0 C070	(- zero) J K L M N O P Q R	D1 D2 D3 D4 D5 D6 D7 D8 D9	51 (U) 52 (U) 43 (U) 54 (U) 45 (U) 57 (U) 58 (U) 49 (U)	7C or 7 E 58 or 5A 5C or 5E 70 or 72 74 or 76 50 or 52 54 or 56 64 or 66 60 or 62	58 19 1A 5B 1C 5D 5E 1F 20
224 225 226* 227* 238* 230* 231* 232* 233* 234 235 236 237 238 239	1110	0000 0001 0010 0011 0100 0101 0101 0110 0101 0101 1001 1010 1010 1011 1101 1110 1111 1110	E0 E1 E2 E3 E4 E5 E6 E7 E8 E9 EA E0 ED EE EF		11 11 11 11 11 11	0 5 0 0 0 0 0 0 0 5 0 5 0 5 0 5 0 5 0 5	8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8	2 1 2 3 4 5 6 7 2 3 4 5 6 7	2820 7010 2800 2400 2200 2040 2040 2020 2010 6830 6430 6430 6430 6130 6080 6070	S T U V W X Y Z	E2 E3 E4 E5 E6 E7 E8 E9	32 (U) 23 (U) 34 (U) 25 (U) 26 (U) 37 (U) 38 (U) 29 (U)	98 or 9A 9C or 9E B0 or 82 B4 or 86 90 or 92 94 or 96 A4 or A6 A0 or A2	0D 4F 10 51 52 13 54
240* 241* 242* 243* 244* 245* 246* 247* 248* 249* 250 251 252 253 254 255	1111	0000 0001 0010 0011 0101 0101 0101 0110 0111 1000 1001 1011 1010 1011 1011 1010 1011 1100 1101 1110 1111 1111	F0 F1 F2 F3 F4 F5 F6 F7 F8 F7 F8 F9 FA FB FC FD FE FF	12 12 12 12 12 12	11 11 11 11 11 11	0 50 50 50 50 50 50 50 50 50 50 50 50 50	8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8	1 2 3 4 5 6 7 2 3 4 5 6 7	2000 1000 0800 0200 0100 00400 0040 0040 0020 0010 E830 E430 E430 E130 E130 E080 E070	0 1 2 3 4 5 6 7 8 9	F0 F1 F2 F3 F4 F5 F6 F7 F8 F9	1A (L) 01 (L) 02 (L) 13 (L) 04 (L) 15 (L) 16 (L) 07 (L) 08 (L) 19 (L)	C4 FC D8 F0 F4 D0 D4 E4 E0	49 40 01 43 04 45 46 07 08

Appendix E. Core Requirements of Subroutines

Communications Adapter subroutine core requirements are listed in the publication <u>IBM</u> <u>1130 Synchronous Communications Adapter Subroutines</u>, GC26-3706. MTCA subroutine core requirements are listed in the publication <u>IBM</u> <u>1130</u> <u>Computing System</u>, <u>Multiple Terminal</u> <u>Communications Adapter (MTCA), Input/Output Control System (IOCS)</u> <u>Subroutines</u>, GC34-0015. <u>1627</u> Plotter subroutine core requirements are included in the publication <u>IBM</u> <u>1130/1800</u> <u>Plotter Subroutines</u>.

Standard	Extended	Standard	Extended
FADD/FADDX FSUB/FSUBX	EADD/EADDX ESUB/ESUBX 98	<u>C/PT System</u>	C/PT System
FMPY/FMPYX 52	EMPY/EMPYX 46	WARI/WARIX 32	VARI/VARIX 32
FDIV/FDIVX 86	EDIV/EDIVX 78	WIAR/WIARX 36	VIAR/VIARX 36
FLD/FLDX 1 54	ELD/ELDX) 46	WIF 26	VIF 26
FSTO/FSTOX (ESTO/ESTOX	WIIF 24	VIIF 24
FLOAT 10	10	WGOTO 22	VGOTO 22
IFIX 40	40	WFIO/WIOI/WIOAI/	VFIO/VIOI/VIOAI/
NORM 42	42	WIOF/WIOAF/	VIOF/VIOAF/
FSBR/FSBRX 24	ESBR/ESBRX 24	WIOFX/WCOMP/ 854	VIOFX/VCOMP/ 864
FDVR/FDVRX 28	EDVR/EDVRX 28	WWRT/WRED/	VWRT/VRED/
SNR 8	8	WIOIX	VIOIX
FABS/FAVL 12	EABS/EAVL 12)
IABS 16	16		
FGEPT 22	EGETP 22		
FARC 34	34		
XMDS 28			
	68		
XSQR 52	52		
XMU 00	66		
	/4		
FCOS/FCOSN 118	ESCOS/ESCOSN 138		
FATN/FATAN 130	EATN/EATAN 148		
FSQR/FSQRT 70	ESQR/ESQRT 76		
FLN/FALOG 136	ELN/EALOG 148		
FEXP/FEXPN 118	EEXP/EXPN 140		
FAXI/FAXIX 78	EAXI/EAXIX 82		
FAXB/FAXBX 54	EAXB/EAXBX 54	DM2 System	DM2 System
FTNH/FTANH 54	ETNH/ETANH 46	SDFIO/SDAF/SDAI/	
FDTB (dec. to bin.) 446	446	SDCOM/SDF/SDFX/ 694	694
DMTD0/DMTX0 412	412	SDI/SDIX/SDRED/	
DMPD1/DMPX1 520	520	SDWKI '	
DMP80 102	102	SUFIND /8	/8
DATSW 34	34	SFAR/SFARA 32	SEAR/SEARX 32
DVCHK 16	16		1100
FCTST 30	30	SIOF/SIOAF/SIOFX/ 1194	1190
LOAD 138	138	SCOMP/SWRI/SRED/	5515 00
OVERF 18	18		JEIF 28
SLITE, SLITT 70	70	SCOTO 22	22
TSTOP 6	6	SIAP/SIAPY 24	30
TSTRT 6	6		24
ISIGN 24	24	JIIF 24	/36
FSIGN 34	ESIGN 34	UFIC 738	

Figure 25. Core Requirements of Arithmetic and Functional Subroutines

Subroutines	No. Core Locations		Uses
	DM2 System	C/PT System	(DM2 System)
CARDO CARDI READO READI PNCHO PNCHI OMPRI PAPTI PAPTN DISKO DISK1+ DISKN+ WRTYO TYPEO PLOTI PRNTI PRNTI ILSOO ILSOI ILSOI ILSOI ILSOI ILSO3 ILSO4+ ILSX2 ILSX3 ILSX4	254 258 96 110 206 218 336 226 306 - 418 688 124 278 186 424 308 22 26 17 28 32 26 17 28 32 24 30 24 34 44	242 246 - - - - 254 294 356 620 808 124 296 216 386 - 18 18 18 18 18 18 18 30 - - - - -	ILS00, ILS04 ILS04, ILS04 ILS04, ILS00 ILS04, ILS00 ILS04, ILS00 ILS04 ILS04 ILS04 ILS02 ILS02 ILS02 ILS04 ILS04 ILS03 ILS01 ILS04
SPIRO SPIRI SPIRN		48 62 62	
FLIPR PAUSE STOP	102 22 12	- 22 8	DISKZ, DISK1, DISKN
SUB SC SUBIN TTEST/TSET DISKZ+ CARDZ PAPTZ PRNTZ TYPEZ WRYTZ READZ PNCHZ PRNZ HOLEZ GETAD EBCTB HOLTB	30 32 16 238 176 226 218 106 62 58 66 186 64 186 64 16 60 54	30 32 16 - 80 202 176 82 66 54 14 54 54	ILS02 ILS04, ILS00 ILS04 ILS01 ILS04 ILS04 ILS04 ILS04, ILS00 ILS04
SYSUP FSLEN/FSYSU	1338 535	-	FSLEN/FSYSU DISKZ

* Part of Resident Monitor

Figure 26. Core Requirements of Miscellaneous and ISS Subroutines

Conversion Subroutines	N Ci Loca DM2 System	Uses	
BINDC DCBIN BINHX HXBIN HOLEB HOLPR EBPRT PAPEB PAPPR ZIPCO SPEED HOLL EBPA PRTY EBCCP EBHOL EBPT3 CPEBC CPHOL CPPT3 HLEBC HOLCP HLPT3 PT3CP PT3CP PT3CP PT3CP PT3CP PT40L	72 88 44 66 134 100 102 246 244 192 162 334 80 80 128 128 128 128 128 128 128 128	72 88 44 66 134 100 102 246 244 192 - - - - - - - - - - - - -	HOLL, EBPA HOLL, PRTY EBPA, PRTY EBPA, HOLL EBPA, PRTY

Figure 27. Core Requirements of Conversion Subroutines

	No.	
Subroutines	Core Locations	Uses
	450	DICKZ
SEQUP, SEQIO, SEQUE	458	DISKZ
ISLOO ISLD ISLDC	639	DISKZ
ISADO, ISAD, ISADC	1799	DISKZ
ISEQO, ISETL, ISEQ,	721	DISKZ
ISEQC		
ISRDO, ISRD ISRDC	460	DISKZ
RPG Decimal Arithmetic		
RGADD, RGSUB, RGNCP	464	
RGMLT	320	RGBTD, RGDTB
RGDIV	815	
	118	RGBTD
KGBID, KGDIB	112	
RPG Sterling and Edit		
RGSTI	258	RGBTD, RGDTB
RGSTO	464	RGBTD, RGDTB
RGEDT	315	
RPG Move		
RGMV1, RGMV5	148	
RGMV2	179	
RGMV3	48	
RGMV4	86	
RPG Compare		
RGCMP	82	
RPG Inidcators		
RGSI1	68	
RGS12	78	
RGSI3, RGSI4	40	
RGSI5	92	
RPG Miscellaneous		
RGTSZ	72	
RGCVB	86	
RGERR	70	
RGBLK	58	
ALTSE (user-written)	(variable)	

Figure 28. Core Requirements of RPG Subroutines (DM2 only) Execution times for the Synchronous Communications Adapter subroutines are listed in the adapter subroutine manual.

CONVERSION SUBROUTINES (see Figure 29).

Basic Definitions

- 1. All times are based on 3.6-µsec instruction cycle.
- The table ordering for codes is as follows (except SPEED)

Standard set: blank, +, &, -, 0-9, A-Z, other special

Extended set: stanlard, non-FORTRAN special, control

- 3. Maximum number of characters checked varies with the set.
 - Standard set Except SPEED: 49 SPEED only: 16

Extended set Except SPEED: 74 SPEED only: 45

4. Conversion times given are

Best time: Found as first character in set

Worst time, standard set: Found as last character in set

Worst time, extended set: Not found in set

5. Time per character is best time, plus table look-up time multiplied by the number of characters to be skipped.

Example: If best = 211, look-up = 45.5 and character is fourth in table (-) Then, character time = 347.5 = 211 + 3(45.5)

1130 ISS TIMES (see Figures 30 and 31)

Basic Definitions

 Only CPU time used by ISS (including transfer vector BSC L) and ILS (including forced BSI I) is given. All the remaining time, minus cycle steals, is available to the user.

2. ILS time is included in ISS interrupt processing calculations

C/PT System

ILS00 - CARDO (col), CARD1 (col)

ILSO1 - PRNT1

ILS02 - DISKO, DISK1, DISKN

ILSO3 - PLOT1

ILS04 - CARDO (op complt), CARD1 (op complt) WRTY0, TYPE0, PAPT1, PAPTN

		Time,	ľ		
	Initial—		W	orst	Table
Subroutine	ization	Best	Std. Set	Extd. Set	Look- Up
BINDC DCBIN BINHX HXBIN	1130 1110 620 760	- - -	- - -		- - -
HOLPR EBPRT	430 420	211 207	2395 2487	3533 3675	45.5 47.5
HOLEB EBCDIC output EBCDIC input	550 550	159 161	2343 2441	3481 3629	45.5 47.5
SPEED Packed EBCDIC output Unpacked EBCDIC output Packed EBCDIC input Unpacked EBCDIC input	250 270 240 240	270 260 394 404	- 1594 1604	- 3914 3924	- 80.0 80.0
ZIPCO (DM2 only) All codes except IBM Card Code	270	270	-	-	-
IBM Card Code input IBM Card Code output	270 270	374 435	-	-	-
PAPPR Per shift char. input Per graphic char. input Per control char. input	580	180 427 407	- 2707 2687	- 3895 3875	- 47.5 47.5
PAPHL PTTC/8 input Per shift char, input Per graphic char, input Per control char, input PTTC/8 output Per control char, output Per shift cranshic char, output Per shift cranshic char, output	490 490	180 306 296 266 316	2482 2472 - 2492 2622	- 3870 3860 3830 3880 4010	- 49.5 49.5 49.5 49.5
PAPEB PTTC/8 input Per graphic char. input Per graphic char. input Per control char. input PTTC/8 output Control char. output Per graphic char. output Per shift/graphic char. output	440 440	190 366 386 296 346 476	2522 2542 2562 - 2522 2652	- 3930 3950 3860 3910 4040	49.5 49.5 49.5 49.5 49.5 49.5

Figure 29. Execution Times of Conversion Subroutines

DM2 System

ILS00 - CARD0 (col), CARD1 (col) PNCH0 (col), PNCH1 (col) ILS01 - PRNT1 ILS02 - DISK1, DISKN ILS03 - PLOT1, PLOTX ILS04 - CARD0 (op complt), CARD1 (op complt), PNCH0 (op complt), PNCH1 (op complt), READ0, READ1, WRIY0,

TYPEO, PAPT1, PAPTN, PAPTX, PRNT3, OMPR 1

Note: In the DM2 system, the Z subroutines are considered to be ISSs and therefore use the appropriate ILSs, e.g., PRNTZ uses ILS01.

3. All times are based on a 3.6-µsec instruction cycle.

Subjoutine and Function	Times (µsec) (n word count)	Subroutine and Function	Times (μsec) (n = word count)
ILSOO	112	PRNTI	100
ILSOI	134	Priot	
ILS02	112	l rim	44142 + 5971.2 (n-1)*
	148		*subtract 11.4 for each word where 1 char. does not ma*ch; 22.8 where heat share do not
LARDU Test	165		motch,
Read	14930 + 38.5 (n)		
Punch	763 + 185 (n)	Print Numeric .	25950 + 2736.8 (n-1)
Feed	605		+268 ×
Sel. Stock. CARDI	290		x = no, idle cycles before 1st numeric char, on wheels is
Test	165		reached
Read	14972 + 38.5 (n)		
Punch	800 + 190 (n)	Control	700
Feed	640	Single space	708
Sel. Stack.	325	Triple space	1288
WRIYO		Skip to channel 12	676*
Test	165	Skip to channel 1	936*
Print	228 + 734 (n)		· · · · · · · · · · · · · · · · · · ·
τγρεο	1/5		*odd 208 for each channel crossed before correct one reached
lest Baard a stat	100 495 + c (925 + 49 5) + 200 = +	DISK0	
Keod print	1595 h + 1224 c	Test	178
	1373 0 1224 0	Read	1492
	€ = sum of char. times for each	Write	
	graphic	Without RBC	1778
	y no. char. skipped in table	With RBC	2050
	look-up	Sook	1062
	a EUM character	1 to center	1076
	c = backspace character	By addi	1502
		DISKI	
Print	344 + 920 (n)	Test	178
PAPTI		Read	900 + 760 x + 478 y
Test	152		
Reod	432 + 808* (n)		x = no. sectors y = no. seeks after 1st sector
	*add +112 if check	Write	
Punch	480 + 680* (n)	Without RBC Write	1292 + 660 x + 822 y
	*add +96 if check	With RBC Write Imm	1562 + 1098 x + 908 y 660 + 622 x + 476 y
PAPTN	· · · · ·	Seek	l
Test	176	1 to center	1072
Read	408 + 952* (n)	By addr	1468
	add +112 if check	Test	178
Punch	464 + 840* (n)	Keod	408 + 652 x + 1012 y
BLOT!	*add +64 if check		x = no. sectors y = no. seeks after 1st sector
FLUTI Test	130	Write	
Print	418 = if char is 0-9	Without RBC	1516 + 610 x + 926 y
	472 = if char is A	Write	1700 1000 1117
	624 = if char is B	With KBC	1/28 + 1022 x + 11/8 y
	698 + 752 = if char is C	Seek	040 + 000 x + 282 y
	224 per dup, of	1 to center	1076
	motion	By addr	1478
			1

Figure 30. Execution Times of 1130 ISS (C/PT System)

Subroutine and Function	Times (µsec) (n = word count)	Subroutine and Function	Times (μsec) (n = word count)
ILS00 ILS01 ILS02 ILS03 ILS04	112 134 102 112 163	PLOTI (Cont'd)	678 + { 752 = if char is C 224 = per dup. of previous pen motion
CARDO Test Read Punch Feed Sel, Stack.	165 14930 + 38.5 (n) 763 + 185 (n) 605 290	PRNT1 Test Print	188 44142 + 5971.2 (n-1)* *subtract 11.4 for each word where 1 char. does not match; 22.8 where both char. do not match.
CARDI Test Read Punch Feed Sel. Stack.	165 14972 + 38.5 (n) 800 + 190 (n) 640 325	Print Numeric	25950 + 2736.8 (n-1) + 268 x x = no. idle cycles before 1st numeric char. on wheels is
READO Test Read Feed READI Test Read Feed	173 546 523 173 576 553	Control Single space Double space Triple space Skip to channel 12 Skip to channel 1	708 998 1288 676* 936* *add 208 for each channel crossed
PNCHU Test Punch Feed PNCH1 Test Punch Feed WRTY0	165 763 + 185 (n) 605 165 800 + 190 (n) 640	PRNT3 Test Print Control Single Space Double Space Triple Space Skip to channel 12 Skip to channel 12	before correct one reached 183 3743 + 45 (n-1) 785 6746 12704 817 817
Test Print TYPE0 Test Read print	165 228 + 734 (n) 165 685 + € (825 - 48.5y) + 390 o + 1595 b + 1224 c	OMPR1 Test Feed Read	227 710 805 + 286 x c c = no. of chars. programmed to be read
Print	<pre></pre>	Disconnect Sel. Stack. DISK1 Test Read	506 495 158 1021 + 491 x + 1226 y x = no. sectors
PAPT 1 Test Read	152 432 + 808* (n) *add + 112 if check	Write Without RBC Write With RBC Write Imm	y = no. seeks after 1st sector 1035 + 491 x + 1226 y 1829 + 982 x + 2452 y 689 + 491 x + 489 y
Punch	480 + 680* (n) *add + 96 if check	Seek 1 to-center By addr	1843 2056
PAPTN Test Read	176 408 + 952* (n) *add + 112 if check	DISKN Test Read	244 1500 + 725 x + 1973 y x = no. sectors
Punch	464 + 840* (n) *add + 64 if check	Write Without RBC	y = no. seeks after 1st sector 1500 + 725 x + 1973 y
PLOT1 Test Print	130 678 +	Write With RBC Write Imm Seek 1 to-center By addr	2599 + 1450 x + 3947 y 1085 + 725 x + 1707 y 1871 2151

Figure 31. Execution Times of 1130 ISS (DM2 System)

ARITHMETIC AND FUNCTION SUBROUTINES

The execution times of the arithmetic and function subroutines are shown in Figure 32. All times are based on a $3.6-\mu$ sec instruction cycle; the times containing a decimal point are milliseconds, all other are microseconds.

SPIR (C/PT SYSTEM)

The SPIRx subroutines take 220 μ secs (3.6- μ sec instruction cycle) plus the DISKx time to read sector 0000.

STANDARD		EXTENDED	
FADD/FADDX]	460	EADD/EADDX]	440
FSUB/FSUBX	560	ESUB∕ESUBX ∫	490
FMPY/FMPYX	560	EMPY/EMPYX	790
FDIV/FDIVX	766	EDIV/EDIVX	2060
FLD/FLDX 】	180	ELD/ELDX	160
FSTO∕FSTOX ∫	180	ESTO∕ESTOX ∫	170
FLOAT	330	-	330
IF1X	140		140
NORM	260		260
FSBR/FSBRX	650	ESBR/ESBRX	740
FDVR/FDVRX	1090	EDVR/EDVRX	2520
SNR	80		80
FABS/FAVL	50	EABS/EAVL	60
IABS	100		100
FGETP	330	EGETP	320
FARC	60		60
XMDS	260		
FIXI/FIXIX	465		465
XSQR 550 av. (860	max.)	550 av. (860	max.)
XMD	520		520
XDD	1760		1760
FSIN/FSINE	3.0	ESIN/ESINE	5.4
FCOS/FCOSN	3.4	ECOS/ECOSN	5.9
FATAN/FATN	5.2	EATAN/EATN	8.9
FSQRT/FSQR	4.5	ESQRT/ESQR	10.4
FALOG/FLN	5.1	EALOG/ELN	8.0
FEXP/FXPN	2.0	EEXP/EXPN	4.4
FAXI/FAXIX	3.8	EAXI/EAXIX	4.7
FAXB/FAXBX	8.0	EAXB/EAXBX	13.3
FTANH/FTNH	4.3	ETANH/ETNH	8.1
FBTD (bin. to dec.)	40.0		40.0
FDTB (dec. to bin.)	20.0	1	20.0
		I	-

Figure 32. Execution Times of Arithmetic and Function Subroutines

Re-enterable Code

Re-enterable code is defined as cole that can be executed by more than one program at a time and that does not modify itself. Such code makes it possible for the programmer to write subroutines that can be called from more than one level of program operation; that is, from the mainline level (no interrupt) and an interrupt priority level or from two different interrupt priority levels. Two problem areas in writing re-enterable code are (1) obtaining temporary storage, and (2) modifying storage locations and/or instructions.

It is necessary, in this discussion of re-enterable code, to point out the following facts about the 1130 and its method of operation:

- Instructions have direct and indirect addressing. The operand of an instruction can address a location that contains either the value or the address of the location that contains the value to be addressed, multiplied, etc..
- Index registers occupy storage locations that can be addressed.
- Register housekeeping is performed for interrupts. IBM Disk Monitor interrupt-programming saves and restores the index registers, accumulator, and accumulator extension.
- Interrupts on same or lower level of priority are inhibited. Once the CPU has executed the hardware-forced branch for a level of interrupt priority, no hardware-forced branch for that level, or a lower, level can occur until the programmer exits from the level.
- Storage can be modified by a single instruction (MDX instruction) that cannot be interrupted.
- The subroutine call instruction (BSI instruction) is not re-enterable. The call instruction stores the return link (address of next instruction following the call) in a storage location. This return-link storage location cannot be varied by the subroutine. Therefore, a second call to the same subroutine stores the return link for the second call in the same location where the return link for the first call was

stored. (The subroutine-call instruction is also the instruction executed for the hardware-forced branch that initiates processing for a level of interrupt priority.)

- Index istructions cannot be indexed. The index instructions (load index, store index, modify index) cannot specify an index register to address the storage location from which the register is to be loaded or modified, or into which the register is to be stored.
- There are no register-to-register instructions. The index registers and accumulator must be loaded from, stored into, or modified from core storage.
- There is no indirect addressing for load-index and modify-index instructions. These two instructions have only immediate operands, and directly-addressed operands.
- There is no instruction to inhibit interrupts. There is no mask instruction to selectively or completely inhibit levels of interrupt, and no instruction to force an interrupt level on.

The definition of re-enterable code given earlier can be extended to include code that modifies itself as long as the modification does not affect the output of the code. Such an extension permits the code to be executed by more than one program at a time. Using this extended definition of re-enterable code, the remainder of this discussion illustrates how re-enterable code can be written for the 1130.

The Lisk Monitor, the Card/Paper Tape System, and their subroutines are <u>not</u> re-enterable. This does not prevent the user from writing his own re-enterable subroutines as long as these subroutines do not call, either directly or indirectly (for example, LINK), any Lisk Monitor or Card/Paper Tape subroutines.

For discussion purposes, there are two areas in writing re-enterable code: (1) getting to (calling) the code, and (2) writing the code. Assuming the existing assemblers and compilers, re-enterable code can only be written in assembler language. However, re-enterable subroutines may be called by either assembler or FORTRAN language programs as described below.

CALLING A RE-ENTERABLE SUBROUTINE

The subroutine calls (LIBF and CALL) cause the following BSI instructions to be generated:

	CALL	Generated Code	System		
	CALL subr	BSI I subr TV locat	Disk Monitor System,Version 2		
	CALL subr	BSI L subr	Card/Paper Tape System		
l	LIBF subr	BSI 3 subr TV disp	both systems		

For a re-enterable call, the suproutine call instruction should be preceded by another instruction which places the return link in a location saved and stored by interrupt programming, such as in an index register, the accumulator, or the accumulator extension. Through conventions agreed upon between the calling program and subroutine, the re-enterable subroutine called expects the return link to be in a pre-defined register and ignores the return link stored by the subroutine-call (branch) instruction. The added instruction in the calling sequence can be a load-index or load-accumulator instruction, or even a load-double (accumulator and accumulator extension) instruction. This combination (load instruction + subroutine call) gives the programmer a re-enterable call that can be used in 1130 programming.

The re-enterable call (two instructions) can be generated for the assembler user by (a) writing and then using a macro, or (b) by actually coding the two instructions. For example, if index register 2 is selected for the return link, the following code could be used:

Actual Coding or Macro						
LDX L2	*+2	r 	RCALL	subr		
CALL	subr	! 				

The RCALL macro is defined by the following code:

Labei 21	25		Opera 27	tion 30		F 32	1		35		40	45	Operan 50	- ds & Remort S					70
	-	٩	sm/	20			Γ	Γ					. 1		<u> </u>				
			R,C,A	٩L	Ł				z.				1 F		1				
			LDY	S.		í.	2	Ĺ		+2.		load	XF	2,2	v/r.	et.ur.	n, aa	d,r,	
		L	CAL	1					z,						1.1.1				
		Ľ	MEL	νD				L	L.			ب ب ب ب ب ب			-1 -1 -1				
					L			L			1				1		1.1.1		
		η.		-	E		Ł.		L	. 1					<u></u>				

The FORTRAN user must write a special subroutine in assembler language and then call that subroutine in FORTRAN. That subroutine is not re-enterable. Consequently, there must be a separate special subroutine for each level (mainline or interrupt) from which FORTRAN calls may be executed. To call re-enterable subroutine A with parameters X and Y, the FORTRAN user would name subroutine A in an EXTERNAL statement and then call special subroutine B with parameters A, X, and Y, in that order. Subroutine B would load the pre-defined register with the address of the location immediately following the A parameter and, using the A parameter, would call subroutine A. If subroutine A is called as follows:

CALL A (X, Y)

then this call can be replaced by the following code to obtain a re-enterable call:

EXTERNAL A



If index register 2 is selected for the return link, special subroutine B is defined as follows:

6	sbel			Operation	Π	F	7				-	_	_	_	_					00	eror	da i	s Re	soor	kı										_
21		25	Ē.	27 30		12	33	35				40				45				Ś	0			5	5			_	60			_	65	 _	70
		-		ENT.				8									_				1			+					<u> </u>				_		
B,	<i></i>		\Box	D.C.				÷.	- ,1	6			1	_			_	_			1				1			_	<u> </u>	۰.			. 1	 	
			÷	LDX.		1	2	в								X	R	2	≂.(a e	1,0	(,r	.	0,1	ſ.,	ζ,	.A	6	4		٩.				
				M.D.X.			2	a								х	R.	2	÷.1	ac	1 _i a	ŗ		0,1	Ċ.	e	a	d	Ľ.	د_	Ś.,		_	 	
				8. 5.C		Ż,	2	- 1	ē.				1.			e	2	<i>i</i> .,	t.,	1	10	ι.	2	Al			A		<u> </u>					 	
				E,ND						,											1				1							i		 	
L.			ų.					1					1				ب				L	-			1				_ 1				. 1		

The re-enterable calling sequence allows the return address stored by the call (BSI) to be modified by the interrupt without affecting subroutine operation since a re-enterable subroutine ignores the effective address (EA) location and uses the contents of XR2 as the return address.

OBTAINING TEMPORARY STORAGE

The temporary storage locations that are easy to use are the areas saved and restored by interrupt programming: index registers, accumulator, and accumulator extension. There are times when these are not adequate:

- When there are not enough registers
- When registers must be loaded with or modified by calculated values (variable rather than constant value)
- When registers must be loaded from, stored into, or modified by locations addressed via index registers

Work areas in storage may be assigned to each subroutine or program (common to many subroutines) to provide temporary storage for each level of operation. Such areas may be used for storage of intermediate results, parameters, data, calculated addresses, etc.. These areas may be accessed via index registers or address constants.

Current 1130 interrupt programming does not provide for level or program work areas. Such areas can be provided for each level by modifying the interrupt programming, by requiring locations in COMMON, or by other changes. However, it might be easier for the user to establish work areas within each re-enterable subroutine as it is written, rather than to modify already-written programs and systems.

A number (X) of subroutine work areas of length (N+1) can be defined: where X is the number of work area words needed for subroutine execution, and N is the number of interrupt levels. The subroutine increments the address of each area by 1 for each entry and decrements the addresses by 1 for each exit. Any instruction can then directly or indirectly reference the area.

Access via index register. If index register 2 is used to locate a 4-word area to be used for up to three concurrent entries, then word 1 might be used for the address of a parameter and word 2 for an intermediate value, as shown in the following code:

سر ا	. 1		Operation	Ł	ŀ	l,	1	1				Operands & Re	marks			
,	25		27 30	L	32	33	L	35	40		45	50	55	60	65	70
			MOX	Г	Ĺ	Γ	Γ	INS	TR+1	. 4.	inc.	addr.	work.	air.e	aut	<u></u>
INS	TR	Π	LDX	E	į,	2	I	WOF	KA-A	÷	XR2	c, c, u, r, r, e,	nti an	1.1.14.	.a.r.e.a.	الانتخبيت
					L		L				ب ب ب ب	،	ب سب الله	• • • •		الحمل
.			•	E	1	Γ.	L	1.				و و و و	ا الماليات		ــــــــــــــــــــــــــــــــــــــ	المصلية
—		Ē		Ľ			L	Lu				ويعرب فيترجي	ساست التروين			L
			LD.	Γ	L	2	I	2.		A	add,	value.	ulm)	tora	ddr.	المسمع
			A	L	Ľ	Г	Γ	הסא	DR.C.		_ <u></u>	s.t.a.n.t.	(AU			لاست
		ł	STO.	E	Γ	a	I	1.			SIG.V.	e, ,c ₁ a,1,c,	ul ate	d, ja,a	dr. 1	استنب
—		E.	LD.	E	1	2	I	1.			giet.	p.a.r.a.m	atier.		+ 4-4-4-4	المستحير
				Ε	Г	Г	T	1		And A		ويد ويت المريد			<u></u>	المسادية
		2		E	J	T	ľ	1.					حداجه		. <u></u>	المسلم
		F		2	1	Τ	E					<u></u>		ب الساب	· · · · · · · ·	استبي
		E	MD.X.	ľ	V	Т	T.	INS	ST.R.+	1	4, d,e,	c,r, w,o,r,	k are	a ae	l.d.r.	السديد
		F		Т	T	Т	F				e.x.	i.t. 1			- بــل بــــ	المسل
		l		F		T	T									· · · · · ·
		F		T	T	T	T		1							ا- بـــــ
		F		T	1	1	Ť									<u></u>
		ŧ.	D.C.	1	ł	t	I,				a.d.d.	rc.o.n.s	tant			
WO	RK.A	1	8.5.5.	1		T	l	34	4		WOT	k, ar.e.a				
		į,		1		T	1									
<u> </u>		ľ		1		T	l	A								• • • •

Access via Address Constants. Assume two words needed for each of three concurrent entries. Note that an address constant is required for each word. Index register 2 is used here to access call parameters rather than work area words.

	_		-		-	_						
Lobel		Operation	1	F	1			Operands &	Remarks			
21 25		27 30	L	32	33		35 40 45	50	55	60	65	70
	П		T		Γ							
		61100	r	t	t				A.1. 0			4
	Н	S.U.S.A	t.	-		-		1 anser	1910 A		macrio	, seld
	Н	****	H		H				L J. M	فاحقت فسنله ساده		
	-	e	Ļ.	⊢	-	ι.		· · · · · · · · · · · · · · · · · · ·			حادية معادية	
	Н	• · · ·	┢	⊢	+	-		يديد فيتبيب		а. а. а. А		
		MDX.	L	L			TEMPA, 2 in.c	r. add	r, wor	K. WO.F	d. 1.	
have dealers	1		Ŀ	L.	LI	Ł	1					
		MDX.	Ľ	4			58100+1, 2 inc	r. ad.d.	r. wio.r	K. WO.F	d. 2 .	
			Г									
			r.									
	Η		۲		t	1	──╀╌┼╌┼╌┼╌┽ ─┼╌┵╺┷╺┷╺┷	+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-		****		. است ک
	Н		t	ŀ		F		ي ي باب محمد	المطيعات	ي و به الم	• • • • • • •	ب ب
<u> </u>	H	SIX	┢	μ	1	-	ILMPA SITO	r.e. X.R.	L . L M.	WOLK.	w.o.r.d.	<u></u>
		LD.	L	1			TEMPAac.c	=, c, o, n, 1,	entis,	<u>XR1 .</u>	جيب	
<u> </u>	ί.	* I. I. I.	L									
	E	•	L									
			Г	Г	П	1					· · · · ·	<u> </u>
F		STO	t	7	t		SBI ØØ±1 ada		1 L.L.	1. i. i. i		<u></u>
60.00	H	<u>,,,,,,</u>	t	ť,	H	-		art that's	1, 1, 1, 1, 1V	S. 11	<u>o</u> , <i>r</i> , o , <i>r</i> , o , <i>r</i> , o , <i>r</i> , o , <i>r</i> , o , <i>r</i> , o , <i>r</i> , o , <i>r</i> , o , <i>r</i> , o , <i>r</i> , o , <i>r</i> , o , <i>r</i> , o , <i>r</i> , o , <i>r</i> , o , <i>r</i> , o , <i>r</i> , o , <i>r</i> , o , <i>r</i> , o , <i>r</i> , o , <i>r</i> , o , <i>r</i> , o , <i>r</i> , o , <i>r</i> , o , <i>r</i> , o , <i>r</i> , <i>s</i> , <i>r</i> , <i>s</i> , <i>r</i> , <i>s</i> , <i>r</i> , <i>s</i> , <i>r</i> , <i>s</i> , <i>r</i> , <i>s</i> , <i>r</i> , <i>s</i> , <i>r</i> , <i>s</i> , <i>r</i> , <i>s</i> , <i>r</i> , <i>s</i> , <i>r</i> , <i>s</i> , <i>r</i> , <i>s</i>	ــــــــــــــــــــــــــــــــــــــ
30,1,0,0	Н	L.L.	┝	4	+ +	⊢	THE MART TILL A HOUSE	<u>,≊,c,o,n,t</u> ,	antis.	LALL		
	Н	• • • •	┝	-	+	-	╶╋╼┺╼┺╼╬╍╬╸┢╸╘╶┶═┶╌┹╼╊╼┲╌	فالمتك فالمتكرية	· · · · · · · · · · · · · · · · · · ·			<u> </u>
J		e	L		1	Ļ.,		+ - - - - -		او	<u> </u>	حبب
	L	e	L				L					1
		L.D	F		2		Ø	ain 1	st. ca	11. 00	ramet	
	Π	S.L.A.	Г				8	1-145	t.i.f.u	2nd b	uto	
	П	STO	r	,		-	TEMPA Ser					****
	Н	5,70	h	ŕ	H	-	LENGALL SIGN	ਦਾ ਸਾਹਾਣਾ	LT TI G		e	
	Н	<u>μ</u>	÷		2	-	<u>1</u>	<u>a</u> ., <u>n</u> , 2,		l'i pa	ramet	e.r. 1
	μ	AND.	Ľ.			-	HOOFF.	late.	2nd_b	witler		
here	Ċ.	O.R.	Ľ	1		Ľ.,	T _I EMPA com	bine y	vith.	1,s.t. b	4, t.e.	1
	i											
	Π	• • • •	x			E.	1					
	Г		R									
	۲	MOY	Ľ	,	Н		TEMPA = 2 das		-1 1-1-1-1-1-1-1-1-1-1-1-1-1-1-1-1-1-1-	4		
	3	MOV	h	7	í	H		u		BI MOIL	a	
		in the second	t,	Þ	-	H		<u>c, r, a</u> ,a,a,	a,r, <u>w</u> o	r K. WO	r.a. ie	• • • • •
			۲	-	H	н	<u></u>	deter i v		<u></u>	<u></u>	<u> 1</u>
	Н	<u></u>	┡	-		_		دريسار والسبو	Sec. 1. 1. 1	ستدلد فالع		<u></u>
	1		L	L.,		H	L					
TEMPA		D.C					TEMR-A ada	IL OIL .	curre.	n.t. w.o.	r.k. w.o	r.d. ,1
T.E.MP.		8,S,S.	Ĺ				2,*3,	ce. fe	r3. A	ntrie	5	
HØØFF		D.C.	П				OOFF	k to	isola		abters	
		1	Н		Н	М			1-161.76		M DI MO	arr.
	Н		h	H	Η			≝⊶∔≁≁	<u></u>		· · · · · ·	
	ш	L	L		L.I	L.,	1. 1. 1. 1. 1. 1. <u>1. 1. 1. 1. 1. 1. 1. 1. 1. 1. 1. 1. 1. 1</u>	للم الم الم الم				المسعم

Note: Interrupts (and the subsequent call from interrupt processing) can occur after any instruction; therefore, each direct reference to the work area should address a different area. Otherwise, at any time, a sequence that should address a particular word in the work area can end up referencing different words in the work area and overlaying the contents of words used for a previous, and yet unfinished call. The following example avoids this problem:

Label	Ē	Operation	1	F	1.	Γ				Operands	A Remorks			
21 25		27 30	22	22	33	L	35	40	45	50	55	60	<u>65</u>	70
	14	MDX.		Ľ.	L	L	ALS.	$T_{1} + I_{1}$	1. inc.	r, w _i o,r	k, are	a, add	C	البياريا
<u></u>	L	••••••	L	L.	⊢	F	<u> </u>							
	L	e	H	L	┝	┢	1						LLLL	لاست
<u> </u>	Ļ	مبيده	H	Ł	-	Ļ	بيد	، ب ب		بملبب	يت الم	والمساهدة والمساد		لاحب
Later to the	L	<u>S.T.O.</u>	┢	Ľ	-	Ļ	RE.S.	T.1.+1.	s_a.v.	e, v,a,/	<u>u,e, , ,</u>	يتطاريك الم	_ السالي ال	المساسط
hur	2	• • • •	L	L	1	L	4							لاست
	Ļ		L	L	L	L				<u></u>	والمساحد المر	بببي	للبلب	لىب
	Ľ		L	L	L	L		سلسب					<u></u>	ل ــ
RE.S.T.1	L	LD.	L	Ł		L	WOR	KA1.	get.	v.a.l.u	e			للتمتيت
have	Ľ	••••	L	L	⊢	L								الدار الما
<u> </u>		• • • • •	Ļ	L	L	L			و المنظور الم				بالمتعالم المناسب	لمست
	L	e	H	-	┢	L	<u> </u>	، ا					يتد المعدي	المعمد
RESTZ	Ļ	LD.	L	μ	L	Ł	RE S	T.L.t.L.	gie.t.	valu	<u>а. </u>			ليتنت
	L		L	L		F		والمتحاد والم				ب ال ال ال		لىت
1	μ		Ŀ	1		l	1			$(\mathbf{r}_{1},\mathbf{r}_{2},\mathbf{r}_{3})$	с к. н . к	$\mathbf{r}_{i} \in \{1, \dots, n\}$	•	•

MODIFYING STORAGE OR INSTRUCTIONS

Storage and/or instructions can be modified in a re-enterable subroutine if the sequence shown in Figure 32.1 is used.

- Save the location or instruction to be modified in temporary storage (work area, index register, accumulator, or accumulator extension)
- 2. Modify the location or instruction
- 3. Execute, using the modified location or instruction
- 4. Restore the location or instruction from temporary storage

An example best illustrates the techniques of temporary storage and storage/instruction modification. The



Figure 32.1 Modifying Storage or Instructions

coding examples below illustrate four ways of loading index register 1 with the address of table D, assuming the following address chain and that only location A is directly accessible.

L cha	1	Г	Overation	Г	F	-	Г	Γ	Operands & Remarks
21	25		27 30		32	33	Ł	35	s 40 45 50 55 60 65 <u>70</u>
A			DC.		Г	Г		a	• • • • • • • • • • • • • • • • • • •
		Г		Γ			Г		
		Γ	• • • •	E	Г	Г		Ι,	1
				Г		Г	Γ	Γ,	
B		Г	DC .	Г	Г	Г	Г	Ic.	· · · · · · · · · · · · · · · · · · ·
		Г	DC	T	Г	Γ	Г	6	
		r	DC.	F	F	F	F	E	
		r		Г	T	Г	r	Γ.	
		F		Г	T	T	r	1.	
		۲	<u> </u>	r	t	r	t	17	<u></u>
<u> </u>		r	866	t	t	t	t	5	
and the second		t	0,57,57	t	t	t	t	ť.	
<u> </u>		ŀ	· · · ·-	t	t-	t	t	Ľ	
<u> </u>	· · · ·	ł	• • - • · · · ·	t	┢╌	┢	t	۲	<u></u>
	• • •	┝	<u>.</u>	⊦	ł	┢╍	ł	1.1	<u></u>
$\rho_{}$		┝	8.55	┝	╀	⊢	⊦	r	<u> </u>
.		┡	••••	⊦	┞	┢	₽	≁	
<u> </u>	<u> </u>	┡	ــــــــــــــــــــــــــــــــــــــ	Ļ.	⊢	╞	┝	≁	<u>▋_▋▋↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓ </u>
<u> </u>		Ļ	<u></u>	┡	⊢	1	┡	1	<u></u>
pøø	Ø,1	F	p.c.	Ļ	┡	Ļ	Ł	Р	<u>, , , , , , , , , , , , , , , , , , , </u>
L		L		L		L_	L	, ا	والمتعادية المتحديث فتعصينا فتتعاد والمتحد والمتحد والمتحد والمحاد والمحاد والمحاد والمحاد والمحاد والمحاد والم

The load-index instruction has one less level of indirect addressing than the load-accumulator instruction. Using the symbol A as an operand, an indirect load-accumulator instruction can obtain the address of C, while an indirect load-index instruction can only obtain the address of B. The coding examples illustrated below show that using work area words is the most expensive method, both in number of instructions required, and in the time it takes to execute those instructions. The technique selected depends on the temporary storage available at the time (accumulator extension, or second index register, or neither).

Example 1. Using work area words:

labe!	F	Operation	Γ	F	7	П				Operands & Ren	ncaria ș			
21 25	Ł.	27 30		32	33		3540		45	50	55	60	65	70
	Γ	MOX.		Ł	-		1.N.S.T.1.+1	u.t.	inc.r.	addr.	WO.r.K.	area	A	ليستيات
	E	MOX.		L			1,NS,T,2,+1	int.	inc.r.	add,r,	WO, r.K.	ar,ea	B	أسلحاه السلي
	Γ							Ĺ					1.1	ل ب ب ب
	L	•••••	Ľ	L		L		<u></u>				بتنقد		
	[E								لستنت
	Γ	LA	E	L			A		a10, c, =, c	addr, I	<u>a</u>			المتعالية
	L	A					00001.		سبب	بيبات				المعال
		5.7.0	L	1		Ľ	1,NST,1,+1	here	,q1C,C,=10	dd,r, E	<u>3+1/</u>			لينتب
1.NS.T.1	Γ	LD.	Г	1		Γ	WORKA-1	L	accipie	ddr. 1	2	سيب		للعبابات
	Г	S.T.O.	Γ	1/	Г	E	1.N.S.T.2.+1							
INST2	L	LOX.	E	1	1	Γ	WORKE,-1		inde,	(, <u>(</u> , e.g.	,1,=,a,d,d	dr. D.		لسيسب
	E		Γ	Γ		Ľ	4					وللم المراجع	11	ي ب ب
	E		5	Ţ	L	E		- I - I - I		L. L. L. L.			<u>د</u> د.	ىبىت
	Γ		Γ			Ε			، ب ب ب					
	E	MDX.	Γ	4	I	Γ.	INST.1.+1	i	1, d,e,c,i	. <u>a</u> ,d,d,	r, w.o.r.	k, _i a,r,e,	a. 1	4
	E	M.D.X.	Γ	L		L	1NST2.+1	1	deci	. a.d.d.	r. w.o.r.	k gree	a	3
	Г		r	Γ	Γ	Γ.	1						<u></u>	المتناب
	Г		E	T	Γ	Г	1							
	F		F	1	T	Г	1				1 1			
WORKA	Ľ	8,55	ſ	Г	Г	Г	n_{1}		work.	ar ea	1			ليبين
WORKE	Ľ	8,5,5,	ſ	1	Γ	Г	n		WO, r.K.	a,r.e.a.	2			
	_			_	***	_								

Example 3. Using a second index register:

	_		_	_	_	_								
tabel .	1	Operation		F	l٦	Ľ	l			Operands &	Remark1			- 1
25	r.	27 30	L	32	33		35	40	45	50	55	60	65	
		LOX.	Г	7	2	T	INSTI	dil	Save	<u>, cıo, n</u> ı	tents	Let 1	<u>NST11+1</u>	in in
	1		Г		Г	Γ	1		XR2	سبي		حداث ب		
	_	LD.	Γ				A		d1C,C=.0	a,d,d,r	B	ب ا		
	Г	A	Г		Γ	L	00001		مبيار		ee			· · · · -
	Γ	5,7.0.	Γ		Γ	Ľ	1.N.S.T.1	+1, 1, 1, 1	_a,c,c,=,	<u>a,d,d,r</u> ,	B+11	المار		4- 4-4 -7
INST.1	Г	LOX		/	1		#1#.		inde.	<u>x, ;r,e</u> ,	معراب و	ddr_C	يبعب	لمعب
· · · · ·	C	ST.X.	Γ		N	1	INST	<u>1,+1</u>	<u>,r,e,s,t,</u>	<u>o,r,e</u> ,	c.o.n.t.e	<u>n, f s</u> 0	f. 1 N.S.T	БĿ±Ц
	F.		Г	Ľ	Г	Т					الم الم الم	<u></u>		لمعبب

Example 4. Taking advantage of the fact that 1130 index register occupy storage locations that can be addressed:

Lobel	Γ	Operation	Γ	F	T	Π			Operands i	Remarks			
2125		27 30		32	33		35 40	45	50	55	60	65	70
		LD.					A		a,a,d,d,r	ALL			
		S,T,O,		L			1	XR.	$l_{i=a,d_id_ir}$	BL			Later I
		LO			1		ta a contra	<u></u>	s,a,a,d,d,r				
		S,T,O,		4			1		1,=,a,d,d,r				
		Luci		L							1.1.1.1.1.		

1800 COMPATIBILITY

Each MDX instruction used to increment or decrement addresses must be immediately followed by a NOP instruction because of the skip that occurs if the addresses cross the 32K boundary (positive value less than 32K, negative value equal to or greater than 32K. The example 4 technique of modifying storage and/or instructions cannot be written since the 1800 index registers do not occupy addressable storage locations.

Example 2.	Using	the	accumulator
extension:			

Lobel	Operation	П	F	1	1			Operands	& Remarks			
21 25	27 30	b	2	33	35	40	45	50	55	60	65	70
	La.	П	l		LING	S.T.1.+11	,	, c.e.n.t	entis	0.1.11	VS.T.1+1	<u> </u>
	RTE	Ц	7	\downarrow	16		Sigiv	e, ,110,	a.c.c)e	and in the	ــــــــــــــــــــــــــــــــــــــ	دىت
	40.	L			A.		, <u>a</u> icic	=,a,d,d,r	·			الرب ب
	A	П	1		DØ.	0.0.1		حال ب		-		المسل
	S, T,O,	П			I.N.	<u>s, T, 1, +, 1</u>	. <u>a</u> ic.c	,⇒, <u>a,d,d,r</u>	B+1			ليبيب
INST1	LOX	E1/	1	1	*	**	, , ind	ex, re	.g. 1 = 0	dd,r	2	الم ال
Land Land	RT.E.	П			1,6		<u></u>			سبب	ب ب ب ب	بب
	S.T.O.	Ш	1	1	L.N.	ST.1.+11.	, rje s	tore.	, c, o, n, t, e	<u>, n, t, s, </u> , o	of, INS	$T_1 + 1$
		IΤ	I	Т	1							

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